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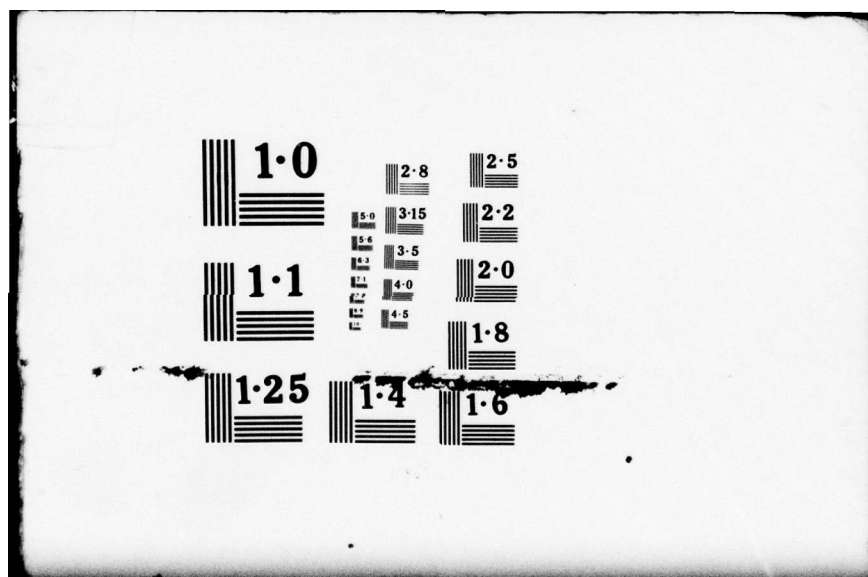
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THE EFFECTS OF FINITE WORDLENGTH.

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THESIS 17

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Master's thesis,

11
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12
185p.

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SOFTWARE TOOL(S) FOR EVALUATING THE
EFFECTS OF FINITE WORDLENGTH

THESIS

Presented to the Faculty of the School of Engineering
of the Air Force Institute of Technology

Air University

in Partial Fulfillment of the
Requirements for the Degree of

Master of Science

by

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Captain

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Graduate Electrical Engineering

December 1977

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Preface

The n-bit simulation tool has been developed with the hopes that it can be of some help to engineers in the evaluation of the effects of varying wordlength upon their algorithms.

I would like to thank Professor Richard and Mr Stan Musick for their comments and help during the design phases of the n-bit simulation tool. Also, I owe a great deal of thanks to my two advisors, Dr Peter Maybeck and Captain J B Peterson, for their on-going encouragement and advice throughout this period. They were always there when I needed them. Lastly, I would especially like to thank my wife, Karen, who was also my typist. I would never have made it without you. Thank you for your patience, understanding, support, and your endless flow of love, not to mention your outstanding work.

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Abstract

↓ In the development of estimation and control systems, the algorithms developed are, in general, very sensitive to the wordlength of the onboard computer. The representation of the desired design can be greatly affected by the quantization resulting from truncation or round off. The result can be even more severe than numerical imprecision: numerical instability can be generated that render algorithms totally useless.

An important part of the algorithm development is the determination of the appropriate wordlength. The proper wordlength can be determined by running the algorithm on a simulated n-bit machine for several values of n. The results can then be evaluated against the performance specification to determine the wordlength and precision requirements.

The objective of this thesis was to process algorithms written in FORTRAN on the Control Data Corporation (CDC) 6600/CYBER 74 computer systems such that the results obtained were similar to those obtainable on an n-bit machine. The problem of n-bit simulation was addressed from two levels: the assembly language level and the FORTRAN language level. Each level involved designing a preprocessor, which would modify the algorithm's code so that the numerical effects of n-bit wordlength could be realized. The assembly

(continued on p 1x)

(cont. fr p viii)

language (COMPASS) level approach failed, however, an n-bit simulation tool was successfully developed and implemented at the FORTRAN level. Several user options were incorporated into the n-bit simulation tool to enable the user to simulate the characteristics of various computer types.

I. Introduction

In the development of estimation and control systems, the algorithms developed are, in general, very sensitive to the wordlength of the onboard computer. The representation of the desired design can be greatly affected by the quantization resulting from truncation or round off. The result can be even more severe than numerical imprecision: numerical instabilities can be generated that render algorithms totally useless.

An important part of the algorithm development is the determination of the appropriate wordlength. The proper wordlength can be determined by running the algorithm on a simulated n -bit machine for several values of n . The results can then be evaluated against the performance specification to determine the wordlength and precision requirements.

The objective of this thesis is to process algorithms written in FORTRAN on the Control Data Corporation (CDC) 6600/CYBER 74 computer systems such that the results obtained are similar to those obtainable on an n -bit machine. Then the results would reflect the degree of accuracy that would have resulted from actually processing the algorithm on a computer with that n -bit wordlength. For the purposes of this report, n -bit simulation will refer to the effect n -bit wordlength would have upon the numerical accuracy of an algorithm being processed on an n -bit wordlength machine.

This type of tool would be useful in the initial stages of the selection process for an onboard computer that would process the desired algorithms. It is essential to know the full impact of the wordlength of a computer upon the algorithms before selecting the computer.

For a hypothetical example, suppose the Air Force just purchased several mini-computers having 16-bit wordlengths. It was thereafter discovered that these computers did not furnish the degree of accuracy required by the algorithms. Considerable time may be expended rewriting the algorithms to process properly on the purchased computers, if that is at all possible. In fact, it may be the case that the minimum wordlength which could meet the requirements would be 24-bits or that a 16-bit wordlength computer with double precision capabilities would be required. Application of the n-bit simulation tool before purchase could have saved the government significant time and money.

Similarly, at the other end of the scale, this software tool could save government expenditures for computers with much longer wordlengths than the algorithms could ever require. Computer costs go up rapidly with longer wordlengths. Tailoring of the computer wordlength to the algorithm requirements could therefore eliminate some unnecessary expenditures.

Through the utilization of the n-bit simulation tool, an algorithm can be processed with several different wordlengths. The results can be compared to results generated

by the full 60-bit wordlength (which is longer than any foreseeable onboard computer wordlength) of the CDC 6600/CYBER 74 computer systems. The degree of accuracy lost due to each n-bit wordlength can then be determined.

Example of Its Application

Figure 1 can be used to demonstrate a typical application of the n-bit simulation tool. The Kalman Filter represents the algorithms which would be processed by the onboard computer. However, this procedure did not correctly reflect the degree of accuracy the onboard computer (having a shorter wordlength) would achieve. The original algorithm (programed in FORTRAN) would be modified by the n-bit simulation tool

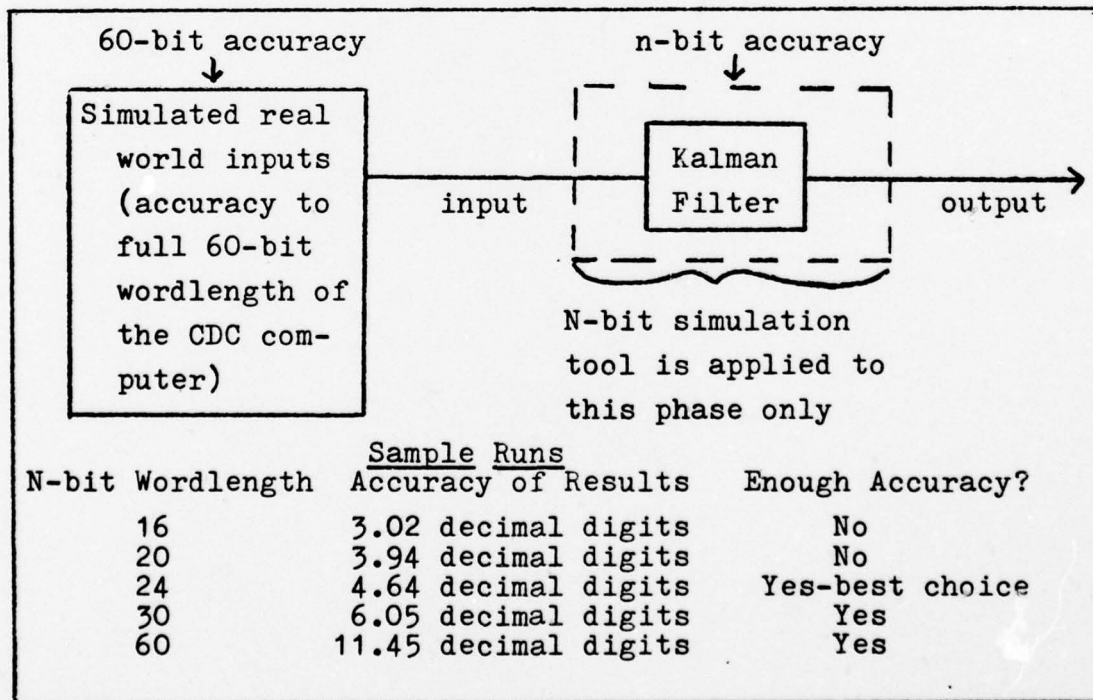


Figure 1 N-bit Simulator Application

so that its results would show the exact amount of significance that an n-bit wordlength computer would have produced. It may be executed several more times with varying wordlengths until sufficient information has been gained.

Requirements

The requirement for this thesis is to develop an n-bit simulation tool which when applied to a FORTRAN programed algorithm will produce the near exact arithmetic accuracy as a computer of the specified n-bit wordlength. Five options were considered necessary to provide the user with a versatile n-bit simulator which could be used to simulate the characteristics of various computer types:

- 1) Allow the user to specify n, the number of bits per wordlength,
- 2) Permit specification of floating point word characteristics:
 - a) The number of bits to be used for the exponent and mantissa
 - b) Position of the implied decimal point with respect to the mantissa, either to its left (making the mantissa a fractional representation) or to its right (making the mantissa a fixed point representation),
- 3) Give the user an option to have arithmetic operations performed with either truncation or rounding effects,

- 4) Allow the user to specify whether arithmetic calculations are performed in single, double, or triple wordlength precision with the results being stored as single precision quantities. This option essentially simulates the effects of a computer having a single, double, or triple wordlength accumulator but only single precision storage variables, and
- 5) Allow the user to specify different portions of an algorithm for different wordlengths.

Another characteristic of the n-bit simulation tool is that when an overflow is detected for a given n-bit wordlength, the maximum value that can be represented with the n-bit word should be substituted for the overflowed word and an overflow message printed. Overflows are treated differently depending upon the computer type, but it was felt maximum number replacement for overflows would better demonstrate the numerical effects of shorter wordlength upon an algorithm than simply truncating the most significant bits of the value that overflowed. When a word's value exceeds the maximum positive number the word could contain, the maximum positive number would be substituted into the word. If the word's value exceeds the maximum negative number, the maximum negative number would be substituted into the word.

Some general goals of this thesis were to provide the user with a n-bit simulation tool that would be reliable, be easy to use, be maintainable, provide an option to communicate overflow error messages when they occur, and to employee

structured design techniques in programing so the resulting design can attain certain desirable characteristics.

Assumptions

Some fundamental assumptions upon which this thesis was based are listed below:

- 1) Applying n-bit simulation effects to all arithmetic expressions is sufficient to simulate an entire program effectively in n-bit precision. The major purpose of this thesis is to provide a tool which can be used to evaluate the effects that varying wordlengths have upon arithmetic accuracy of an algorithm.
- 2) The algorithms which are to be n-bit simulated, are programable in American National Standards Institute (ANSI) FORTRAN and are executable on the CDC 6600 or CDC CYBER 74 computer systems. Potential users currently use the FORTRAN language and execute programs on these computer systems regularly.
- 3) The maximum fixed point wordlength desired to be simulated is 48 bits. This is the largest fixed point representation the CDC system can handle. In addition, the largest exponent and mantissa that would be required are 11 and 48 respectively. These are also maximums for CDC single precision floating point representations. Within the foreseeable future, onboard computers would not exceed any of these maximums.

- 4) Additional core storage and execution time required for n-bit simulation will not be a limiting factor in the design of the n-bit simulation tool. This is partially true because this n-bit simulation tool is expected to be used only several times each year by any projected user.
- 5) The quantities to be n-bit simulated are single precision numeric quantities (i.e. no logical, alphanumeric, or double precision variables and quantities are used within the program being simulated). Again, this is because this thesis is addressing the problem of numerical precision and the stability of mathematical algorithms.
- 6) All floating point values to be n-bit simulated will have three characteristics in common with CDC floating point representations:
 - a) Mantissa will be normalized (i.e. the most significant bit of the mantissa is always a one, except in the case of zero).
 - b) Exponents will represent powers of two, which would be multiplied with the mantissa to obtain the floating point value of the word, for instance, not by powers of 16 for hexadecimal machines.
 - c) Computers being simulated represent negative numbers with 1's complement representation, where a negative number is represented by a 1

in the sign bit and the 1's complement of its magnitude in the other bit positions. Floating point exponents for 2's complement representation have a positive exponent range of one less than the negative exponent (for instance, 2^{-128} to 2^{+127} , whereas 1's complement would range from 2^{-127} to 2^{+127}).

General Approach

This thesis study addressed the n-bit simulation problem at two levels. Each level involved designing a preprocessor which would modify the algorithm's code so that the numerical effects of the n-bit wordlength could be realized. The first approach attempted to accomplish n-bit simulation through modification of the COMPASS assembly language representation produced from a FORTRAN compilation of the algorithm. The COMPASS code would be modified by an n-bit simulation preprocessor so that every arithmetic operation would be replaced by COMPASS code sequence which, when executed, would generate the exact accuracy attainable on a computer of the n-bit wordlength specified by the user. This approach, however, was unsuccessful in meeting the requirements.

The second level approach was able to perform n-bit wordlength affects upon a FORTRAN algorithm by modification of the FORTRAN code. This modification of the FORTRAN code requires a preprocessor which analyses the arithmetic statements within a program and replaces them with subroutine

calls. The subroutines perform the arithmetic operations while simulating n-bit wordlength affects. This approach proved successful in solving the n-bit simulation problem and is discussed in detail in Chapter 3.

Chapter Synopsis

Chapter 2 of this thesis report contains a summary of the analysis and design accomplished at the COMPASS level. A successful prototype of the n-bit simulation tool could not be achieved at the COMPASS level. The reasons for the failure are elaborated in Chapter 2. A working prototype of the n-bit simulator was achieved by way of the FORTRAN level approach. This approach and the resulting system design are discussed in Chapter 3. Chapter 4 briefly summarizes the effects that the Structured Design techniques had upon the system design. The next chapter summarizes the testing performed to test, debug, and validate the n-bit simulation system. Finally, Chapter 6 draws some conclusions and makes recommendations concerning the n-bit simulation tool.

II. COMPASS Analysis And Design

The first approach to solving the n-bit simulation problem was at the COMPASS assembly language level. This approach attempted to accomplish the numerical effects of n-bit wordlength upon an algorithm through modification of COMPASS code. The first section of this chapter explains why the COMPASS level approach was undertaken prior to the FORTRAN level approach. The following sections describe an analysis of the COMPASS code (produced by FORTRAN compilation) with respect to the n-bit simulation problem. Finally, the approach taken to solve the n-bit simulation problem and the major problems encountered which resulted in the failure of the COMPASS level approach are described.

The COMPASS level approach was anticipated to be less complicated than the FORTRAN level approach to the n-bit simulation problem for a number of reasons. In a single FORTRAN language statement several types of arithmetic operations and several data types may be mixed together in many different ways. Each data type would have to be handled differently, and the precedence by which the arithmetic operations are supposed to be performed would require further analysis and handling. Also, arithmetic expressions may occur in almost any type of FORTRAN statement, which could seriously complicate n-bit simulation performance.

On the other hand, at the COMPASS level, each arithmetic operation could be accomplished by a single instruction

involving only equivalent data types. The arithmetic precedence order of execution and conversion of data types has already been accomplished by the FORTRAN compiler. Furthermore, in contrast to many other assembly languages, COMPASS does not include base addressing and relative addressing. This greatly simplifies the insertion of additional COMPASS instructions into the original code, which would enable each arithmetic instruction to be performed while simulating the n-bit wordlength effects.

N-bit Simulation Accomplishment

N-bit simulation for an entire COMPASS coded algorithm can be achieved by simulating the effects of n-bit wordlength for each arithmetic instruction. A typical COMPASS fixed point arithmetic instruction is IX7 X3+X4. This instruction adds the contents of register X3 to the contents of register X4 and stores their sum into register X7. To simulate the effects of n-bit wordlength, the contents of registers X3 and X4 would be modified to contain the equivalent n-bit wordlength values (assuming they have not already). Then the two registers would be added together and the resulting contents of the X7 register would be similarly converted to the n-bit equivalent.

Fixed Point N-bit Simulation

N-bit simulation is accomplished differently for fixed point and floating point number representations. The range of fixed point values which could be represented is limited

Floating Point N-bit Simulation

To n-bit simulate a floating point word type, the correct number of significant bits for the mantissa must be maintained and the exponent should not be allowed to exceed the equivalent exponent range representable in an n-bit floating point word. CDC floating point data types have normalized mantissa. Therefore, in order to maintain the specified number of significant mantissa bits, only the number of bits specified for it would be retained immediately to the right of the least significant exponent bit (bit #12). This can be accomplished by right shifting (with sign extension) the contents of the word followed by a circular left shift. Right shifting with sign extension displaces the contents of a word X positions to the right, where X is the number of bits computed from subtracting the specified number of mantissa bits from 48. Forty-eight is the length of the CDC 60-bit floating point word mantissa. Those bits shifted off the right-end are lost. All positions vacated on the left by the right shift are replaced with the value of the sign bit. Figure 3 (a) shows an example of a 60-bit floating point word value and (b) shows the resulting contents of (a) after a right-shift of 40-bits (for an 8 bit mantissa). In a circular left-shift, bits shifted off the left end of the word replace those from the right end. In this manner, the significant bits that had been shifted off the right end from the original word are replaced with zeroes for positive numbers and ones for negative numbers. Therefore, for this

(a)	1721652525237366553
(b)	1721652525000000000

Figure 4 N-bit Decimal Representation

the value left is equal to 3.32971; a significance of .00029 has been lost in the 15-bit mantissa representation.

Overflow of the exponent could be detected by comparing the maximum/minimum values possible for a floating point word with the specified bit lengths for the exponent and mantissa. Again overflow detection would result in the maximum/minimum value being stored over the old value. In addition to the case where the exponent gets too large, the case of exponent underflow must also be addressed. This could occur when a value is computed which is so small that its most significant digit is not within the range of the specified exponent (although it is within the CDC's large range). For example, 5 bits may be specified for the n-bit floating point exponent and 6 bits for the mantissa. This exponent could represent an exponent range of ± 16 . Therefore, the maximum values representable with that floating point word would be ± 63 (the maximum mantissa value) times 2^{+16} if the implied decimal point is to the right of the least significant mantissa bit.

If the implied decimal point is to the left of the mantissa, the maximum would be .984375 times 2^{+16} (shown in Figure 5). The minimum significance representable would be ± 32 times 2^{-16} and $\pm .5$ times 2^{-16} respectively for the two

different decimal point positions. Underflow for the 12-bit floating point word would occur if a 60-bit division resulted in a quotient less than 2^{-24} . This value would not be within the range representable by a word with a 5 bit exponent. Zero would replace the value that overflowed before the program executed its next operator.

$0111111111111111_A = .984375 * 2^{+16}$ $0111111111111111_A = 63 * 2^{+16}$ $000000100000_A = .5 * 2^{-16}$ $000000100000_A = 32 * 2^{-16}$
--

Figure 5 N-bit Floating Point Value Range for n Equals 12

Now that the strategy has been defined for simulating the effects of n-bit wordlength for arithmetic instructions and the fixed or floating point word values involved, the following section will describe the arithmetic instructions which the n-bit simulation would be performed upon.

The Arithmetic Instructions

Many COMPASS code listings, resulting from the three FORTRAN optimizing compilers, were examined to investigate the handling of FORTRAN assignment statements and arithmetic expressions. The arithmetic operations addition, subtraction, multiplication, and division were found to be handled by the COMPASS instructions in Table I, where X_i , X_j , and X_k may be any combination of the eight X registers ($X_1 - X_8$).

Operation type	COMPASS Instruction		Octal code for operation type
	COMPASS mnemonic	Address field	
fixed pt add	IXi	$X_j + X_k$	36
floating pt add	FXi	$X_j + X_k$	30
fixed pt subtract	IXi	$X_j - X_k$	37
floating pt subtract	FXi	$X_j - X_k$	31
fixed pt multiply	DXi	$X_j * X_k$	42
floating pt multiply	FXi	$X_j * X_k$	40
fixed pt divide	*FXi	X_j / X_k	44
floating pt divide	FXi	X_j / X_k	44
*fixed point values in registers X_j and X_k are converted to their floating point equivalents before the division is performed. The quotient X_i is then converted back to a fixed point representation.			

Table I Fixed And Floating Point Arithmetic Instructions

Each fixed and floating point single precision arithmetic operation generated by the FORTRAN compiler was performed by a unique instruction with the exception of division. As noted in Table I, fixed point division is performed by a sequence of instructions. The divisor and dividend are converted to their floating point equivalents before a floating point divide is performed. The quotient then is converted back to its fixed point equivalent representation. This sequence of instructions for the fixed point quotient $X_1 = X_2/X_3$ is shown below:

PX2	B0,X2	[Pack X2]	} Converting X2 and X3 to floating point
PX3	B0,X3	[Pack X3]	
NX0	B0,X3	[Normalize X3]	
FX1	X2/X0	[Divide]	} Converting result back to fixed point
UX1	B7,X1	[Unpack quotient]	
LX1	B7,X1	[Shift to integer position]	

Since fixed point and floating point divisions share a common instruction ($FXi Xj/Xk$) and each type requires a different method to simulate the effects of n-bit wordlength, a choice would have to be made not to allow one or the other of the division types to occur at the FORTRAN program level. Floating point was expected to be the more common usage for algorithms, so fixed point divisions would be restricted from n-bit simulated algorithms. If the fixed point division sequence always appeared as consecutive instructions, fixed and floating point division could be differentiated; but the FORTRAN optimizing compilers order the COMPASS instructions for the most efficient execution, and consequently the COMPASS code generated for a sequence of FORTRAN statements may be completely intermixed.

One other potential problem could occur for the fixed point multiplication instruction $DXi Xj*Xk$ shown in Table I. $DXi Xj*Xk$ is actually the COMPASS double precision multiply instruction. However, this conflict of usage would not occur since the algorithms for the n-bit simulator would not require double precision FORTRAN variables (Assumption 5 states this in Chapter I).

Now that the COMPASS instructions, which perform the arithmetic computations have been described, two methods of incorporating modifications for each arithmetic instruction will be described in the following section.

Two Approaches to Incorporating N-bit Simulation

Two fundamental approaches to implementation of the n-bit simulation modifications were considered. The first approach would replace each arithmetic instruction with subroutine calls. The second approach would substitute, in-line, the necessary instructions to perform the n-bit simulation effects.

The subroutine approach would require an argument list to accompany each subroutine call indicating the registers involved with that instruction. This argument list would key the handling subroutine to manipulate the designated registers. Uniquely identifying each instruction/register combination with a separate subroutine would require 512 different subroutines per arithmetic instruction type (three registers are used for each arithmetic instruction and each register may be any one of eight possible X registers, making 8^3 or 512 possible combinations). Clearly that would have been too many routines. So an argument list would be necessary indicating the registers involved, and a technique would have to be devised in each subroutine to handle them correctly.

In contrast, the in-line code approach would be simpler and save subroutine linkage time, but would require more core space. If branching would be necessary within the substituted code sequence, some method would have to be devised to provide unique labels within each substituted sequence. For the initial COMPASS level investigation, the in-line approach was chosen.

Methods for Isolating Arithmetic Instructions

In both cases previously stated, the first obstacle was trying to isolate arithmetic operation instructions of Table I from the rest of the COMPASS code. Three techniques were considered.

The first technique involved examining the raw COMPASS object code (machine language). This would be a complicated procedure due to problems involved distinguishing data from instructions. Also labels, already assigned to fixed locations, are transparent from their associated object code. Injecting additional COMPASS object code to perform n-bit simulation would lead to misalignment of labels with respect to their associated object code locations. When examining the object code, a table containing all possible COMPASS instruction codes and the associated instruction length would have to be maintained. In addition, one must account for no-operation instructions (NOOPS). NOOPS are used to fill in remaining portions of 60-bit words which are not completely filled by instructions due to labels or other conditions. It would be a cumbersome task indeed. Consequently, this technique was quickly ruled out.

The second technique involved examining the COMPASS source code generated from the FORTRAN compilation of the algorithm. Distinguishing the arithmetic operation instructions would require a table lookup of each COMPASS code mnemonic. When matches occurred, a sequence of instructions would be inserted to perform the n-bit wordlength effects

for that instruction type. A significant advantage of this method is that the COMPASS assembler would do the work of realigning all COMPASS object code and label pointers. COMPASS source code would be accessed by using the E option on the FTN control card. The sequence of control cards in Figure 6 could be used to access the COMPASS source code generated from a FORTRAN compilation. The execution of NBIT (in Figure 6) would modify the original COMPASS source code (the COMP file) for n-bit simulation producing the revised n-bit simulated code in file COMP2. COMP2 would then be assembled and executed.

```
FTN, E = COMP.  
NBIT, COMP,,, COMP2.  
REWIND, COMP2.  
COMPASS, I = COMP2, S = FTNMAC, B = BIN.  
LDSET (LIB = FORTRAN/SYSIO).  
BIN.
```

Figure 6 N-bit Simulation Control Card Sequence

The third method for isolating and handling the arithmetic instructions would take advantage of some sophisticated features of the COMPASS assembler called operation definitions (OPDEFs) and IF conditionals (IFCs). Both are pseudo instruction types. The OPDEF consists of a sequence of COMPASS source code that is saved and assembled when an instruction type within a given routine matches the syntax specification. It usually includes parameters to be substituted for formal parameters, so that the object code generated can vary with

each assembly of the definition and whatever combination of registers that are involved in the instruction (REF 3: 5-27). Using this feature, the assembler would do the work by replacing the arithmetic instructions with the n-bit simulation sequence.

Figure 7 shows a COMPASS routine in which an OPDEF was inserted, then assembled with the rest of the COMPASS code. This particular OPDEF syntactically defines the general instruction type $IX_i X_j + X_k$. P1, P2, and P3 used within the OPDEF would represent the registers X_i , X_j , and X_k respectively. In the sequence of original COMPASS code, the $IX7 X0 + X4$ instruction would match this syntax and is replaced by this sequence when it is assembled as shown in the after assembly code sequence. After assembly, the object code (shown on the left) is produced by the COMPASS assembler upon detection of the $IX7 X0 + X4$ instruction. It can be seen after the assembly that the OPDEF for the fixed point add instruction replaced the original one.

Two additional pseudo instructions were also used in this example: SET (REF 3: 4-36) and CPOP (REF 3: 6-7). CPOP was used to disguise the fixed point addition instruction from its own OPDEF. Otherwise, if an instruction were used within its own OPDEF definition it would result in the assembler getting caught in an infinite loop of replacing an instruction with itself, which triggers the process again, and so on.

Location	Code Generated	Label	Operation	Variable
<hr/>				
(Before assembly)			IDENT	TEST
Inserted into original COMFASS code	NBITS	IXX+X	SET	16
			OPDEF	P1,P2,P3
	XAA+A		LX.P2	59-NBITS
			LX.P3	59-NBITS
			AX.P2	59-NBITS
			AX.P3	59-NBITS
			CPOF	0,360B,132B
			XA,P1	A.P1+A.P3
			LX.P1	59-NBITS
			AX.P1	59-NBITS
			ENDM	
Original COM- PASS code from FORTRAN compi- lation			SA5	A0
			SX7	116100B
			IX2	X3-X4
			IX7	X0+X4
			DX7	X4*X4
END				
<hr/>				
(After assembly)			IDENT	TEST
0 54500 7170116100 36234	NBITS	IXX+X	SET	16
			OPDEF	P1,P2,P3
	XAA+A		LX.P2	59-NBITS
			LX.P3	59-NBITS
			AX.P2	59-NBITS
			AX.P3	59-NBITS
			CPOP	0,360B,132B
			XA.P1	A.P1+A.P3
			LX.P1	59-NBITS
			AX.P1	59-NBITS
			ENDM	
1 20053 20453 21053 21453	XAA+A		SA5	A0
			SX7	116100B
			IX2	X3-X4
			IX7	X0+X4
			LX.0	59-NBITS
			LX.4	59-NBITS
			AX.0	59-NBITS
			AX.4	59-NBITS
			CPOP	0,360B,132B
			XA.7	A.0+A.4
2 36704 20753 21753 42744			LX.7	59-NBITS
			AX.7	59-NBITS
			DX7	X4*X4
			END	
<hr/>				
			Substituted by assembler for IX7 X0+X4	

Substituted
by
assembler
for
IX7 X0+X4

Figure 7 OPDEF Application

The other pseudo instruction mentioned earlier, IFC, could be used to incorporate the user options, for example, truncation or rounding effects. Each effect would require somewhat different code. Both code sequences could be incorporated into the arithmetic instruction OPDEF, but only the one whose option was on would be invoked as part of the OPDEF, upon assembly.

This package of COMPASS pseudo instructions could greatly simplify the n-bit simulation set up. However, a means would have to be devised to get the OPDEFs inserted within each routine, since they are effective only for the routine in which they were defined. If branching were required within an OPDEF to perform n-bit simulation effects, a special element called *L (REF 3: 2-9) and the SET pseudo could be used within each OPDEF to offer each invocation of the OPDEF a unique label. Figure 8 shows an example of its usage within an OPDEF. The SET *L + 2 instruction would put the correct location into the symbol table for N (shown at location 15). When ZR X.P1,N referenced it, the value (a location) currently in the symbol table would be put in the object code. The SA.P1 P,1 instruction (where the label is actually needed) would be positioned a known number of locations down from the

Location	Label	Instruction
15	N	SET *L + 2 ZR X.P1,N (wants to jump to SA.P1 P,1)
16		(other instructions)
17		SA.P1 P,1 [needs the label]

Figure 8 Example Of Labeling Within An OPDEF

SET instruction. Since each OPDEF invocation would occur at a different location, each reference would be unique, since the value of N would change with the location of each invocation.

Problems Encountered With The COMPASS Level Approach

After further analysis of the COMPASS level approach, three problems were uncovered. These problems will be discussed in the following subsections.

Multiples Of Two

Upon further analysis, it was discovered that when the FORTRAN optimizing compiler 1 or 2 was used, fixed point multiplications and divisions by multiples of two were not converted into arithmetic multiply or divide instructions, but were accomplished by shifting the register contents left or right the number of bits required to obtain the proper result. However, the zero level compiler did convert such occurrences into the expected arithmetic instructions. Therefore, a user would either have to compile the algorithm using the level 0 optimizer which would result in less efficient execution or be restricted from using multiples of two in arithmetic equations, since those arithmetic operations could not be n-bit simulated. The user could still use a multiple of two in an equation for optimization level 1, but not as a constant within the equation. It could overcome the optimizers intelligence by assigning that multiple of two (constant) to a variable just prior to the expression as shown

below. This procedure would force the compiler to generate the COMPASS arithmetic instruction.

(desired equation: $IRSLT = I1/8 * I2$)

ITEMP = 8

IRSLT = I1 / ITEMP * I2

The problem with multiples of two could be overcome by putting an added restriction upon the user. However, the problems described in the following subsections are of a more serious nature and could not be resolved.

No Double Precision Arithmetic Accuracy

Another problem that became evident quickly was that no FORTRAN arithmetic expression could be evaluated in double or triple precision, then stored as a single precision value. This is because at the COMPASS level the last arithmetic operation of a compiled FORTRAN equation is indistinguishable from the first. There is no way at the COMPASS level to detect the first operation generated from a FORTRAN arithmetic statement. The COMPASS code generated from individual FORTRAN statements is transparent at the COMPASS level and in many cases the code from several FORTRAN statements are intermixed together so the computer can execute with greater efficiency. Thus all operations would be limited to single precision and that user option specified in the requirements would not be available.

Requirements For Spare Registers

A much more serious problem evolved when it was determined

that extra registers would be needed, in addition to the ones involved in the arithmetic instruction, in order to perform the required n-bit simulation effects. These extra registers would be required for storing the maximum and minimum values of the n-bit fixed and floating point words. To determine overflow conditions these values would have to be compared to the contents of the registers involved in the arithmetic instruction. In order to make the comparisons, the values must be stored in a 60-bit register. However, due to the nature of the COMPASS level approach, nothing can be presumed to be known about any register before that instruction or after. At the time the arithmetic instruction is executed, any or all of the 16 X and B registers may contain values which were set up by the FORTRAN compiler to be used later. Modification of a register's contents could be disastrous.

Since the status of all registers outside the ones involved in the arithmetic instruction is not known, all registers not involved in the arithmetic operation instruction must have their values restored to what they were just prior to the operation. Either a register saving/restoring technique would have to be employed or else some register with known or fixed values would have to be found which could be used as an extra register then restored afterwards.

Several possibilities were investigated and some techniques were devised, but none which were reliable 100% of the time. The different avenues taken are described more fully in Appendix A. Literature searches in CDC manuals

could not produce any amount of detail about the register handling conventions of the FORTRAN compilers. Nor could any register saving/restoring convention be found. Apparently, the compilers do not save all registers before jumping to external subroutines, as many computer systems do (e.g. IBM). Instead it saves only those registers it knows through its intelligence that it will need later.

Without extra registers, no rounding effects could be performed and no overflows could be detected. It logically follows that without overflow detection, no overflow messages could be printed nor would there be a way or means to substitute maximum values for values that overflow. The only action that could be performed would be to limit fixed point values and floating point mantissas to the specified number of bits of precision (through left and right register shifts, to accomplish the truncation).

Conclusions

The more the COMPASS level approach was explored, the more limitations and shortcomings were uncovered. Single assignment statements could not be detected, double and triple precision operations could not be performed as options, and no completely reliable scheme could be devised to generate extra registers required for overflow detection and handling. Therefore, the COMPASS level approach was aborted. The n-bit simulation problem was next addressed at the FORTRAN level as discussed in the following chapter.

III. FORTRAN Analysis And Design

After failing to establish an operational n-bit simulation at the COMPASS level, a FORTRAN level approach to the n-bit simulation problem was endeavored. The FORTRAN level approach attempted to accomplish n-bit simulation by modifying the user's FORTRAN-programed algorithms. This chapter will describe an analysis of the FORTRAN language with respect to the n-bit simulation problem. Also it will describe the features of the FORTRAN algorithm preprocessor and n-bit simulation subroutines that were designed to accomplish the n-bit simulation.

FORTRAN Problem Analysis

A basic assumption for this thesis was that n-bit wordlength effects could essentially be simulated by modifying only the arithmetic expressions within a program. This is because the n-bit simulation tool was devised to enable a user to evaluate the numerical stability and precision of algorithms as a function of wordlength. In Standard ANSI FORTRAN, arithmetic expressions may appear in almost all statement types. N-bit simulation of arithmetic expressions for all possible statement types would complicate the FORTRAN level approach to a considerable extent. Therefore, it was decided that the initial FORTRAN n-bit simulation tool developed would limit n-bit wordlength effects to only those expressions occurring with FORTRAN assignment statements.

This would be the minimum requirement to n-bit simulate an entire algorithm and would still not unduly inhibit the user from programing in his normal manner.

To simulate the effects of n-bit wordlength for an arithmetic statement, each arithmetic operation and its associated operands would, first, have to be isolated from the rest of the statement. Then each operand would have to be modified so that it represents the n-bit wordlength equivalent (if it has not already). Next, after performing the operation, the result must be similarly modified. If the operation would have resulted in an overflow condition for the n-bit computer, an overflow message should be printed to the user and the overflow value should be replaced with the maximum value representable with an n-bit word (as specified in the n-bit simulation requirements of Chapter I).

To correctly evaluate an arithmetic statement, the order of precedence in which the FORTRAN arithmetic operations are normally performed would have to be preserved. Also the variable types (fixed and floating point) would have to be determined then handled appropriately and labels accompanying arithmetic statements would have to be handled so that the execution sequences of the algorithm would not be affected.

Figure 9 shows the Backus Normal Form (BNF) for all possible syntactic representations of an arithmetic assignment statement. Any arithmetic assignment statement may consist of unaries, integer (fixed point) and real (floating

<code><arithmetic statement></code>	<code>::= <identifier> = <expression></code>
<code><identifier></code>	<code>::= <letter> <identifier> <letter> <identifier> <digit> <function></code>
<code><expression></code>	<code>::= <sign> <term> <expression> <sign'> <term></code>
<code><term></code>	<code>::= <factor> <term> <mul-op> <factor></code>
<code><factor></code>	<code>::= <f2> <factor> <exponentiation> <f2></code>
<code><f2></code>	<code>::= <identifier> <literal> (<expression>)</code>
<code><function></code>	<code>::= <identifier> (<arguments>)</code>
<code><arguments></code>	<code>::= <expression> <arguments> , <expression></code>
<code><literal></code>	<code>::= <I2> <real #></code>
<code><real #></code>	<code>::= <I2> . <I2> . <I2> . <I2></code>
<code><I2></code>	<code>::= <digit> <I2> <digit></code>
<code><letter></code>	<code>::= A B C D E F G H I J K L M N O P Q R S T U V W X Y Z</code>
<code><digit></code>	<code>::= 0 1 2 3 4 5 6 7 8 9</code>
<code><sign></code>	<code>::= <sign'> <null></code>
<code><sign'></code>	<code>::= + -</code>
<code><mul-op></code>	<code>::= * /</code>
<code><exponentiation></code>	<code>::= **</code>

Figure 9 Backus Normal Form For Arithmetic Statement

point) variables, integer and real literals, arithmetic operations, parentheses, array items, and function subroutine calls. Some examples of various arithmetic assignment statements are shown in Figure 10. Boolean IF statements may also contain arithmetic assignment statements. The arithmetic statement would be executed if the logical or relational expression in the IF statement predicate is true. The basic form of a Boolean IF statement with an arithmetic statement is:

IF (Logical or relational expression) <arithmetic stmt>

Label field	Col. 7
	A = B*C*(LONG1-5.324/(-24))
1	B = FUNC(4*J,-K,F2-LAST)+NUMB1
	C = -W**5
25	D = ARRAY1(5,I)*5.014327
	E = +(0.5**3**TWO)/DIVDER
4	IF(A.EQ.(10*R))F = 2000+45.+3*REALNUM
	RESULT = FINAL

Figure 10 Examples Of Legal Arithmetic Statements

Evaluation Of An Arithmetic Statement Through Reverse Polish Notation

The task of simulating the effects of an n-bit wordlength machine requires that the result of a computation for an arithmetic statement be modified to reflect its n-bit wordlength equivalent. This necessitates the isolation of each arithmetic operation and the operands the operation is acting upon.

Therefore, each FORTRAN arithmetic statement must be broken down into a series of arithmetic statements, each performing a single arithmetic operation. The final result from the series of single computation statements would be algebraically equal to the result computed from the whole, complex arithmetic statement from which they were derived.

An example of breaking a complex statement into a series of single computation statements is shown in Figure 11. Equation (a) of Figure 11 could be accomplished just as well through the sequence of single operation arithmetic statements shown in (b). However, the intermediate variables used in (b) must agree with the dominant data type (mode) of each computation. That is, when an expression contains both fixed and floating point quantities, the final result would have to be of the floating point data type since floating point is dominant over fixed point in FORTRAN.

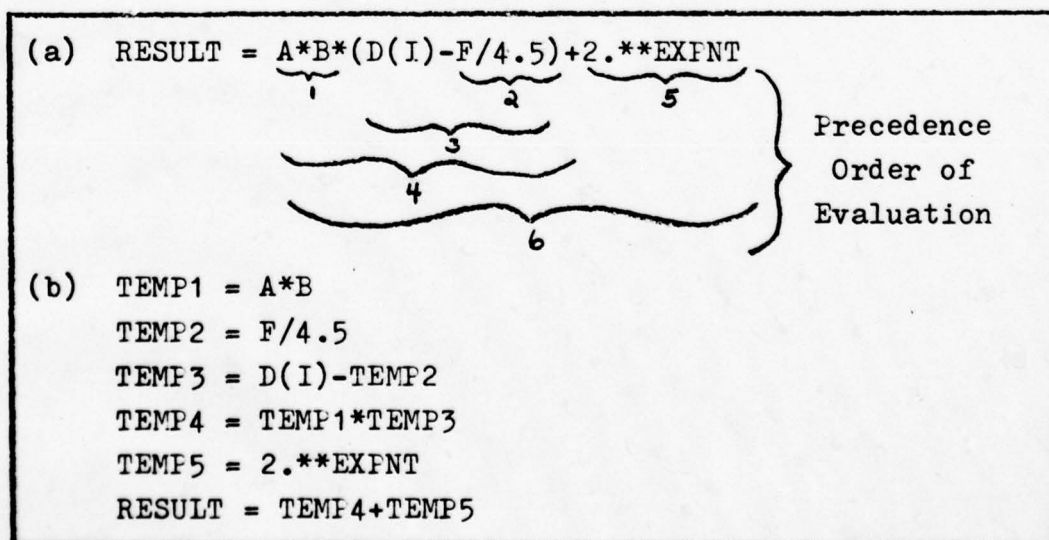


Figure 11 Equivalent Computation Of A Complex Arithmetic Equation

When a FORTRAN compiler analyzes FORTRAN arithmetic statements, it determines the proper sequence in which the computations should be performed and the necessary mode conversions needed for the differing variable data types. Then the compiler generates the assembly language instructions which will do the job; forming intermediate results between computations until the final assignment is made.

An approach similar to one a FORTRAN compiler might use for evaluating an arithmetic statement is the Early Operator Reverse Polish Notation (REF 5: 55). Polish notation is a mathematical notation that provides for the representation of complex expressions in a non-ambiguous manner, without relying on hierarchial delimiters such as parenthesis. The basic difference between standard mathematical notation and Polish notation is the relative position of the operator to its operands.

Figure 12 shows an example of a FORTRAN arithmetic statement and its equivalent Reverse Polish Notation representation. To evaluate the Polish string, it is scanned left to right until an operator is found. Then that operator is used on the preceding two operands-replacing the two

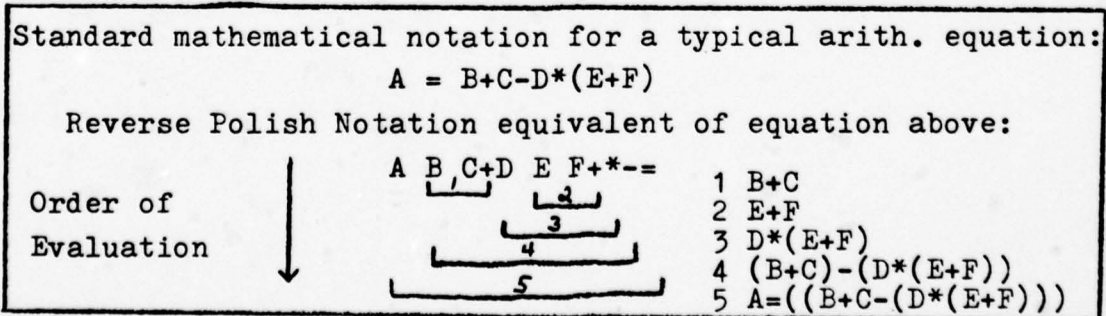


Figure 12 Example Of Reverse Polish Notation

by the result and continuing the scan. Each arithmetic operator has a hierarchical precedence value. Table II shows the precedence relationship among the arithmetic operations of the FORTRAN language. Operations with higher precedence in the hierarchy are performed first in the evaluation of an arithmetic equation. Operations of equal precedence are evaluated left to right. The BNF for FORTRAN arithmetic statements in Figure 9 also reflects this precedence of the operators.

	Arithmetic operation	Precedence rank
Higher precedence	+	1
	-	1
	*	2
	/	2
	- (unary)	3
	+ (unary)	3
	**	4

Table II Operator Precedence Order For Evaluating Arithmetic Equations

Through the use of Polish Notation techniques, the process of generating a sequence of single operator statements from a single complex arithmetic statement is greatly simplified. Once an arithmetic statement is in its equivalent Polish Notation form, the Polish string can be scanned left to right for operators. For each operator a function subroutine call is generated with the operands involved passed along as arguments. The subroutine called would perform the computation while simulating the effects of n-bit wordlength.

Additional FORTRAN Language Considerations/Problems

In the analysis of the FORTRAN language with respect to the n-bit simulation problem, several obstacles and potential problems had to be dealt with to handle all cases of arithmetic statements. Besides the variables and arithmetic operations, an arithmetic statement may involve function subprogram references, arrays, labels, and IF statements. An additional consideration is to provide the user a means of tracing down overflow occurrences. These cases will be discussed in the following subsections.

Function References. Function subprogram references with a FORTRAN arithmetic statement present a potential problem since the functions they reference may not be coded in FORTRAN. To be n-bit simulated completely, all code executed must be written in the FORTRAN language. System functions such as SIN, COS, and ATAN (external or FORTRAN library functions (REF 4, I-8)) are not coded in the FORTRAN source language and therefore can not be altered by the FORTRAN level approach to simulate n-bit wordlength effects. A conversion of all intrinsic and external functions into their FORTRAN language equivalents would remedy this problem. However, without a FORTRAN version, the most effective action that can be done to approximate n-bit wordlength affects on an intrinsic function would be to n-bit simulate each argument of the function and then n-bit simulate the function value returned.

Arrays. N-bit simulation for each argument (subscript) of a FORTRAN array item would not really serve any purpose. FORTRAN array arguments must always be fixed point values. They are usually relatively small fixed point values and should be well within the range of an n-bit wordlength computer's fixed point words, in presumably every case. Consequently, it was considered unnecessary overhead to n-bit simulate expressions and single arguments used for array subscripts. Treating the entire array item the same as a single variable was found sufficient.

Since arrays and function subprogram references may have the same syntactical representation in an arithmetic statement but would be n-bit simulated differently, they must be differentiated from each other. To distinguish arrays from functions, each program unit is examined to determine the variables declared as arrays. Arrays may be declared in either DIMENSION, INTEGER, REAL, or COMMON declaration statements. So for each variable in question, all array names declared within that program unit are compared to that variable name. If the variable in question is not an array, then the variable defaults to being treated as a function subprogram reference, and n-bit simulated accordingly.

Labels. Another potential problem area for FORTRAN is where a label is attached to an arithmetic statement which is the last statement to be executed in the DO loop. A label in FORTRAN can be used for two different purposes.

These usages of a label will be described before defining the problem in detail.

A statement label uniquely identifies a statement so it can be referenced by another statement. Labels occur in columns 1-5 of a FORTRAN statement. The actual statement appears in columns 7-72.

Statement number	Label	Statement
1		C = X+2.5
* 2		GO TO 38
3	14	A = B
4	38	C = B*D/5
* Execution of statements 1 and 2 would be followed by a branch statement 4 (through the use of label 38)		

Figure 13 Normal Use Of A Statement Label

Generally, a label marks a statement so that the statement can be branched to for execution, as shown in Figure 13. However, there is an exception in the case of DO loops. A DO loop statement, such as the one shown in Figure 14, is used to establish a sequence of statements to be executed repeatedly for a specified number of times. In the case of Figure 14, a label is associated with the last statement of a DO loop sequence. This label is called a DO loop termination statement label. Upon execution of its associated statement, program execution returns to the first statement of the DO loop sequence to execute the DO loop sequence again. This looping would continue until something causes the program to branch out of the loop. The main difference

between the DO loop termination label and a normal label is that a termination label is used to indicate the last statement of a sequence to be executed and a normal label is used to indicate the first of a sequence.

Statement number	Label	Statement
1		DO 70 L = 1,20
2		A(L) = A(L)+5
* 3	70	C(L) = D(L)+A(L)
* Label 70 indicates last statement executed in the DO loop. After execution of statement 3, L would be incremented and the loop executed once again, until L exceeds the value 20.		

Figure 14 Example Of DO Loop Termination Label

The problem with respect to this effort arises when a complex arithmetic statement is broken up into a sequence of single operation arithmetic statements; especially if the FORTRAN program is analyzed and translated on a statement by statement basis (i.e. each statement is analyzed without knowledge of the context of its use with regard to other statements of the program). Each label type must be handled differently. Figure 15 shows an example of what would happen if a termination label were treated the same as a normal label in the translation of a complex arithmetic statement to its single operation series equivalent. In the example, it is associated with an arithmetic statement. Normal label handling would distort the real meaning of the DO loop. In the translated code of Figure 15, part of the computation is

now positioned outside the DO loop sequence. Thus, the results would not be the same.

Before translation	After translation
DO 5 I = 1,7	DO 5 I = 1,7
5 SUM = A(I)+SUM*2	5 T1 = SUM*2
	SUM = A(I)+T1 ← outside of loop now

Figure 15 Incorrect Translation Of A DO Loop

To assure a consistent meaning of labels for n-bit simulation modifications, the user will always be required to assign DO loop termination labels to CONTINUE statements. A CONTINUE statement can always be used as the last statement in a DO loop sequence without changing the loop's meaning. A CONTINUE statement will never be modified for n-bit simulation since only arithmetic statements are changed. Figure 16 shows how a DO loop termination label could be moved to a CONTINUE statement without changing the performance of the loop. Figure 16 also shows the single operation translation of that loop.

Before translation	After translation
DO 5 I = 1,7	DO 5 I = 1,7
SUM = A(I)+SUM*2	T1 = SUM*2
5 CONTINUE	SUM = A(I)+T1
	5 CONTINUE

Figure 16 Use Of CONTINUE Statement In A DO Loop

IF Statements. Another slight complication of the FORTRAN approach is translating an arithmetic statement contained as part of an IF statement. Figure 17 (a) shows such an IF statement. A procedure had to be developed so that arithmetic statements contained as part of IF statements could be translated into a series of single operation statements without distorting the original program's meaning. Figure 17 (b) restates (a) in such a manner that the arithmetic statement is now separated from the IF statement and could be translated into a sequence of single operation statements, shown in (c), without affecting the original meaning. Two new labels and two GO TO statements were added to the original sequence in (a) to preserve the correct order of execution. The execution paths for (a), (b), and (c) are identical as shown in (d).

Overflow Tracing. Another FORTRAN consideration was assisting the user in tracing down the instances where an overflow of the n-bit wordlength occurred during the n-bit simulation of an algorithm. It would be helpful to the user to indicate the name of the routine and the line number within the routine in which the overflow condition was detected. This would be in terms of the original FORTRAN coded algorithm and not the n-bit simulation modified code. In addition, specifying the operation executed at the time of the overflow might also be helpful.

To assemble the necessary data for these overflow messages, the beginning of each new routine must be detected.

(a) A = FIRST
 IF(A.EQ.C) RESULT = A/5*B
 A = NEXT

(b) A = FIRST
 IF(A.EQ.C) GO TO 1500
 GO TO 1501
1500 RESULT = A/5*B
1501 A = NEXT

(c) A = FIRST
 IF(A.EQ.C) GO TO 1500
 GO TO 1501
1500 TEMP1 = A/5
 RESULT = TEMP1*B
1501 A = NEXT

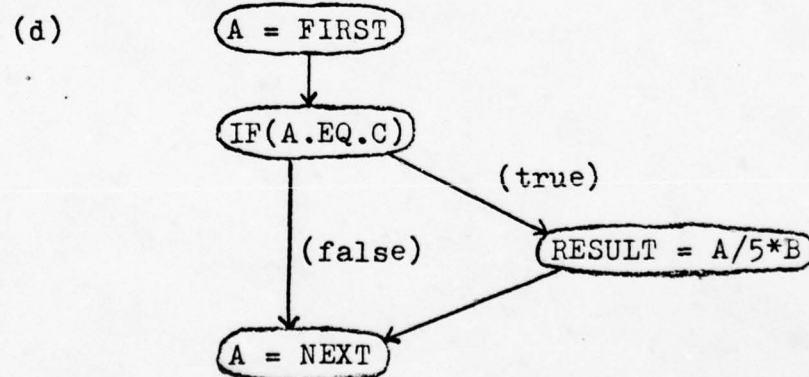


Figure 17 Arithmetic IF Statement Handling

New program units (routines) are initiated by one of the following FORTRAN keywords: PROGRAM, SUBROUTINE, FUNCTION, REAL FUNCTION, or INTEGER FUNCTION. All other routine initiation keywords, such as LOGICAL FUNCTION or DOUBLE PRECISION are not allowed in n-bit simulation programs, by one of the thesis assumptions.

Once the beginning of a new routine is detected the routine name is extracted, the line counter for that routine is initialized to one and incremented with each succeeding statement. When an arithmetic statement is translated into a series of subroutine calls, the routine name and line number containing the arithmetic statement are included as arguments. Therefore, when an n-bit wordlength overflow is detected, this information along with the operation being performed at the time would be printed. Then, the user could easily trace down the statement which was the source of the overflow condition.

The analysis of the FORTRAN language with respect to n-bit simulation now has been described. The following section describes the system developed which incorporated the results of the FORTRAN analysis and enables a FORTRAN algorithm to simulate n-bit wordlength effects.

N-bit Simulation Tool Design

The design of the n-bit simulation tool was divided into two parts. Part one preprocesses the FORTRAN algorithms before they are executed; translating each FORTRAN arithmetic

statement into a series of subroutine calls, each of which performs one of the statements's computations. Part two of the design develops the subroutines which would perform the n-bit wordlength effects.

Preprocessor

The purpose of the preprocessor is to perform a lexical scan and analysis of each statement in the FORTRAN algorithm, then convert the arithmetic statements into a series of n-bit simulation function subroutine calls.

The preprocessor essentially performs two major functions:

- 1) It performs a syntactic analysis for arithmetic statements and converts them into their Reverse Polish Notation equivalent representation. (Further detail concerning the syntax analysis is contained in Appendix D.).
- 2) The preprocessor disassembles the Reverse Polish Notation representation into function subroutine calls, tailored according to the variable data types, and arithmetic operations involved. Each function reference replaces an arithmetic operation from the original statement. If temporary variables are required to store intermediate results, a variable matching the dominant data type involved is created. In addition, for each function reference, the preprocessor supplies the routine name and line number of the FORTRAN statement from which the statement was derived.

The syntax description for an arithmetic statement was shown by the BNF of Figure 9. Each statement fitting this syntax description would be transformed into one or more function subroutine references. Each function reference would replace a single arithmetic operation. Any statement which did not fit the syntax description would not be affected. Figure 18 contains an example of a preprocessed program.

In Figure 18, the temporary variables were created to hold the intermediate computed values. Floating point variables are stored in variables preceded by RTTTT and fixed point variables by ITTTT. The last letter is incremented from A to B to C and so on for each succeeding intermediate variable within each equation. These names will be restricted from the user, but are named so that they are not likely to conflict with any variable name a user might create.

The function reference names indicate the type of operation that is being performed from the original arithmetic statement. For example, statement 8 of the preprocessed program performs the multiplication of 15 and B from the original program statement 8. The last three letters of the function name indicate the type of operation (e.g. ADD for addition, and MPY for multiply). The first two letters indicate the data type of the operands being passed as arguments. II indicates both operands are fixed point variable data types. RR indicates both operands are floating point. R1 and R2 indicate one of the operands being passed is fixed point and the other floating point. The number 1

Statement #	Original FORTRAN Algorithm
1	PROGRAM SAMPLE (INPUT, OUTPUT)
2	DIMENSION A(10), K(30), KEY(7)
3	CALL SETNBIT (24,1,1,1,17,6,0,KEY)
4	B = 25.3
5	READ *, L1
6	L1 = L1
7	CALL SUB1 (A,25,B)
8	NEXT = A(1)+15*B
9	K(L1) = NEXT/50*(A(2)+B)
10	STOP "END OF SAMPLE"
11	END
Original Statement #	Preprocessed Algorithm Version
1	PROGRAM SAMPLE (INPUT,OUTPUT)
2	DIMENSION A(10), K(30), KEY(7)
3	CALL SETNBIT (24,1,1,1,17,6,0,KEY)
4	B = RASGN(25.3, 7HSAMPLE ,6,1,KEY)
5	READ *, L1
6	L1 = IASGN(L1, 7HSAMPLE ,6,1,KEY)
7	CALL SUB1(A,25,B)
8	RTTTTAA = R1MPY(15,B, 7HSAMPLE ,8,0,KEY)
	NEXT = RRADD(A(1),RTTTTAA,7HSAMPLE ,8,1,KEY)
9	ITTTTAA = IIDVD(NEXT,50,7HSAMPLE ,9,0,KEY)
	RTTTTAA = RRADD(A(2),B,7HSAMPLE ,9,0,KEY)
	K(L1) = R1MPY(ITTTTAA,RTTTTAA,7HSAMPLE ,9,1,KEY)
10	STOP "END OF SAMPLE"
11	END

Figure 18 Example Of A N-bit Simulated Algorithm

in R1 indicates the fixed point variable is the first argument of the function, 2 indicates it is the second. Table III contains a list of all function subroutines. It also indicates the arithmetic operation performed and the operand data types involved.

Name of Function	Operation	Operand 1 Data Type	Operand 2 Data Type
IIADD	+	I	I
R1ADD	+	I	R
R2ADD	+	R	I
RRADD	+	R	R
IIMNS	-	I	I
R1MNS	-	I	R
R2MNS	-	R	I
RRMNS	-	R	R
IIMPY	*	I	I
R1MPY	*	I	R
R2MPY	*	R	I
RRMPY	*	R	R
IIDVD	/	I	I
R1DVD	/	I	R
R2DVD	/	R	I
RRDVD	/	R	R
IIEXP	**	I	I
R1EXP	**	I	R
R2EXP	**	R	I
RREXP	**	R	R
IASGN	=	I	-
RASGN	=	R	-

I - Integer (Fixed Point)
R - Real (Floating Point)

Table III N-bit Simulation Function Subroutine Names

There are two forms of n-bit simulation functions: single operation replacements and simple assignment statement replacements (e.g. IASGN and RASGN). The general format of these two function types are:

```
FUNCTION (OPERAND1,OPERAND2,ROUTINE,STMT#,FINALFLG,KEY)  
FUNCTION (OPERAND,ROUTINE,STMT#,FINALFLG,KEY)
```

The OPERANDS are variables which will be n-bit simulated by the function. The ROUTINE is the name of the routine containing the statement and STMT# is the line number of the routine in which the statement occurs.

FINALFLG is used to distinguish intermediate value assignments from final assignments of an arithmetic assignment statement. Zero indicates the result is an intermediate assignment, one indicates a final assignment. Intermediate assignment results will maintain the arithmetic precision specified by the user; either single, double, or triple precision. Final assignments indicate the final computation of the arithmetic statement. The resulting value from a final computation always has single precision accuracy.

The last argument of the function is KEY. KEY is a seven item array containing values computed by the SETNBIT subroutine and is used by the n-bit simulation subroutines to simulate n-bit wordlength effects.

Each arithmetic statement in its original form appears in a COMMENT statement (a non-executable statement type). This COMMENT is printed just prior to the series of function

calls derived from it (Figure 19). This procedure helps improve readability of the preprocessed program code for the user.

Stmt. #	Original Algorithm Statements:
1	A = B+FUNC(A,5*K,K1) where FUNC is a FUNCTION
2 51	IF(A.EQ.C) LAST = FINAL/R*S*25.97
Preprocessed Algorithm Statements:	
C	A = B+FUNC(A,5*K,K1) ←
	RTTTTAA = RASGN(A,7HPROGRAM1,1,1,KEY)
	ITTTTAA = IIMPY(5,K,7HPROGRAM1,1,1,KEY)
1	ITTTTAB = IASGN(K1,7HPROGRAM1,1,1,KEY)
	A = RRADD(B,FUNC(RTTTAA,ITTTAA,ITTTAB),
	* 7HPROGRAM1,1,1,KEY)
C	IF(A.EQ.C) LAST = FINAL/R*S*25.97 ←
51	CONTINUE
	IF(A.EQ.C) GO TO 7901
	GO TO 7902
2	
7901	RTTTTAA = RRDVD(FINAL,R,7HPROGRAM1,2,0,KEY)
	RTTTTAB = RRPY(RTTTAA,S,7HPROGRAM1,2,0,KEY)
	LAST = RRPY(RTTTAB,25.97,7HPROGRAM1,2,1,KEY)
7902	CONTINUE

Figure 19 Preprocessor Output For Two Complex Arithmetic Statements

A brief description of the preprocessor's operation has now been given. A detailed description of the preprocessor is given in Appendix C. Chapter 5 contains a description of the preprocessor design and summarizes the functions of the modules within the preprocessor. The following section describes the n-bit simulation subroutines which are called by the n-bit simulated algorithm built by the preprocessor.

N-bit Simulation Subroutines

The n-bit simulation function subroutines are designed to perform the n-bit wordlength effects upon the operands involved in a single arithmetic operation. A separate function subroutine was developed for each operation and data type combination as shown in Table III. Each routine simulates n-bit wordlength effects according to the options of the user. The user states the options desired by a call to a SET-NBIT routine to be subsequently discussed in detail. The various user options are stated in Figure 20. Options 2, 3, and 4 apply only to floating point computations. The following subsections will show how each option is carried out.

User specifications:

- 1) Total number of bits per word
- 2) Number of bits in exponent and number of bits for mantissa, in addition to the position of the implied decimal point with respect to the mantissa. The sum of the mantissa and exponent bits must be one less than the total number of bits specified in 1). The extra bit is for the sign.
- 3) Rounding or truncation effects for computations.
- 4) Single, double, or triple precision arithmetic accuracy.
- 5) Overflow messages print or suppression.

Figure 20 User Specifications (Options)

Option One. The first option indicates the total number of bits in the wordlength to be simulated. The SENBIT routine computes the maximum value that a fixed point word of n-bit wordlength could contain. The absolute value of

each fixed point result is compared to this maximum. If the maximum is exceeded, overflow has occurred. The value that overflowed is then replaced by the maximum positive value possible if the positive maximum was exceeded, or the maximum negative value possible if the negative maximum was exceeded.

Option Two. The second user option defines the data format to be used for floating point values. The different specifications include the number of mantissa and exponent bits along with implied decimal point with respect to the mantissa (either to its left or to its right). They are used by SETNBIT to compute the maximum and minimum positive floating point values representable for the n-bit word. If the absolute value of a floating point result exceeds these limits, an overflow or underflow condition has occurred. Underflow means the result represented was a smaller value than could be stored in an n-bit word with the given number of exponent bits. Upon detection of underflow, the value that underflowed is replaced with zero. Overflow detection is treated similar to fixed point overflow, except the maximum floating point values are used.

Also, for floating point values, the proper number of significant mantissa bits must be maintained. This is achieved by a right-shift (sign extended) followed by a circular left-shift of the floating point variable. For example, if 18 bits were specified for the mantissa, the rightmost 30 significant bits would be truncated, as shown in Figure 21.

(In Octal)	
Original value	17465760371015777610 <div style="text-align: center;"> </div> exponent 48 bits of sig- nificance
Right-shifted 30	000000000001746576037 <div style="text-align: center;"> </div> exponent 18 bit mantissa
Circular left-shifted 30	17465760370000000000 <div style="text-align: center;"> </div> 18 bit mantissa

Figure 21 N-bit Simulation Of Mantissa

Option Three. The third user option is rounding or truncation effects. If truncation is specified no special action has to be taken, because the CDC 6600/CYBER 74 computers truncate by default. To round, however, a value equal to one half of the least significant bit (LSB) is added to the absolute value of the word. This is accomplished by right-shifting the number of mantissa bits specified; so that the least significant bit is in the second bit position from the right of the CDC 60-bit word (the 59th bit). Then the value one is added (using fixed point addition) to the quantity and the resulting quantity is right-shifted one bit. If the LSB-1 is a one, then the result would be rounded up; if it is not, it would have no effect. Figure 22 shows a case of a 15-bit mantissa in which the result is rounded up. After the one bit right-shift, the floating point value is left-shifted back to its original position.

A special condition may result from this technique of rounding. This condition occurs when the addition of one to

001111100110101111110101110111000110101100011010001000000000

[illegible][illegible][illegible][illegible]

00111110011010111111010111100000000000000000000000000000000000

the right-shifted floating point value results in an increment of the exponent. The n-bit mantissa contains its maximum representation in these cases. Figure 23 shows a rounded value for a floating point value with a 15 bit mantissa. The rounded result, shown in (b) has no significant mantissa bits and would be treated from there on by the computer the same as a floating point zero. However, in a n-bit

(a) 17257777750031152437

15 bit mantissa

(b) 17260000000000000000

(c) 17264000000000000000

computer the round up of the value in (a) would result in an increment of the exponent and the mantissa would have a one in its most significant position, as shown in (c).

Option Four. The fourth user option is to specify single, double, or triple precision arithmetic accuracy. This option is used to simulate the effects of a whole complex equation being computed with multi-word precision, then storing the result as a single precision value. The preprocessor passes a flag which indicates if the n-bit subroutine computation will result in an intermediate value or a final assignment. The SETNBIT routine computes the bit shifting needed to be done for the intermediate values.

A floating point 15-bit word with nine mantissa bits and the triple precision option would be handled as shown in Figure 24. In a triple word accumulator of a 15-bit computer, there would be 39 significant mantissa bits (i.e. 9 bits plus two additional 15-bit words used to extend the mantissa significance). The final value assigned would be handled just as any 15-bit single precision floating number, similar to Figure 21.

Original value	17465760371015777615
Right-shift 9 bits	00017465760371015777
Left-shift 9 bits	17465760371015777000
	39 bit mantissa

Figure 24 Multiple Precision Accuracy

Option Five. The fifth option can be used to print or suppress overflow messages when the n-bit wordlength is

exceeded. A test is made in the overflow handling routine. If the message flag is off (set to zero), no message is printed. If the message flag is on (set to one), a message is printed indicating the line number and routine in which the overflow resulted from and the operation being performed at the time the condition was detected.

N-bit Simulation Sequence

The general sequence of events involved in simulating n-bit wordlength effects for each n-bit simulation routine is listed below:

- 1) Perform the arithmetic operation using the two input operands and storing the result. Input operands are assumed to have n-bit wordlength significance already.
- 2) If the rounding option (which applies only to floating point computations) is used, rounding effects are performed on the result.
- 3) Check the result for overflow. If overflow conditions exist, branch to overflow handler to substitute a value for the overflowed result and print an overflow message (if the message print option is on).
- 4) For floating point quantities, truncate the mantissa of the result to the number of bits specified for the n-bit floating point mantissa. If this is an intermediate result (indicated by the Final Flag input argument), the result's mantissa is truncated

to the precision specified (single, double, or triple).

5) Return to the calling routine.

SETNBIT. The subroutine SETNBIT computes boundary values based on the user options selected (Figure 20). The SETNBIT subroutine is required to be the first executable statement in each program unit. The user specifies the options desired as arguments to the SETNBIT subroutine. Different routines may execute with different options. Each routine's options are translated into seven key values which are localized to each routine in an array called KEY. The array KEY is used by n-bit simulation function subroutines to carry out the n-bit wordlength effects specified by the user. A description of each of the seven values of the KEY array and how each is computed is as follows.

KEY(1) represents the maximum fixed point value representable with n-bits. The maximum value is computed by setting the sign bit (left-most bit in a CDC 60-bit word), then circular left-shifting n, the number of bits in the wordlength, and subtracting one, as shown in Figure 25 (where n equals 15).

(In Octal)	
1. Set sign bit:	40000000000000000000
2. Circular left-shift 15 bits:	0000000000000000040000
3. And subtract one:	000000000000000037777

Figure 25 Computation Of Maximum Fixed Point Values

KEY(2) represents the maximum floating point value possible for an n-bit floating point word with the specified floating point word characteristics (exponent/mantissa bit lengths and the implied decimal point position). This maximum value is computed by determining the maximum range for the n-bit word exponent. This exponent maximum is computed exactly in the same manner as KEY(1). The exponent range would be equal to one-half the maximum exponent representable. For a five-bit exponent and a ten-bit mantissa, the maximum representable would be (in binary) 11111 or 2^{+31} . The minimum exponent would be 00000 or 2^{-31} (assuming a bias of 10000). If the implied decimal point is to the left of the mantissa (making the whole mantissa a fractional value), the maximum exponent is decreased by the number of mantissa bits (ten for this example) in the n-bit word. If the implied decimal point is to the right, the maximum exponent is not affected.

KEY(3) represents the least significant floating point value representable with the specified mantissa, exponent, and implied decimal point position. Figure 26 shows this value for a 16-bit word with a five-bit exponent and ten-bit mantissa with the implied decimal point to the right. To compute this value, the fixed point representation of the mantissa, shown in Figure 26 is multiplied by two to the

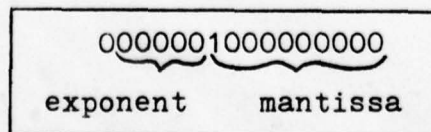


Figure 26 Example Of Least Significant Floating Point Value

maximum negative 31 exponent (-31 for a five-bit exponent). If the mantissa is a fractional representation for this example, the number of mantissa bits (ten) is subtracted from the maximum negative exponent (-31). So the least significant value in this case would be the fixed point mantissa value times $2^{(-31 - 10)}$ or 2^{-41} . This last value would be equivalent to .5 times 2^{-31} . Any value computed in the n-bit simulated algorithm with less significance would constitute an underflow for the n-bit word.

KEY(4) and KEY(5) simply contained the options indicated for rounding/truncation and message printing/suppression, respectively.

KEY(6) indicates the number of bits a mantissa must be right-shifted then left-shifted to retain the proper amount of mantissa significance, which would be 48 minus the number specified. Forty-eight is the number of bits in the mantissa of CDC's 60-bit word.

KEY(7) indicates the number of bits a mantissa must be shifted to retain single, double, or triple precision accuracy. This is computed by adding the number of mantissa bits to another full-wordlength (n-bits), for double precision, or to two wordlengths for triple precision. Then that result is subtracted from 48. For a 16-bit word with a nine-bit mantissa, the resulting quantity for single precision would be 39 or $(48 - 9)$. For double precision, it would be 23 or $(48 - (9 + 16))$. For triple precision, it would be 7 or $(48 - (9 + 16 + 16))$.

Besides precomputing the values for the KEY array, SET-NBIT performs validation checks on the options to verify they are consistent and can be simulated using the CDC 60-bit wordlength. Failure of a validation check results in program termination.

The Overall N-bit Simulation Tool

To use the n-bit simulation tool, a user must first comply with the restrictions which are described in Appendix B. Two major restrictions are to debug the FORTRAN algorithm from all syntax errors prior to execution, and to assign all external values of the program (such as, through a read) to themselves. For the variable L5, the statement $L5 = L5$ would reduce the value in L5 to its n-bit wordlength significance. One additional restriction is that the first executable statement of each program routine must be a call to the SETNBIT subroutine containing the user's options, and the array variable KEY must be dimensioned for seven items. The SETNBIT subroutine call allows the user to specify different n-bit wordlength characteristics for different subprograms within the same set of programs.

The preprocessor will read the FORTRAN algorithm and transform each arithmetic statement into a series of function subroutine calls. The preprocessed program would then be compiled and loaded with the n-bit simulation function subroutines and executed. During program execution, the n-bit function subroutines would perform the arithmetic

operations and assignments while simulating the effects of n-bit wordlength and other user options.

Appendix B contains the User's Guide to the n-bit simulation tool and Appendix C contains a Detailed Description of the preprocessor and the n-bit simulation subroutines. The next chapter analyzes the design of the preprocessor in further detail and describes the modules contained within it.

IV. Structured Design Of The Preprocessor

This chapter describes the final design of the n-bit simulation preprocessor which resulted from the application of structured design techniques. It also discusses how applicable structured design techniques were found to be for a small system design (i.e. the preprocessor) of approximately 1000 source statements.

Structured design is "the art of designing the components of a system and the interrelationship between those components in the best possible way" (REF 9: 7). Basically, it is a form of top-down design. One structured design method is transform analysis. Transform analysis is a strategy that derives initial structure designs which are usually quite good and require only modest restructuring to arrive at the final design. The quality of a design can be measured in terms of its reliability (the number of "bugs" encountered in the program), maintainability (the amount of effort required to fix "bugs" in the program), and modifiability (the cost of changing or extending the program).

Data Flow Description

Using the transform analysis approach, a data flow diagram of the preprocessor was developed, as shown in Figure 27. The preprocessor inputs a FORTRAN statement, then scans the statement for a sequence of characters which could represent an object of an arithmetic statement. An object refers to the variable to which the result of arithmetic

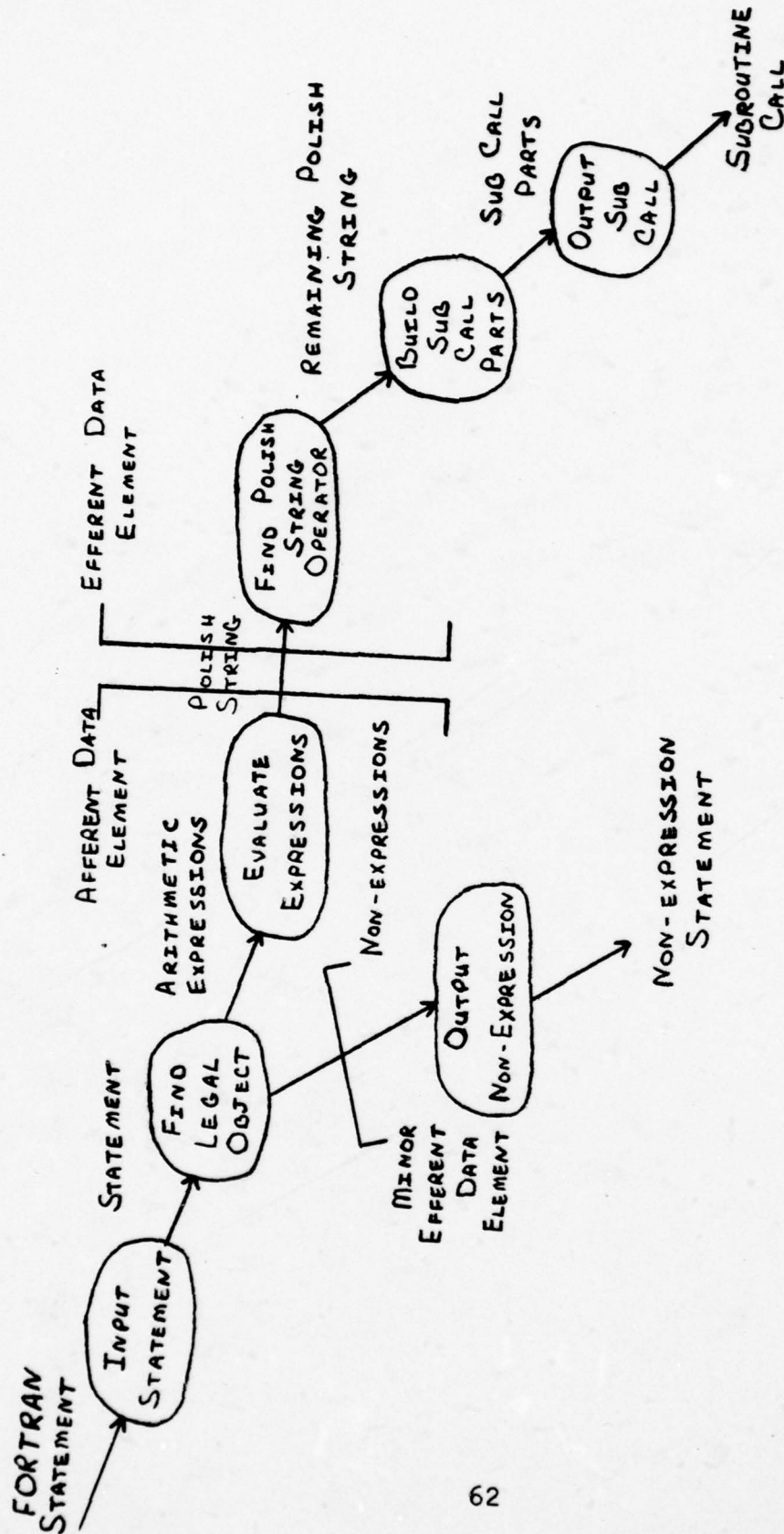


Figure 27 Preprocessor Data Flow Diagram

equation is assigned; where an arithmetic statement is of the general form "object = arithmetic expression". Any statement not having a legal object is regarded as a non-expression and is output in its original form to the n-bit simulated program file.

Upon detection of a legal object, the scan continues to the right side of the '=' in the arithmetic statement to evaluate the expression and convert the statement into its equivalent Polish Notation representation.

The right half of the data flow diagram describes the process in which the Polish Notation string is scanned for operators. Each operator found will result in the building of a function call. To accomplish this, the operands which are being operated upon are found in the Polish string, then merged with the function routine name and other data required to build the function call (such as routine name, line number, and the reference to the KEY array).

The afferent and efferent data elements are defined in the data flow design also. "Afferent data elements are those high-level elements of the data which are furthest removed from the physical input, yet still constitute inputs to the system. Efferent data elements are those furthest removed from the physical outputs which may still be regarded as outgoing." (REF 9: 262). In the case of the preprocessor, the Polish Notation string representation of each arithmetic statement constitutes the afferent as well

as the efferent data elements. All processes to the left of the afferent data element are dedicated towards building the Polish string from the input FORTRAN statement. All those to the right are intent upon converting the Polish string into a series of subroutine calls.

In developing a program structure from the data flow diagram, transform analysis states that major afferent and efferent data elements should be handled by separate modules and that these modules should be immediately subordinate to the main module of the preprocessor. The basic module structure which should be developed from this data flow diagram is shown in Figure 28. The fully factored structure which resulted for the final design is shown in Figure 29. The module definitions for the resulting preprocessor design and the module interfaces are contained in Appendix C. The following section explains the differences between the fully factored structure of Figure 29 and the structure that would be expected to follow from Figure 28.

Departures

Ideally, with the transform analysis strategy, every module of the afferent branch should have only afferent or transform subordinates, where afferent characteristics imply that data is floating inward (actually upward) with respect to the system. Transform modules transform the input data closer to its most abstract input representation

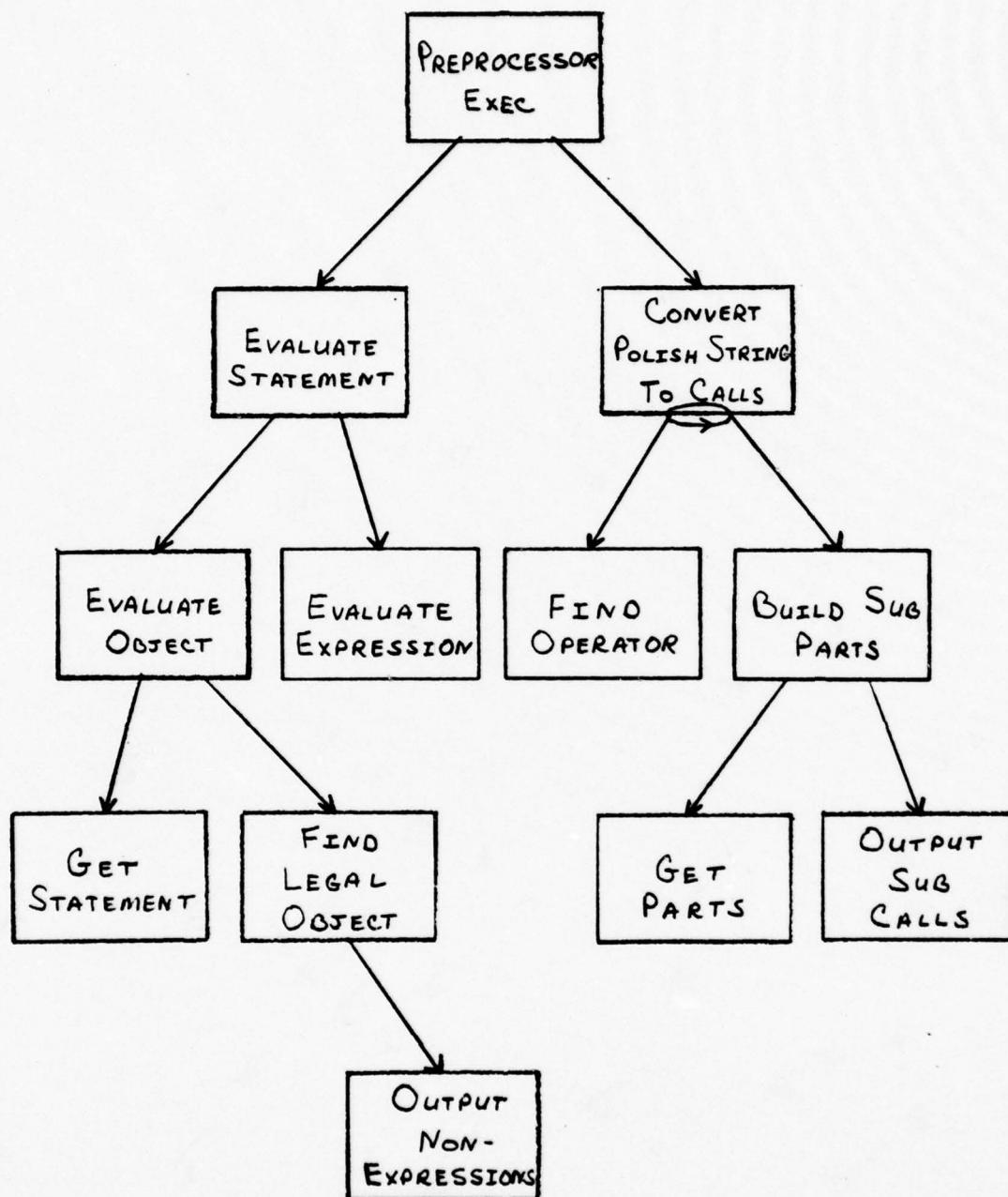
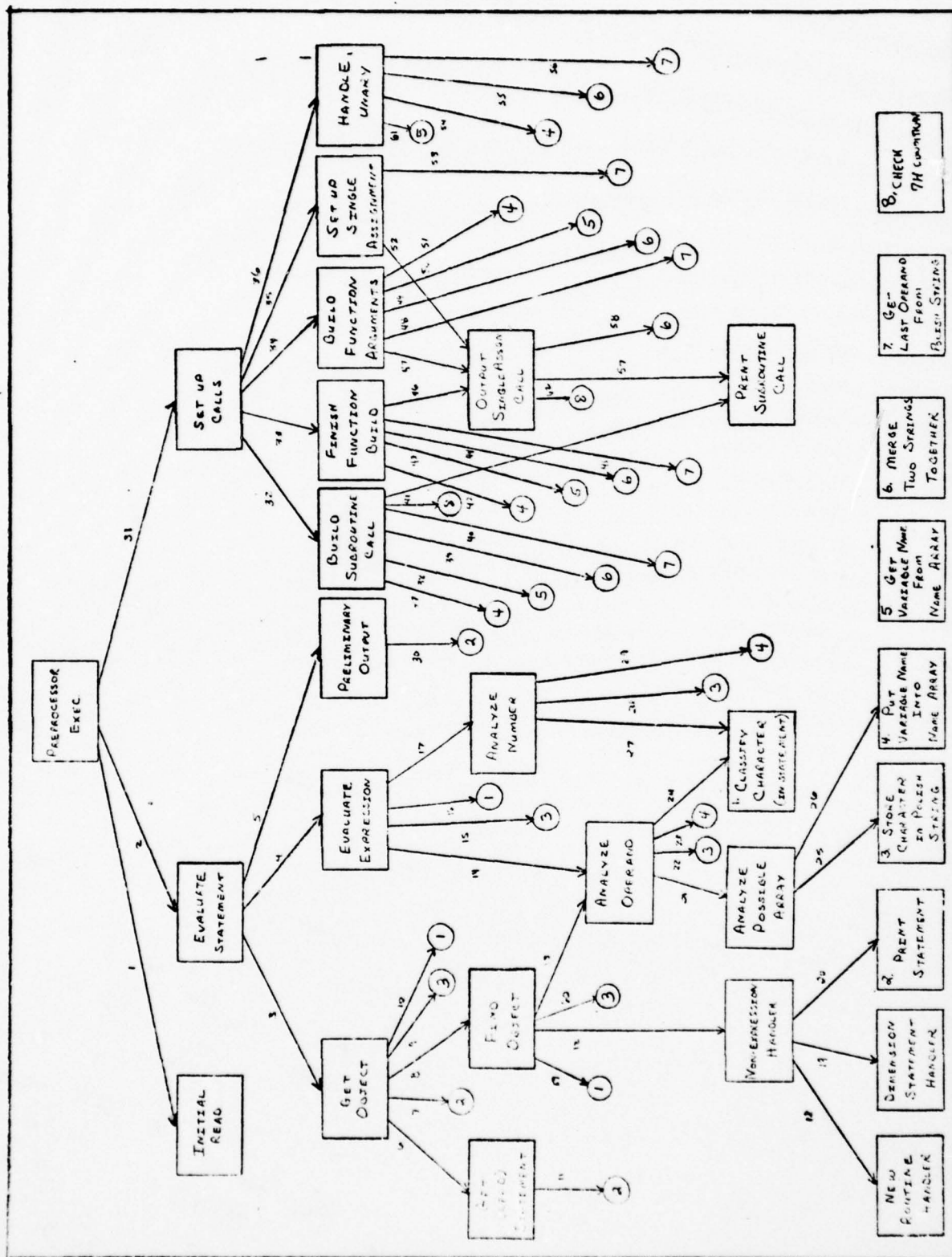


Figure 28 Initial Program Structure For The Preprocessor
Data Flow Diagram



for the system. Efferent (data flowing outward) modules are expected to have only efferent and transform subordinates, where in this case the transform modules transform data closer to its output form. Anything contrary to a consistent data flow (either inward or outward) is termed a departure. The departures of the resulting structure of Figure 29 are discussed in the following paragraphs.

Two of the departures are for two minor efferent modules which appear inside the afferent branch of the pre-processor structure. These are the PRINT-STATEMENT and PRELIMINARY-OUTPUT modules. The PRINT-STATEMENT module is called when any module within the afferent branch detects a non-arithmetic statement. Only arithmetic statements are converted into Polish Notation and subsequently transformed into subroutine calls. All other statements are passed along to the n-bit simulated program file unchanged.

The PRELIMINARY-OUTPUT module is performed once an arithmetic statement has been converted into its Polish Notation equivalent representation. This routine sets up prints to the output file for a comment containing the original arithmetic statement and the preliminary output that is required for labels and IF statements. Perhaps this module should have been included under the efferent branch module structure, but it requires the original input statement. No module in the efferent branch requires the original statement, while almost every afferent branch module does. In order to prevent the original input of the

system from being passed unnecessarily into the efferent branch side, the PRELIMINARY-OUTPUT module was included as the last step of the afferent branch.

The INITIAL-READ module which is directly subordinate to the executive module represents another departure. This module is used to read the first program card image. All succeeding reads are performed by the GET-STATEMENT module, which operates on a 'look-ahead' basis, in order to detect statements that are continued on other cards. The INITIAL-READ module does the initial read so the GET-STATEMENT module does not have to repeatedly test for the initial read on all cards in the program.

The other departures of the final structure in Figure 29 appear in the efferent branch. The module called FINDING-THE-OPERATION in Figure 28 was such a short process that it was merged into the SET-UP-CALLS module. Also, the GET-PARTS module was merged with its subordinate (BUILD-SUB-CALLS) because they were so strongly data coupled. Data coupling measures the amount of data passed as arguments between modules as subroutine parameters. In this case, the GET-PARTS module was very simple and the data coupling between it and BUILD-SUB-CALLS was high enough that a separate module was not justified.

The Effects Of Structured Design

One of the goals of this thesis was to determine how applicable structured design concepts were to small programs.

The importance of structured design in large programs has received considerable publicity, but the programs of less than 5,000 source statements are more common place in the programing world. The n-bit simulation preprocessor is an example of a small program design.

Part of the evaluation of the resulting structured design is based on its comparison to a design developed for the preprocessor before the structured design approach was begun. The previous, or first, program design did not take a formalized approach to structuring the design. It was based upon a flowchart of the preprocessor concept. Many of the modules resulted from blocking off portions of the flowchart. The first design's program modules were fairly intricate and hard to understand. The structured design approach emphasizes that the structure of the program should resemble the structure of the problem, which the first design certainly did not.

The first design consisted of 16 modules that were interdependent on each other. Modules were not designed to be independent. Modules split up functional tasks. When a module called another, certain assumptions were being made. As a result, a program modification to one module may have indirectly affected another. Program modifications could not be performed without analyzing many modules to evaluate the impact of the change. The whole design was essentially tailored to solving this problem, without much consideration to how future changes could be incorporated.

It is probable that as the program was modified it's complexity would go up and it's maintainability down. Though the first design was never carried out to completion, it was clearly inferior to the final structured design.

The structured design consisted of 30 modules. It seemed to be simpler, more straight-forward, and more general approach to the problem. These modules were generally much more functional in nature. In fact, ten of the 30 modules were so generally defined that each was used by three to eight calling modules. A new maintenance programmer could more readily understand the design since the structured design was more functionally broken down.

Another aspect of the structured design was that variables were initialized at the lowest possible level. In the first design, nearly all initializations were performed by one module. Initialization in the first module where the variable is needed is more desirable, because when that module is modified it is possible an initialization may have to change. Since the modules are so far removed, the change might easily be overlooked.

Data coupling of modules was used for the preprocessor interface in all but four cases. In three of the cases, COMMON coupling was chosen over data coupling because in each case only two modules needed to share the data and these modules were separated by several module levels. COMMON coupling of the two modules was believed better than passing the data along as arguments through many modules

that had no need for the information. By COMMON coupling the two modules, it prevented other modules from having needless access to the data.

The other exception to data coupling involved the passing of a control flag as a subroutine argument. This flag was passed down to the BUILD-SUBROUTINE-CALLS module. It indicated the subroutine call being built was the final assignment for the arithmetic statement, and no temporary variable would need to be created for storing the result.

No other control flags were used specifically. Alternate returns from subroutines were used in place of control flags. These alternate returns were used for end-of-file conditions and for signaling the detection of an illegal arithmetic statement. The alternate returns for an illegal arithmetic statement were used to avoid further syntax analysis of the statement. Non-arithmetic statements do not need to be converted to Polish Notation form. Myers (REF 6: 51), however, does not regard this usage of alternate returns as control flags since the return indicates to the superordinate module that the function failed. Myers regards this type of return code as data since it does not tell the calling module what to do, but says 'I've failed' and the calling module can respond however it wants.

One other design improvement over the first design, resulting from the structured design approach, was better scope of effect/scope of control. Scope of effect/scope of control essentially rationalizes that all modules effected

by a decision should be subordinate to (beneath) the module making the decision. A consequence of disregarding the scope of effect/scope of control rule is that decisions are repeated as control flags are passed back to the superordinate to be retested.

The preprocessor design kept the scope of effect for a decision within the scope of control of the module making the decision. Under the first design of the preprocessor programs, the detection of an illegal arithmetic statement resulted in passing a control flag back to the superordinate module which output the statement for that subordinate's decision and several other subordinate modules which detected syntax errors. In the structured design, the statement was output by the module detecting the illegal syntax.

Testing The Quality Of The Design

An important test for rating the quality of a design is the ease with which modifications can be made to the existing program. For one proposed change to the program, the design was found quite adaptable. This change is described in the following subsection. The change is explained, then the modules and interfaces affected are listed. In addition, the impact of the change upon the quality of the present program is evaluated.

One change was to replace all subroutine calls built by the efferent branch with in-line code that would do the n-bit simulation effects. The in-line code would be tailored

to the user options in effect for that particular routine. This would mean fewer decisions would be made (such as testing if the rounding effects option is on). Also subroutine linkage time would be saved since the n-bit simulation code was no longer performed by a subroutine. This would result in improved execution time for the n-bit simulation of an algorithm. This proposed change is described completely in Appendix D, as a future enhancement to the system.

Basically, this change would require modifications to three of the current modules: BUILD-SUBROUTINE-CALLS, OUTPUT-SINGLE-ASSIGNMENT, and NON-EXPRESSION-HANDLER. Also two new modules would need to be developed. The modifications to BUILD-SUBROUTINE-CALLS and OUTPUT-SINGLE-ASSIGNMENT would simply be the replacement of one block of statements within each one. This block of statements presently merges portions of the function subroutine being built. The new code introduced within these routines would merge these same parts already present in a different way and call a new module which would insert these parts into standard positions to build the appropriate series of in-line statements for n-bit simulation.

Four additional interfaces would be introduced into the present code; three of them to join the new modules to their superordinate modules and the other would interface the two new modules through the use of COMMON coupling. No present interfaces would be changed and the strength of each module modified would not be decreased. The changed module,

NON-EXPRESSION-HANDLER, presently detects key statements and calls the appropriate handling routine. The CALL SETNBIT subroutine reference could be detected using the same technique, then a new module to handle the statement could be called. This new module would analyze the arguments of the SETNBIT subroutine and compute the key values normally computed by the SETNBIT subroutine of the present system, then pass the keys as data through an interface to the other new module created for this enhancement.

All changes for this program enhancement could be accomplished quickly, easily with firm confidence that the system would be just as functional and easy to maintain as before.

Conclusion

In conclusion, structured design techniques were found to be valuable in the creation of a better design. The final design had many significant advantages over the previous design which relied on no actual methodology. The structured design resulted in a simpler system since its structure resembled the problem structure and the modules were functional in nature. Simplicity is an important design objective. A program that is simple and easier to understand has a more positive effect on the future maintenance and modification costs of the program.

Even though the structured design analysis took considerably more time, it was felt the time would be regained

in the succeeding stages of the preprocessor system life-cycle. The integration of the preprocessor modules and the debugging went quickly and smoothly once the design was coded. The testing and the results obtained are discussed in the next chapter.

Based on the preprocessor program, the structured design techniques were found to be very applicable to this small program design. Perhaps the cost of poor designing is not as great for a small program as for a large program, but small programs comprise a large percentage of the programming systems around the world. Therefore, the overall impact of structured design could be even greater than applying it to only large systems. The next chapter will review the result from tests performed upon the preprocessor developed from the structured design and the other parts of the n-bit simulation tool.

V. Testing And Results

This chapter reviews the tests performed on the two portions of the n-bit simulation tool, the preprocessor and the n-bit simulation subroutine. It then examines some of the results obtained from two examples of FORTRAN algorithms to demonstrate some of the effects obtainable through the n-bit simulation tool.

Preprocessor

The objective of the preprocessor is to translate a given FORTRAN algorithm into its equivalent n-bit simulation form. This process converts all syntactically correct arithmetic statements into a sequence of one or more function subroutine calls.

A test program including various forms of each FORTRAN statement type was developed to verify that non-arithmetic statements would pass through the preprocessor unaffected and that virtually any form of an arithmetic statement would be correctly translated into the intended sequence of function subroutine calls. This sample program is shown in Figure 30. The n-bit simulated version of this program produced by the preprocessor is shown in Figure 31.

All statement types were processed correctly by the preprocessor. The sequence of function calls for each arithmetic statement reflected the proper precedence order of evaluation for the operators and operands involved. (A comment statement containing the original statement preceded

→	PROGRAM PREPERS(INPUT, OUTPUT (TAPE1=OUTPUT))	000100
	IMPLICIT INTEGER (A-Z)	000110
C	IMPLICIT INTEGER IS NOT ALLOWED IN THE N-BIT SIMULATOR	000120
C	NOR ARE INTEGER OR REAL VARIABLE DECLARATIONS, VARIABLE	000130
C	TYPE SHOULD USE DEFAULT VALUES I-N FOR INTEGER, AND	000140
C	A-H, 0-7 FOR REALS	000150
→	DIMENSION C(20), N(15,2), K(15), KEY(7)	000160
	COMMON A, P, Q, R	000170
	INTEGER A2, A4	000180
	EQUIVALENCE (C5, C6), (I1, O6)	000190
	DATA C/20*1.57, BLNK/10	000200
C	THIS IS ONLY A TEST PROGRAM TO EXAMINE THE PREPROCESSORS	000210
C	HANDLING OF THE VARIOUS STATEMENT TYPES	000220
C	A "CALL SETNBIT" MUST BE THE FIRST EXECUTABLE STATEMENT WITHIN	000224
C	EACH ROUTINE AND KEY(7) MUST BE DIMENSIONED FOR EACH ROUTINE	000225
C	IN THE N-BIT SIMULATED PROGRAM	000226
→	CALL SETNBIT(24, 1, 1, 1, 1, 1, 6, 0, KEY)	000230
	READ *, C5, I1	000240
C	ALL EXTERNAL VALUES COMING INTO A N-BIT SIMULATED ROUTINE	000241
C	MUST BE SET EQUAL TO THEMSELVES- IN ORDER TO GIVE EACH ITS N-BIT	000242
C	SIGNIFICANCE	000243
→	C5=C5	000250
→	I1=I1	000250
C	EVERY EXTERNAL VARIABLE MUST BE ASSIGNED TO ITSELF TO	000270
C	N-BIT SIMULATE ITS ACCURACY	000290
	READ (1,3) (K(L), L=1, 15)	000290
3	FORMAT(15F0.3)	000300
C	THE ARRAY K (EXTERNAL VALUES) MUST BE SET EQUAL TO ITSELF TO	000305
C	GIVE IT ITS N-BIT WORDLENGTH SIGNIFICANCE	000306
→	DO 4 L=1, 15	000310
→	K(L)=K(L)	000320
→	CONTINUE	000330
	M=(23+3*(PRV-1000))	000340
	DO 5 M=1, 15	000350
	C(M)=K(M)*2	000360
	CONTINUE	000370
	PRINT *, " INPUT "	000380
	PRINT *, K	000390
	A1=K(1)+78524/K(2)	000400
	CALL SUB1(25, C, A1)	000410
*	A2=1+2/3*I+C(12)	000420
	CALL MANIPUL(A, P, X, I), RETURNS(13)	000430
	DEFIND 1	000440
	ASSIGN 13 TO ILABEL	000450
	DO 7 L=1, 15	000460
	C(L)=C(L, 1)+2+.37	000470
	CONTINUE	000480
	GO TO (13, 15), 4TLABEL	000490
70	IF(I1.EQ.0) GO TO 9	000500
	WRITE (1, 10) A1, (C(M), M=1, 20, 2)	000510
19	FORMAT(2X, =A10)	000520
	GO TO 13	000530
9	PRINT *, " STARTED LAST PART "	000540
13	B1=C(20)*C(19)+K(4)*(-1.7)	000550
15	A2=C(1)*(A1+K(4))	000560
	(A1-5)+1-000	000570
	STOP " PREPROCESSOR TEST COMPLETE "	000580
	END	000590
	SUBROUTINE SUB1(I1, C, A1)	000600
	COMMON A, P, Q, R	000610
	DIMENSION C(1), KEY(7)	000620
	CALL SETNBIT(24, 1, 1, 1, 1, 1, 6, 0, KEY)	000630

Figure 30 Sample Program

C1=C(I1)/3.14235673	000720
Y=.5	000730
RESULT=SIGN(X)*A1**3.5	000740
NEXT=IFUNC(A1,X,C1*15.3/A1)	000750
I1=-I1+A1-C(I1)-4	000760
IF(I1.E7.4) GO TO 17777	000770
IF(I1.E7.3) A1=NEXT**(-2)+I1**(-2)	000780
17777 RESULT=RESULT+0.5	000790
M=I1+40	000800
SUM=0.	000810
DO 19 K=1,4	000820
C(1)=C(4)+RESULT*3.1454	000830
SUM=C(4)+SUM	000840
IF(SUM.GT.2000.) SUM=SUM/A1-100	000850
19 CONTINUE	000860
RETURN	000870
END	000880
FUNCTION IFUNC(A1,A2,A3)	000890
DIMENSION KEY(7)	000900
CALL SETNBIT(2,1,1,1,17,6,0,KEY)	000910
IFUNC=0	000920
IF(A1.EQ.SHIFT(A2,5).AND.A2.EQ.5.9)	000930
IFUNC=1	000940
RETURN	000950
END	000960
SUBROUTINE MANIPUL(A,B,C,K),RETURNS(N2)	000970
COMMON/NEXT/ NOTHING	000980
DIMENSION KEY(7)	000990
DIMENSION A(14),	001000
C1(4),	001010
C2(6),	001020
LARRAY(29),	001030
MARRAY(23,2)	001040
LOGICAL L2	001050
L2=.TRUE.	001060
IF(C1(2).EQ.K) L2=.FALSE.	001070
I(1,1,1)=25+5-7*K-LARRAY(K)+MARRAY(K,1)	001080
17 IF((LARRAY(7)-1).EQ.LARRAY(7+1)) LARRAY(7)=0.1	001090
RETURN	001100
END	001110
INTEGER FUNCTION IFUNC2(I1)	001120
DIMENSION KEY(7)	001130
CALL SETNBIT(2,1,1,1,17,6,1,KEY)	001140
IFUNC2=FUNC3(I1**2+I1,I1,6)	001150
IF(IFUNC2.GT.100) IFUNC2=0	001160
RETURN	001170
END	001180
REAL FUNCTION FUNC3(I,I2,I3)	001190
DIMENSION KEY(7)	001200
CALL SETNBIT(12,1,3,1,7,5,1,KEY)	001210
FUNC3=I/I2/I3 (I2+I3*(4.3+I2))	001220
IF(FUNC3.LT.7.) GO TO 40	001230
FUNC3=FUNC3*FUNC3	001240
RETURN	001250
END	001260
SUBROUTINE FOROPS(QUINCY,DUM1,DUM2)	001270
DIMENSION A(27),P(1,33)	001280
A1=3+3*	001290
RESULT=A=3	001300
A11=NAVEED0.015+3.	001310
A111111=123.5673	001320
RESULT=IND(2,5,P*3,3)+Y*TOTAL,F,G,H,SEN(X),I)	001330
GOVOK=L/3.567	001340
T11=(4X+3)*(4.3+AA4)	001350
A111111=123.5673+127.5677	001360
A(2,3)=A1+7/57	001370
A111111+2+1+K-1+7+3-1+1)=10.0+A111111*ARGUMENT	001380
RETURN	001390
END	001400

Figure 30 (Continued)

```

PROGRAM PREPRSS(INPUT,OUTPUT,TAPE1=OUTPUT)                                000100
IMPLICIT INTEGER (A-Z)                                                    000110
C IMPLICIT INTEGER IS NOT ALLOWED IN THE N-BIT SIMULATOR                  000120
C   NOR ARE INTEGER OR REAL VARIABLE DECLARATIONS, VARIABLE              000130
C   TYPE SHOULD USE DEFAULT VALUES I-N FOR INTEGER,AND                  000140
C   A-H,I-Z FOR REALS                                                    000150
DIMENSION C(20),D(15,2),K(15),KEY(7)                                     000160
COMMON A,B,DDDD                                                         000170
INTEGER A2,A4                                                            000180
EQUIVALENCE (C5,75),(I1,06)                                           000190
DATA C/20*1.7,3LNK/10H /                                              000200
C THIS IS ONLY A TEST PROGRAM TO EXAMINE THE PREPROCESSORS              000210
C   HANDLING OF THE VARIOUS STATEMENT TYPES                             000220
C   A " CALL SETBIT" MUST BE THE FIRST EXECUTABLE STATEMENT WITHIN      000230
C   EACH ROUTINE AND KEY(7) MUST BE DIMENSIONED FOR EACH ROUTINE       000240
C   IN THE N-BIT SIMULATED PROGRAM                                       000250
CALL SETBIT(2,1,1,1,1,17,6,0,KEY)                                       000260
READ *,C5,I1                                                            000270
C   ALL EXTERNAL VALUES COMING INTO A N-BIT SIMULATED ROUTINE          000280
C   MUST BE SET EQUAL TO THEMSELVES-IN ORDER TO GIVE EACH ITS N-BIT     000290
C   SIGNIFICANCE                                                         000300
C   C5=C5                                                                000310
C   C5=ASGN(C5,74HPREPRSS, 22,0,KEY)                                     000320
C   I1=I1                                                                000330
C   I1=IASGN(I1,74HPREPRSS, 23,0,KEY)                                     000340
C   EVERY EXTERNAL VARIABLE MUST BE ASSIGNED TO ITSELF TO              000350
C   N-BIT SIMULATE ITS ACCURACY                                          000360
READ (1,2) (K(L),L=1,15)                                              000370
FORMAT(15F9.3)                                                         000380
C   THE ARRAY K (EXTERNAL VALUES) MUST BE SET EQUAL TO ITSELF TO      000390
C   GIVE IT ITS N-BIT WORDLENGTH SIGNIFICANCE                           000400
DO 4 L=1,15                                                            000410
C   K(L)=K(L)                                                            000420
C   K(L)=IASGN(K(L),74HPREPRSS, 31,0,KEY)                               000430
CONTINUE                                                                000440
M1=(23+3*(22/-1)*0)                                                    000450
RTTTAA=ROUND(22/1.1,0,74HPREPRSS, 32,0,KEY)                          000460
RTTTAA=ROUND(3,RTTTAA,74HPREPRSS, 33,0,KEY)                          000470
M4=ROUND(22,RTTTAA,74HPREPRSS, 33,1,KEY)                              000480
DO 5 M=1,15                                                            000490
C   C(M)=K(M)*2                                                         000500
C   C(M)=ROUND(K(M),2,74HPREPRSS, 35,1,KEY)                            000510
CONTINUE                                                                000520
PRINT *, " INPUT "                                                     000530
PRINT *,K                                                                000540
C   A1=K(1)+75524/K(2)                                                  000550
RTTTAA=ROUND(75524,K(2),74HPREPRSS, 39,0,KEY)                        000560
A1=IASGN(K(1),RTTTAA,74HPREPRSS, 39,1,KEY)                            000570
CALL SUB(27,0,A1)                                                       000580
M4=M4/2+1+D(12)                                                        000590
RTTTAA=ROUND(1,0,74HPREPRSS, 41,0,KEY)                                000600
RTTTAA=ROUND(RTTTAA,0,74HPREPRSS, 41,0,KEY)                            000610
RTTTAA=ROUND(RTTTAA,1,74HPREPRSS, 41,1,KEY)                            000620
A4=ROUND(RTTTAA,D(12),74HPREPRSS, 41,1,KEY)                            000630
CALL MODUL(4,0,Y,I),RETURNS(13)                                         000640
RTTTAA=1                                                                000650
ASSIGN 13 TO LABEL                                                      000660
DO 7 L=1,15                                                            000670
C   C(L)=C(L)+2.37                                                       000680
C   C(L)=ROUND(C(L),2,74HPREPRSS, 40,1,KEY)                            000690
CONTINUE                                                                000700
C   TO (13,1),ATLABEL                                                  000710

```

Figure 31 N-bit Simulated Version Of Sample Program

30	IF(I1.EQ.0) GO TO 9	000720
	WRITE (1,19) A1,(C(NN),NN=1,20,2)	000730
19	FORMAT(2X,=A1))	000740
	GO TO 13	000750
9	PRINT *, " STARTED LAST PART"	000760
13	CONTINUE	000770
C	A1=C(20)*C(19)*K(-)**(-1.7)	000780
	RTTTTAA=RRMPY(C(20),C(19),7HPRPRSS, 5,0,KEY)	000790
	RTTTTAA=RIEXP(K(4),-1.7,7HPRPRSS, 5,0,KEY)	000800
	A1=RRMPY(RTTTAA,RTTTA,7HPRPRSS, 5,1,KEY)	000810
15	CONTINUE	000820
C	A2=C(1)*(A1+K(-))	000830
C	(A1-5)+1.000	000840
	RTTTTAA=R2ADD(A1,K(4),7HPRPRSS, 56,0,KEY)	000850
	RTTTTAA=RRMPY(C(1),RTTTTAA,7HPRPRSS, 56,0,KEY)	000860
	RTTTTAA=R2MNS(A1,7HPRPRSS, 56,0,KEY)	000870
	RTTTTAA=RRMPY(RTTTAA,RTTTTAA,7HPRPRSS, 56,0,KEY)	000880
	A2=R2ADD(RTTTAA,14000,7HPRPRSS, 56,1,KEY)	000890
	STOP " PREPROCESS TEST COMPLETE"	000900
	END	000910
	SUBROUTINE SUB1(I1,C,A1)	000920
	COMMON A,B,DD	000930
	DIMENSION C(1),KEY(7)	000940
	CALL SETNPT(2,-1,1,1,17,6,0,KEY)	000950
C	C1=C(I1)/3.1423578	000960
	C1=RRDVC(C(I1),3.1423578,7HSUB1, 5,1,KEY)	000970
C	X=1.5	000980
	Y=ASGN(C.5,7HSUB1, 5,0,KEY)	000990
C	RESULT=SIGN(X)+A1-3.5	001000
	RTTTTAA=ASGN(X,7HSUB1, 7,0,KEY)	001010
	RTTTTAA=ASGN(SIN(RTTTAA),7HSUB1, 7,0,KEY)	001020
	RTTTTAA=RIEXP(A1,7.5,7HSUB1, 7,0,KEY)	001030
	RESULT=RRMPY(RTTTAA,RTTTTAA,7HSUB1, 7,1,KEY)	001040
C	NEXT=IFUNC(A1,X,C1*15.3/A1)	001050
	RTTTTAA=ASGN(X,7HSUB1, 3,0,KEY)	001060
	RTTTTAA=ASGN(A1,7HSUB1, 3,0,KEY)	001070
	RTTTTAA=RRMPY(C1,15.3,7HSUB1, 3,0,KEY)	001080
	RTTTTAA=RRDVC(RTTTAA,A1,7HSUB1, 3,0,KEY)	001090
	NEXT=IASGN(IFUNC(RTTTAA,RTTTTAA,RTTTTAA),7HSUB1, 3,1,KEY)	001100
C	T1=-I1+A1-C(I1)-4	001110
	RTTTTAA=R2ADD(-I1,A1,7HSUB1, 9,0,KEY)	001120
	RTTTTAA=R2MNS(RTTTAA,C(I1),7HSUB1, 9,0,KEY)	001130
	T1=R2MNS(RTTTAA,7HSUB1, 9,0,KEY)	001140
	IF(I1.EQ.0) GO TO 1777	001150
C	IF(I1.EQ.0) A1=NEXT*(-2)+I1*(-2)	001160
	IF(I1.EQ.0)	001170
	GO TO 7911	001180
	GO TO 7912	001190
7911	CONTINUE	001200
	RTTTTAA=RIEXP(NEXT,-2,7HSUB1, 11,0,KEY)	001210
	RTTTTAA=RIEXP(I1,-2,7HSUB1, 11,0,KEY)	001220
	A1=I1ADD(RTTTAA,RTTTTAA,7HSUB1, 11,1,KEY)	001230
7912	CONTINUE	001240
47777	CONTINUE	001250
C	RESULT=RESULT+.5	001260
	RESULT=R2ADD(RESULT,0.5,7HSUB1, 12,1,KEY)	001270
C	NEXT=1.0	001280
	NEXT=I1ADD(I1,0.5,7HSUB1, 13,1,KEY)	001290
C	END=)	001300
	SUB=ASGN(0.,7HSUB1, 1,0,KEY)	001310
	DO 19 K=1,4	001320
C	C(4)=C(4)+RESULT*.1454	001330
	RTTTTAA=RRMPY(RESULT,0.1454,7HSUB1, 14,0,KEY)	001340
	C(4)=R2ADD(C(4),RTTTTAA,7HSUB1, 14,0,KEY)	001350
C	END=END+.5	001360
	RTTTTAA=RRDVC(END,0.5,7HSUB1, 15,1,KEY)	001370

Figure 31 (Continued)

C	IF(SUM.GT.2000.)	SUM=SUM/A1-100	001380
	IF(SUM.GT.2000.)		001390
	GO TO 7913		001400
	GO TO 7914		001410
7913	CONTINUE		001420
	RTTTAA=RRVDS(SUM,A1,7HSUB1 , 18,C,KEY)		001430
	SUM=R2MNS(RTTTAA,100,7HSUB1 , 18,1,KEY)		001440
7914	CONTINUE		001450
19	CONTINUE		001460
	RETURN		001470
	END		001480
	FUNCTION IFJNC(A1,A2,A3)		001490
	OTIENSION KEY(7)		001500
	CALL SETNBIT(24,1,1,1,17,6,0,KEY)		001510
C	IFJNC=0		001520
	IFUNC=IASGN(0,7HIFUNC , 4,0,KEY)		001530
C	IF(A1.EQ.SHIFT(A2,5).AND.A2.EQ.5.9)		001540
C	IFUNC=1		001550
	IF(A1.EQ.SHIFT(A2,5).AND.A2.EQ.5.9)		001560
	GO TO 7911		001570
	GO TO 7912		001580
7911	CONTINUE		001590
	IFJNC=IASGN(1,7HIFUNC , 6,0,KEY)		001600
7912	CONTINUE		001610
	RETURN		001620
	END		001630
	SUBROUTINE MANIPUL(A,3,C,K), RETURNS(N2)		001640
	COMMON/NEXT/ NOTHING		001650
	OTIENSION KEY(7)		001660
	OTIENSION A(15),	LARRAY(29),	I(5,3,3),
	C1(4),	C2(8),	HARRAY(23,2)
	LOGICAL L2		001690
	L2=.TRUE.		001700
	IF(C1(2).EQ.K) L2=.FALSE.		001710
C	I(1,1,1)=25+5-7*K LARRAY(K)*HARRAY(K,1)		001720
	ITTTAA=I1ADD(25,5,7HMANIPUL , 9,0,KEY)		001730
	ITTTAB=IIMPY(7,K,7HMANIPUL , 9,0,KEY)		001740
	ITTTAC=IIMPY(ITTTAA,LARRAY(K),7HMANIPUL , 9,0,KEY)		001750
	ITTTAD=IIMPY(ITTTAC,HARRAY(K,1),7HMANIPUL , 9,0,KEY)		001760
	I(1,1,1)=I1MNS(ITTTAA,ITTTAD,7HMANIPUL , 9,1,KEY)		001770
1	CONTINUE		001780
C	IF((LARRAY(7)-1).EQ.LARRAY(7+1)) LARRAY(7)=0.1		001790
	IF((LARRAY(7)-1).EQ.LARRAY(7+1))		001800
	GO TO 7911		001810
	GO TO 7912		001820
7911	CONTINUE		001830
	LARRAY(7)=IASGN(0.1,7HMANIPUL , 10,0,KEY)		001840
7912	CONTINUE		001850
	RETURN		001860
	END		001870
	INTEGER FUNCTION IFUNC2(I1)		001880
	OTIENSION KEY(7)		001890
	CALL SETNBIT(24,1,1,1,17,6,1,KEY)		001900
C	IFUNC2=IFUNC2(I1+11,I1,5)		001910
	ITTTA4=IIMPY(I1,4,7HIFUNC2 , 4,1,KEY)		001920
	ITTTA1=I1ADD(ITTTA4,I1,7HIFUNC2 , 4,0,KEY)		001930
	ITTTA2=IASGN(I1,7HIFUNC2 , 4,0,KEY)		001940
	ITTTA3=IASGN(5,7HIFUNC2 , 4,0,KEY)		001950
	IFUNC2=IASGN(1,IFUNC2(ITTTA4,ITTTA2,ITTTA3),7HIFUNC2 , 4,1,KEY)		001960
C	IF(IFUNC2(11,5)) IFUNC2=0		001970
	IF(IFUNC2(11,5))		001980
	GO TO 7911		001990
	GO TO 7912		002000
7911	CONTINUE		002010
	IFUNC2=IASGN(0,7HIFUNC2 , 4,1,KEY)		002020
7912	CONTINUE		002030

Figure 31 (Continued)

the n-bit simulation function calls generated). For example, the arithmetic statement in line 420 of Figure 30 translates into the statements shown in lines 590 through 630 of Figure 31. In this example, statements that would not normally appear if n-bit simulation were not performed are designated with a pointer in Figure 30.

Certain extremes were introduced into the original program to demonstrate the versatility and range of statements that could be handled by the preprocessor. A few of these extreme cases shown in Figure 31 were continuation lines (line 830), function references with greater than six arguments (line 2270), long array argument lists (line 2560), long integer and real numbers (line 2490), and multiple IF statements within a given routine (lines 920-1490). The latter case was used to verify that unique labels were being generated from a series of modified IF statements.

N-bit Simulation Subroutines

The preprocessor produced an n-bit simulated algorithm which referenced up to 22 different n-bit simulation function subroutines. Testing these n-bit simulation subroutines consisted of two phases. The first phase validated the key values which were computed by the SETNBIT subroutine. The second phase involved testing the 22 n-bit simulation function subroutines to verify the n-bit wordlength effects and other user options were being properly performed.

SETNBIT Key Values Verification

A critical portion of testing the n-bit simulation tool involved verifying the correctness of the key values computed by the SETNBIT subroutine. This is important because the performance of all 22 n-bit simulation subroutines are based upon the computed key values. The seven key values produced by the SETNBIT subroutine, based on the user options, are listed below:

- 1) Maximum positive fixed point value possible
- 2) Maximum positive floating point value possible
- 3) Least significant positive floating point value possible
- 4) Rounding or truncation effects flag
- 5) Overflow message printing or suppression flag
- 6) Number of least significant bits (LSB) to be truncated from floating point values for final assignments
- 7) Number of LSB to be truncated from floating point values for intermediate assignments

Figure 32 shows key values computed for various user options. Figure 32(a) has specifications for a 24-bit floating point word with 17 mantissa bits, six exponent bits, and an implied decimal point to the left of the mantissa. In a 24-bit word, the maximum floating point value representable would look like:

011111111111111111111111
sign expo- mantissa
bit nent implied decimal point

AD-A055 777

AIR FORCE INST OF TECH WRIGHT-PATTERSON AFB OHIO SCH--ETC F/G 9/2
SOFTWARE TOOL(S) FOR EVALUATING THE EFFECTS OF FINITE WORDLENGT--ETC(U)
DEC 77 G A KLEIN

UNCLASSIFIED

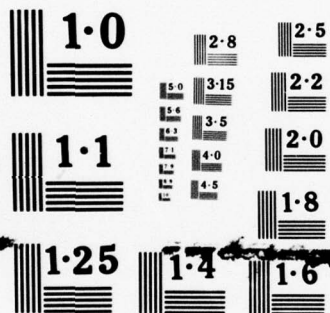
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(a)

N-BIT SPECIFICATIONS ARE:
NUMBER OF BITS/WORD = 24

TRUNCATE (0) OR ROUND (1) = 0

PRECISION = 1

OVERFLOW MESSAGES (1) / NO MESSAGES (0) = 1

MANTISSA BITS = 17

EXPONENT BITS = 6

DECIMAL PT LEFT (0) OR RIGHT (1) = 0

KEY VALUES IN OCTAL

KEY(1) = 00000000000037777777
KEY(2) = 17567777760000000000
KEY(3) = 16604000000000000000
KEY(4) = 00000000000000000000
KEY(5) = 00000000000000000001
KEY(6) = 00000000000000000037
KEY(7) = 00000000000000000037

(IN DECIMAL)

KEY(1) = MAXIMUM INTEGER = 8388607
KEY(2) = MAXIMUM FLOATING POINT VALUE POSSIBLE = $2.147467264E+9$
KEY(3) = LEAST SIGNIFICANT FLOATING PT POSSIBLE = $2.328306436539E-10$
KEY(4) = ROUNDING (1) / TRUNCATION (0) EFFECTS = 0
KEY(5) = OVERFLOW MESSAGE PRINT (1) / SUPPRESS (0) = 1
KEY(6) = # BITS TRUNCATED FROM RIGHT OF FLTG PT FOR
FINAL ASSIGNMENT = 31
KEY(7) = # BITS TRUNCATED FROM RIGHT OF FLTG PT FOR
INTERMEDIATE ASSIGNMENT = 31

Figure 32 Key Values For Various User Options

(b)

N-BIT SPECIFICATIONS ARE:
NUMBER OF BITS/WORD = 24

TRUNCATE (0) OR ROUND (1) = 0

PRECISION = 1

OVERFLOW MESSAGES (1) / NO MESSAGES (0) = 1

MANTISSA BITS = 17

EXPONENT BITS = 6

DECIMAL PT LEFT (0) OR RIGHT (1) = 1

KEY VALUES IN OCTAL

KEY(1) = 00000000000037777777
KEY(2) = 20007777760000000000
KEY(3) = 17014000000000000000
KEY(4) = 00000000000000000000
KEY(5) = 00000000000000000001
KEY(6) = 00000000000000000037
KEY(7) = 00000000000000000037

(IN DECIMAL)

KEY(1) = MAXIMUM INTEGER = 8388607

KEY(2) = MAXIMUM FLOATING POINT VALUE POSSIBLE = $2.81472829227E+14$

KEY(3) = LEAST SIGNIFICANT FLOATING PT POSSIBLE = .000030517578125

KEY(4) = ROUNDING (1) / TRUNCATION (0) EFFECTS = 0

KEY(5) = OVERFLOW MESSAGE PRINT (1) / SUPPRESS (0) = 1

KEY(6) = # BITS TRUNCATED FROM RIGHT OF FLTG PT FOR
FINAL ASSIGNMENT = 31

KEY(7) = # BITS TRUNCATED FROM RIGHT OF FLTG PT FOR
INTERMEDIATE ASSIGNMENT = 31

Figure 32 (Continued)

(c) .

N-BIT SPECIFICATIONS ARE:
NUMBER OF BITS/WORD =24

TRUNCATE (0) OR ROUND (1) =1

PRECISION = 2

OVERFLOW MESSAGES (1) / NO MESSAGES (0) =1

* MANTISSA BITS =17

* EXPONENT BITS =6

DECIMAL PT LEFT (0) OR RIGHT (1) = 0

KEY VALUES IN OCTAL

KEY(1)= 000000000000037777777
KEY(2)= 17567777777777777600
KEY(3)= 16604000000000000000
KEY(4)= 00000000000000000001
KEY(5)= 00000000000000000001
KEY(6)= 00000000000000000037
KEY(7)= 00000000000000000007

(IN DECIMAL)

KEY(1)= MAXIMUM INTEGER = 9388607

KEY(2)= MAXIMUM FLOATING POINT VALUE POSSIBLE = 2.147483647999E+9

KEY(3)= LEAST SIGNIFICANT FLOATING PT POSSIBLE = 2.328306436539E-10

KEY(4)= ROUNDING (1) / TRUNCATION (0) EFFECTS = 1

KEY(5)= OVERFLOW MESSAGE PRINT (1) / SUPPRESS (0) = 1

KEY(6)= * BITS TRUNCATED FROM RIGHT OF FLT6 PT FOR
FINAL ASSIGNMENT = 31

KEY(7)= * BITS TRUNCATED FROM RIGHT OF FLT6 PT FOR
INTERMEDIATE ASSIGNMENT = 7

Figure 32 (Continued)

The maximum exponent possible for six bits is ± 31 , assuming an exponent bias of 100000_2 . Therefore, the value shown would be equal to the contents of the mantissa times 2^{+31} . The mantissa above represents the fractional value of $(2^{17} - 1)/2^{17}$. When multiplied out, this fraction times 2^{+31} equals 2147467264., which is consistent with the value computed by the SETNBIT routine for KEY(2).

The least significant floating point value possible would look like

0000000^A100000000000000000

The mantissa in this case is equal to .5 and the exponent is equal to -31 (with respect to the bias 100000_2). So the least significant number possible is .5 times 2^{-31} or $2.328306436539 \times 10^{-10}$. This is consistent with the value computed for KEY(3) by SETNBIT.

The verification of the fixed point maximum can be accomplished by looking at the octal representation of KEY(1) in Figure 32(a). The least significant 23 bits are ones; this translates to 8388607 for its equivalent decimal representation. In a 24-bit fixed point word, the largest value would have 1's in the least significant 13 bit positions.

Figure 32(b) has the same specifications as (a) with the exception that the implied decimal point is to the right of the mantissa. The maximum floating point number possible for this representation is the maximum mantissa, which has a fixed point value of $2^{17} - 1$ or 131071, times 2^{+31} ,

yielding 2.81472829227 times 10^{14} . This value is also exactly the same as the computed value for KEY(2) of Figure 32(b). The least significant floating point value also computes to the exact value shown in KEY(3) of (b). The maximum fixed point value does not change with the moving of the decimal point since only floating point values are affected.

In Figure 32(c), double precision arithmetic accuracy was specified. KEY(7) contains the number 7, which is the number of bit positions that should be truncated off the CDC's 48-bit mantissa. This means 41 significant mantissa bits are retained. The single precision mantissa of 17 bits (shown in KEY(6)) plus extra wordlength of 24 bits equals 41.

Testing Of The N-bit Simulation Subroutines

A unique function subroutine was developed for each arithmetic operation and each possible combination of input variable data types. Twenty of the 22 function subroutines developed perform arithmetic operations and two perform simple assignments.

The 20 function subroutines can be divided into two general classes: those which produce floating point values and those which produce fixed point values. The subroutines of a given class basically differ only in the first statement they execute. That statement performs the arithmetic operation being simulated. After which, they are constructed the same. For instance, for the operations $A*B$ and $C+D$ the first statement in each handling subroutine would perform the required operation upon the operands and store the result

in the variable RESULT. From there on, exactly the same n-bit simulation effects are performed upon RESULT for both subroutines. Therefore, verification of the results produced by each combination of user options for one subroutine of each class should sufficiently prove the other subroutines of that class perform correctly also.

Figure 33 shows the results from two sets of fixed point and floating point additions. For each case within a set, the octal representation of the two values being added are shown by INPUT1 and INPUT2. The NORMAL OUTPUT represents the result of the computation using the full 60-bit precision of the CDC computer. The N-BIT OUTPUT represents the value resulting from the n-bit simulation of the two inputs being added together. If an overflow condition occurs, a message is printed (e.g. case 1C), and the maximum value possible replaces the old value. The key values precede each case. Case one is executed for the options: 16-bits, rounding, single precision accuracy, overflow messages, nine mantissa bits, and six exponent bits with the implied decimal point to the left of the mantissa. Cases 1A through 1H were performed with fixed point addition. Cases 1I through 1P were performed with floating point addition. Case 1O shows an example where underflow for the 16-bit word was detected. Case 1N shows that 9 bits of significance were maintained in the mantissa and also shows a case of round up.

In order to verify correct detection of overflow, many borderline cases were examined (e.g. cases 1A through 1D).

SET 1

```

KEY(1)= 00000000000000007777 MAXIMUM FIXED POINT
KEY(2)= 17557770000000000000 MAXIMUM FLOATING POINT
KEY(3)= 16800000000000000000 LEAST SIGNIFICANT BITS POINT
KEY(4)= 00000000000000000001 ROUNDING(1) OF TRUNCATION(0)
KEY(5)= 00000000000000000001 OVERFLOW MESSAGE OPTION
KEY(6)= 00000000000000000147
KEY(7)= 00000000000000000047
  INTEGER INTEGER ADDITION
  MAX-15 + 15
INPUT1=00000000000000007776 INPUT2=00000000000000000017
NORMAL OUTPUT=00000000000000007777
N-BIT OUTPUT=00000000000000007777

  MAX-1 +0
INPUT1=00000000000000007776 INPUT2=00000000000000000000
NORMAL OUTPUT=00000000000000007776
N-BIT OUTPUT=00000000000000007776

  MAX-1 + 2
INPUT1=00000000000000007776 INPUT2=00000000000000000002
NORMAL OUTPUT=00000000000000007778
*** OVERFLOW IN TESTING LINE # 1 FOR ADDITION ***
N-BIT OUTPUT=00000000000000007777

  MIN+3 + -2
INPUT1=77777777777777770000 INPUT2=77777777777777777775
NORMAL OUTPUT=77777777777777770001
N-BIT OUTPUT=77777777777777770001

  MAX + 3000
INPUT1=00000000000000007777 INPUT2=00000000000000000070
NORMAL OUTPUT=00000000000000007777
*** OVERFLOW IN TESTING LINE # 1 FOR ADDITION ***
N-BIT OUTPUT=00000000000000007777

  -MAX + -20
INPUT1=77777777777777770000 INPUT2=77777777777777777753
NORMAL OUTPUT=77777777777777770001
*** OVERFLOW IN TESTING LINE # 1 FOR ADDITION ***
N-BIT OUTPUT=77777777777777770000

  MAX + - MAX
INPUT1=00000000000000007777 INPUT2=77777777777777770000
NORMAL OUTPUT=00000000000000000000
N-BIT OUTPUT=00000000000000000000

  MAX-1 + 35
INPUT1=00000000000000007776 INPUT2=00000000000000000045
NORMAL OUTPUT=00000000000000007777
*** OVERFLOW IN TESTING LINE # 1 FOR ADDITION ***
N-BIT OUTPUT=00000000000000007777

```

Figure 33 Examples Of N-bit Simulation Additions

1I RMAX+ 0
 INPUT1=17567770000000000000 INPUT2=00000000000000000000
 NORMAL OUTPUT=17567770000000000000
 N-BIT OUTPUT=17567770000000000000

1J RMAX - A BIG ONE
 INPUT1=17567770000000000000 INPUT2=60220017777777777777
 NORMAL OUTPUT=17564000000000000000
 N-BIT OUTPUT=17564000000000000000

1K RMAX + BIG ENOUGH ONE (OVERFLOW EXPECTED)
 INPUT1=17567770000000000000 INPUT2=17557760000000000000
 NORMAL OUTPUT=17515770000000000000
 **** OVERFLOW IN TESTPG LINE # 2 FOR ADDITION ***
 N-BIT OUTPUT=17567770000000000000

1L -RMAX+ -RBIG
 INPUT1=60210007777777777777 INPUT2=60220017777777777777
 NORMAL OUTPUT=60200000000000000000
 **** OVERFLOW IN TESTPG LINE # 2 FOR ADDITION ***
 N-BIT OUTPUT=60210007777777777777

1M RMAX + A DECIMAL
 INPUT1=17567770000000000000 INPUT2=1717600446722743250
 NORMAL OUTPUT=17567770000000000000
 N-BIT OUTPUT=17567770000000000000

1N 5000.13+ .00145
 INPUT1=17344704041217270244 INPUT2=17065740675512106347
 NORMAL OUTPUT=17344704041515034100
 N-BIT OUTPUT=17344710000000000000
 ↑

1O RMIN*2 + -RMIN*2/2
 INPUT1=16614000000000000000 INPUT2=61171777777777777777
 NORMAL OUTPUT=16674000000000000000
 **** OVERFLOW IN TESTPG LINE # 2 FOR ADDITION ***
 N-BIT OUTPUT=16600000000000000000

1P 125.7163 + 45.7751
 INPUT1=17305015555755764420 INPUT2=17255530145451275170
 NORMAL OUTPUT=17305015555755764420
 N-BIT OUTPUT=17305015555755764420

Figure 33 (Continued)

SET 2

```
KEY(1)= 0000000000000077777 MAXIMUM FIXED POINT  
KEY(2)= 17568777777777777690 MAXIMUM FLOATING POINT  
KEY(3)= 1660400000000000000 LEAST SIGNIFICANT FLTg POINT  
KEY(4)= 0000000000000000000 ROUNDING-1) OR TRUNCATION-0)  
KEY(5)= 0000000000000000000 OVERFLOW MESSAGE OPTION  
KEY(6)= 0000000000000000047  
KEY(7)= 0000000000000000007
```

```

      RMWZ+ 0
      INPUT1=17567777777777777777600 INPUT2=0000000000000000000000
2A      NORMAL OUTPUT=17567777777777777777600
      N-BIT OUTPUT=17567777777777777777600

```

```

2B  SMAX - A BIG ONE
    INPUT1=175677777777777777600 INPUT2=602260177777777777777777
    NORMAL OUTPUT=175640077777777777600
    N-BIT OUTPUT=175640077777777777600

```

```

2C      RMAX + BIG ENOUGH ONE OVERFLOW EXPECTED
      INPUT1=17562777777777777777000 INPUT2=1756778000000000000000
      NORMAL OUTPUT=1757507377777777777700
      *** OVERFLOW IN TESTP6 LINE #      2 FOR ADDITION ***
      4-BIT OUTPUT=1756277777777777777700

```

```

2D      -RMAX+ -PBIG
INPUT1=60210000000000000177 INPUT2=6022001777777777777777
NORMAL OUTPUT=60202004000000000007
*** OVERFLOW IN TESTF6 LINE #      2 FOR ADDITION ***
N-BIT OUTPUT=60210000000000000177

```

```

2E  * RMAX + A DECIMAL
    INPUT1=1756777777777777600 INPUT2=17176630446722749850
    NORMAL OUTPUT=1757400000000019151
    **** OVERFLOW IN TEST96 LINE #      2 FOR ADDITION ****
    N-BIT OUTPUT=1756777777777777600

```

```

5000.134.00145
2F INPUT1=17344704041217270244 INPUT2=17025740675512106347
NORMAL OUTPUT=17344704041515324130
N-BIT OUTPUT=17344704041515324200

```

```

      RMIN=2 + J*(RMIN-2)
2G INPUT1=166140000000000000 INPUT2=6117177777777777777
NORMAL OUTPUT=165740000000000000
*** UNDERFLOW IN TESTPG LINE # 2 FOR ADDITION ***
N-EIT OUTPUT=000000000000000000

```

[illegible]

Figure 33 (Continued)

Also computations which result in values well within the range of the 16-bit word are shown in Set 1.

Set 2 has the same options as Set 1, with the exception of triple precision arithmetic accuracy. Cases 2F and 2H show examples where the resulting value was rounded up. Also these cases show that 41 significant bits of the mantissa are retained for the double precision arithmetic accuracy.

N-bit Simulation Tool Effects

After verifying the two major parts of the system (i.e. preprocessor and n-bit simulation subroutines), the n-bit simulation tool was tested as a whole. The sample FORTRAN algorithm to be n-bit simulated is shown in Figure 34. It performs very simple computations designed to demonstrate losses of significance for shorter wordlengths. The n-bit simulated version of the program (built by the preprocessor) is shown in Figure 35.

The n-bit simulated program was then executed with various options. The results from these various executions are shown in Figure 36 (b), (c), (d), and (e). The results from the execution of the original FORTRAN algorithm without n-bit simulation effects is shown in Figure 36(a).

Figure 36(b) used a 24 bit wordlength with 17 mantissa bits (with the implied decimal point to its left) and 6 bits for the exponent. It was performed with single precision arithmetic accuracy and with rounding effects. The resulting output for TOTAL was 376784. while the original program

PROGRAM TEST4(INPUT,OUTPUT,TAPE1=OUTPUT,TAPE2=INPUT)	000100
DIMENSION K(10),KEY(7)	000110
B1=55210.	000160
B2=7.327	000170
DO 8 I=1,50	000180
TOTAL=R1/B2+TOTAL	000190
CONTINUE	000200
PRINT *," TOTAL=",TOTAL	000210
HALF=TOTAL-TOTAL/2.	000220
PRINT *," HALF TOTAL=",HALF.	000230
X1=.5791123	000240
X2=.5790995	000250
DIFF=X1-X2	000260
WTDIFF=DIFF*1000.	000270
PRINT *," DIFF=",DIFF," WTDIFF=",WTDIFF	000280
X3=.97265	000290
X4=.9857134	000300
X34DIFF=X3-X4	000310
RATIO=WTDIFF/X34DIFF*100	000320
PRINT *," X34DIFF=",X34DIFF," RATIO =",RATIO	000330
FUNC=FUNC1(23456.34,5+45.,1,0,RATIO)*1.25	000340
PRINT *," FUNC=",FUNC	000350
CALL SUB1(RATIO,X34DIFF-.453,DIFF,TOTAL)	000360
PRINT *," SUB RATIO=",RATIO	000370
STOP " TEST4 COMPLETE"	000390
END	000400
FUNCTION FUNC1(A,B,I,J,R)	000410
DIMENSION KEY(7)	000420
R1=A**(-2)+B**(-2)	000450
FUNC1=R1/3000.	000460
IF(FUNC1.LT..00025) FUNC1=.00025	000470
RETURN	000480
END	000490
SUBROUTINE SUB1(R1,R2,R3,R4)	000500
DIMENSION KEY(7)	000510
P1=R1+R2+R3+R4	000540
R1=R1**(-2)	000550
RETURN	000560
END	000570

Figure 34 Example Of An Algorithm

020GR14 TEST4(INPUT,OUTPUT,TAPE1=OUTPUT,TAPE2=INPUT)	000100
014ENSION K(10),KEY(7)	000110
R1=55210.	000120
R1=RASGN(55210.,7HTEST4 , 3,0,KEY)	000130
R2=7.327	000140
R2=RASGN(7.327,7HTEST4 , 4,0,KEY)	000150
00 8 I=1.50	000160
TOTAL=R1/R2+TOTAL	000170
RTTTAA=R2DVO(R1,R2,7HTEST4 , 6,0,KEY)	000180
TOTAL=RRADD(RTTTAA,TOTAL,7HTEST4 , 6,1,KEY)	000190
CONTINUE	000200
PRINT *, " TOTAL=",TOTAL	000210
HALF=(TOTAL-TOTAL/2	000220
RTTTAA=R2DVO(TOTAL,2,7HTEST4 , 9,0,KEY)	000230
HALF=RRMNS(TOTAL,RTTTAA,7HTEST4 , 9,1,KEY)	000240
PRINT *, " HALF=",HALF	000250
X1=.5791123	000260
X1=RASGN(.5791123,7HTEST4 , 11,0,KEY)	000270
X2=.5790995	000280
X2=RASGN(.5790995,7HTEST4 , 12,0,KEY)	000290
DIFF=X1-X2	000300
DIFF=RRMNS(X1,X2,7HTEST4 , 13,1,KEY)	000310
WTDIFF=DIFF*100.	000320
WTDIFF=R2DVO(DIFF,100.,7HTEST4 , 14,1,KEY)	000330
PRINT *, " DIFF=",DIFF," WTDIFF=",WTDIFF	000340
X3=.97265	000350
X3=RASGN(.97265,7HTEST4 , 16,0,KEY)	000360
X4=.9857134	000370
X4=RASGN(.9857134,7HTEST4 , 17,0,KEY)	000380
X3DIFF=X3-X4	000390
X3DIFF=RRMNS(X3,X4,7HTEST4 , 18,1,KEY)	000400
RATIO=WTDIFF/X34DIFF*100	000410
RTTTAA=R2DVO(WTDIFF,X34DIFF,7HTEST4 , 19,0,KEY)	000420
RATIO=R2DVO(RTTTAA,100,7HTEST4 , 19,1,KEY)	000430
PRINT *, " X34DIFF=",X34DIFF," RATIO=",RATIO	000440
FUNC=FUNC1(23455.34,5*45.,1,1,RATIO)*1.25	000450
RTTTAA=R2DVO(5.45.,7HTEST4 , 21,0,KEY)	000460
RTTTAB=RASGN(23455.34,7HTEST4 , 21,0,KEY)	000470
ITTTAA=IASGN(1,7HTEST4 , 21,0,KEY)	000480
ITTTAB=IASGN(0,7HTEST4 , 21,0,KEY)	000490
RTTTAC=RASGN(RATIO,7HTEST4 , 21,0,KEY)	000500
RTTTAD=RASGN(FUNC1(RTTTAA,RTTTAB,ITTTAA,ITTTAB,RTTTAC),	000510
7HTEST4 , 21,0,KEY)	000520
FUNC=RRMNS(RTTTAD,1.25,7HTEST4 , 21,1,KEY)	000530
PRINT *, " FUNC=",FUNC	000540
CALL SUB1(RATIO,X34DIFF-.453,DIFF,TOTAL)	000550
PRINT *, " SUB RATIO=",RATIO	000560
STOP " TEST4 COMPLETE"	000570
END	000580
FUNCTION FUNC1(A,B,I,J,P)	000590
014ENSION KEY(7)	000600
P1=A*(-2)+B*(-2)	000610
RTTTAA=R2DVO(A,-2,7HFUNC1 , 3,0,KEY)	000620
RTTTAB=R2DVO(B,-2,7HFUNC1 , 3,0,KEY)	000630
P1=RRADD(RTTTAA,RTTTAB,7HFUNC1 , 3,1,KEY)	000640
FUNC1=P1/3000.	000650
FUNC1=R2DVO(R1,3000.,7HFUNC1 , 4,1,KEY)	000660
IF(FUNC1.LT..00025) FUNC1=.00025	000670
IF(FUNC1.LT..00025)	000680
*GO TO 7911	000690
*GO TO 7912	000700
914 CONTINUE	000710

Figure 35 N-bit Simulated Algorithm

FUN31=KASGN(.000025,7HFUNC1 , 5,0,KEY)	000720
912 CONTINUE	000730
RETURN	000740
END	000750
SUBROUTINE SUB1(R1,R2,R3,R4)	000760
DIMENSION KEY(7)	000770
P1=R1+R2+R3+R4	000780
PTTTTAA=RRADD(R1,R2,7HSUB1 , 3,0,KEY)	000790
RTTTTAA=RRADD(RTTTAA,R3,7HSUB1 , 3,0,KEY)	000800
R1=RRADD(RTTTAA,R4,7HSUB1 , 3,1,KEY)	000810
R1=R1**(-2)	000820
R1=R2EXP(R1,-2,7HSUB1 , 4,1,KEY)	000830
RETURN	000840
END	

Figure 35 (Continued)

```

24.0,1,1,17.6,0

TOTAL=376688.
HALF=188344.
DIFF=.0000152587890625 MTDIFF=.0152587890625
X34DIFF=-.0130615234375 RATID=-116.8212890625
FUNC=.00003124959766865
**** UNDERFLOW IN SUB1 LINE # 4 FOR EXPONENT ***
SUB RATID=0.

24.1,1,1,17.6,0

TOTAL=376784.
HALF=188392.
DIFF=.0000152587890625 MTDIFF=.0152587890625
X34DIFF=-.0130615234375 RATID=-116.8232421875
FUNC=.00003125068332994
SUB RATID=7.048250871823E-12

24.0,2,1,17.6,0

TOTAL=376688.
HALF=188344.
DIFF=.0000152587890625 MTDIFF=.0152587890625
X34DIFF=-.0130615234375 RATID=-116.822265625
FUNC=.00003124959766865
**** UNDERFLOW IN SUB1 LINE # 4 FOR EXPONENT ***
SUB RATID=0.

16.0,1,0,9.6,0

TOTAL=362496.
HALF=181248.
DIFF=0. MTDIFF=0.
X34DIFF=-.013671875 RATID=0.
FUNC=.0000311362457275
SUB RATID=0.

16.0,3,0,9.6,0

TOTAL=362496.
HALF=181248.
DIFF=0. MTDIFF=0.
X34DIFF=-.013671875 RATID=0.
FUNC=.0000311362457275
SUB RATID=0.

48.0,1,1,41.6,0

TOTAL=376757.1993959
HALF=188378.5996979
DIFF=.00001280000014958 MTDIFF=.012800000014958
X34DIFF=-.01306340000019 RATID=-97.98368073691
FUNC=.00003124999999998
**** UNDERFLOW IN SUB1 LINE # 4 FOR EXPONENT ***
SUB RATID=0.

60.0,1,1,48.11,1

TOTAL=376757.1993959
HALF=188378.5996979
DIFF=.00001279999999998 MTDIFF=.012799999999998
X34DIFF=-.0130634 RATID=-97.98367956899
FUNC=.00003125
SUB RATID=7.04816251027E-12

```

Figure 26 Sample Outputs

obtained a result of 376757.1993995, a difference of approximately +24. When truncation effects were applied (shown in (c)) to the n-bit simulated program, the results were slightly less accurate, TOTAL equalled 376688, a difference of -69. Employing double precision arithmetic accuracy with single precision storage (shown in (d)) did not impact the accuracy. However, double precision 24-bit storage variables, shown in (f), increased the accuracy of this algorithm's computation to almost the exact accuracy accomplished by the original algorithm in (a).

Double precision storage can be simulated by keeping the exponent the same length as in the single precision 24-bit word, but increasing the mantissa and the overall n-bit length by 24 (another full word of accuracy). The only caution a user must exercise in simulating double precision variables in this manner, is that the fixed point maximum values computed for detecting fixed point overflow will not be correct. Therefore, fixed point values that would have resulted in overflow in an n-bit machine may not be detected by the n-bit simulation tool.

The results shown in (g) of Figure 36 are for the full 60-bits of the CDC word. The n-bit simulated 60-bit results were exactly the same as those produced from the original algorithm and the CDC's full 60-bit accuracy.

For another example of n-bit wordlength effects, the small FORTRAN algorithm shown in Figure 37 was n-bit simulated for various combinations of user options. The original

```

PROGRAM TEST1(INPUT,OUTPUT)
SUM=0.
A=.00011579
R=2.123
C=1.0319
DO 5 I=1,100
SUM=A+3*C+SUM
5 CONTINUE
PRINT *, " SUM=",SUM
STOP " TEST1 COMPLETE"
END

SUM=219.6930889999

```

Figure 37 Example Of A Second Algorithm

algorithm resulted in a SUM of 219.693088999. Numerous other bit lengths were tried and their results are shown in Figure 38. Particularly interesting is the varying effect rounding has upon the results of this algorithm at different bit lengths. A 12-bit word with rounding resulted in a SUM equal to 11, while without rounding SUM equalled 128 as shown in (a) and (d) respectively of Figure 38. Figure 38 (b) and (e) show a case where rounding increased the accuracy of the SUM from 204.5 to 213.5 for a 16-bit wordlength. Cases (f), (g), (h), and (i) show the results from various wordlength and exponent, mantissa combinations.

Cases (j) and (k) show user options which were incorrect and flagged as erroneous by the SETNBIT subroutine. Case (j) had too many mantissa and exponent bits specified for a 16-bit wordlength and case (k) specified 64-bits, but 60-bits is the maximum wordlength this n-bit simulation tool can simulate on the CDC machine.

Cases (l) and (m) contrast the results obtainable from this algorithm with a single word storage 24-bit wordlength machine and a 24-bit machine with double word storage capabilities, both involving single precision arithmetic accuracy.

It would be up to the user to determine the best choice of bit length and characteristics to provide sufficient accuracy for the algorithms being considered.

```

N-BIT SPECIFICATIONS ***
NUMBER OF BITS=12

IPONCE (0) OF ALGO (1)=1
PRECISION =1
OVERFLOW MESSAGE=1 (1) AND MESSAGE=1
# MESSAGE BITS =6
# EXPONENT BITS =5
DECIMAL PT LEFT (0) OF RIGHT (1)=0
KEYS= 0000000000000000 175-7700000000000000
1700400000000000 0000000000000000 0000000000000000
0000000000000000 0000000000000000
SUM=11.

N-BIT SPECIFICATIONS ***
NUMBER OF BITS=12

IPONCE (0) OF ALGO (1)=0
PRECISION =1
OVERFLOW MESSAGE=1 (1) AND MESSAGE=1
# MESSAGE BITS =9
# EXPONENT BITS =8
DECIMAL PT LEFT (0) OF RIGHT (1)=0
KEYS= 0000000000000000 175-7700000000000000
15-0000000000000000 0000000000000000 0000000000000000
0000000000000000 0000000000000000
SUM=104.7

```

Figure 38 Results From N-bit Simulation
Of Second Algorithm


```

N-BIT SPECIFICATION: 16
NUMBER OF BITS: 16

TRUNCATE (0) OF ROUND (1)=1
PRECISION =1
OVERFLOW MESSAGE (1) AND MESSAGE (0)=1
(e) # MANTISSA BITS =9
# EXPONENT BITS =6
DECIMAL AT LEFT (0) OF -16-1(1)=0
KEY1= 0000000000000000 1777777777777777
1000000000000000 0000000000000000 1000000000000000
0000000000000000 0000000000000000
SUN=-15.1

N-BIT SPECIFICATION: 16
NUMBER OF BITS: 20

TRUNCATE (0) OF ROUND (1)=1
PRECISION =1
OVERFLOW MESSAGE (1) AND MESSAGE (0)=1
(f) # MANTISSA BITS =12
# EXPONENT BITS =7
DECIMAL AT LEFT (0) OF -16-1(1)=0
KEY1= 0000000000000000 17777777 2017777777777777
1000000000000000 0000000000000000 0000000000000000
0000000000000000 0000000000000000
SUN=-15.15

```

Figure 38 (Continued)

(g)

PRECISION = 1.

* 35-44126-5115-240

$$4. \quad \exists x \in \mathbb{N} \quad \forall y \in \mathbb{N} \quad x + y = 7$$
$$\text{Defining } F(\text{Left}) \text{ or } F(\text{Right})=0$$

(h)

$$T_k(u) = \begin{cases} 0 & \text{if } |u| \geq 1 \\ 1 - |u| & \text{if } |u| < 1 \end{cases}$$

FOUNDED = 1

OVERFLOW MESSAGE(S) TO /dev/null MESSAGE(S) TO=1

• 0901-11214-1173 =43

* EXERCISE SET 10

[illegible]

(i)

$$\text{TRUNCATE TABLE T1 SET TABLESPACE TBS1;$$
$$P_{\text{PCLIS12}} = 1$$

QUESTION: Why is $\sum_{i=1}^n (x_i - \bar{x})^2 = \sum_{i=1}^n x_i^2 - n\bar{x}^2$ ($n \geq 1$)?

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$$H^1(\mathbb{R}^n, \mathbb{R}) \cong \mathbb{R}^n \quad \text{for } n \geq 1.$$
$$[T_1] = [T_2] = \dots = 1 \quad \text{and} \quad [T_1] + [T_2] + \dots + [T_n] = 1 \quad \text{and} \quad [T_1] + [T_2] + \dots + [T_n] = 1$$

19.693.83.8533

105

N-BIT SPECIFICATIONS ARE:
 NUMBER OF BITS/WORD = 16
 TRUNCATE (0) OR ROUND (1) = 1
 PRECISION = 1
 OVERFLOW MESSAGE(0) AND MESSAGE(0) = 1
 * MANTISSA BITS = 10
 * EXPONENT BITS = 7
 DECIMAL PT LEFT(0) OR RIGHT(1) = 0

MANTISSA + EXPONENT + 1 SHOULD EQUAL N-BITS
 ERROR- ILLEGAL DETECTED DATA PARAMETER FOUND
 STOP BAD INPUT TO DETECT ROUTINE
 .280 OF SECOND EXECUTION TIME

N-BIT SPECIFICATIONS ARE:
 NUMBER OF BITS/WORD = 64
 TRUNCATE (0) OR ROUND (1) = 0
 PRECISION = 1
 OVERFLOW MESSAGE(0) AND MESSAGE(0) = 1
 * MANTISSA BITS = 53
 * EXPONENT BITS = 10
 DECIMAL PT LEFT(0) OR RIGHT(1) = 1

TOO MANY BITS SPECIFIED: MANTISSA IS 54 --- 53 WERE SPECIFIED
 ERROR- ILLEGAL DETECTED DATA PARAMETER FOUND
 STOP BAD INPUT TO DETECT ROUTINE
 .031 OF SECOND EXECUTION TIME

Figure 38 (Continued)


```

N-BIT SPECIFICATION: N=8
NUMBER OF BITS/WORD= 24

TRUNCATE (0) OF ROUND (1)=1
PRECISION =1

(1) OVERFLOW MESSAGE(1) AND MESSAGE(0)=17
# MANTISSA BITS =17
# EXPONENT BITS =6
DECIMAL AT LEFT OF RIGHT(1)=0

KEYS= 000000000000000000000000 12500000000000000000000000
15000000000000000000000000000000 00000000000000000000000000000000
00000000000000000000000000000000 00000000000000000000000000000000
100=119.71404875

N-BIT SPECIFICATION: N=8
NUMBER OF BITS/WORD= 48

TRUNCATE (0) OF ROUND (1)=0
PRECISION =1

(m) OVERFLOW MESSAGE(1) AND MESSAGE(0)=1
# MANTISSA BITS =41
# EXPONENT BITS =6
DECIMAL AT LEFT OF RIGHT(1)=0

KEYS= 000000000000000000000000 12500000000000000000000000
15000000000000000000000000000000 00000000000000000000000000000000
00000000000000000000000000000000 00000000000000000000000000000000
100=119.6930664974

```

Figure 38 (Continued)

Core And Execution Costs Of The N-bit Simulation Tool

Although not regarded originally as a limiting factor in the design of the n-bit simulation tool, the additional core and execution time requirements for n-bit simulating an algorithm may inhibit some users from readily using the tool. The resulting n-bit simulated version of a FORTRAN algorithm requires substantially more execution time and core space than the original algorithm.

For the algorithms shown in Figures 30, 34, and 37 the core space required for the n-bit simulated version was four to five times that used by the original algorithms. Moreover, the n-bit simulation subroutines added an extra 2000 words to the core size requirements. However, for the types of algorithms for which this tool is intended, the original algorithms usually do not require more than 5000 words. An expansion of these kinds of algorithms by a factor of four or five would probably not impact the user very much. Those algorithms that did exceed 5000 probably have much of their core space reserved for dimensioned arrays. The n-bit simulation tool would increase these algorithm's core space requirements to a lesser degree because the core expansion factor is dependent upon the number of arithmetic computations within an algorithm.

Execution time for an n-bit simulated algorithm may be a more significant problem. The programs in Figure 34 and 35 were modified so that the statements within the program were executed 500 times. The original algorithm required

.277 CPU seconds for execution while the n-bit simulated version required 2.4 CPU seconds, an increase of almost a factor of ten. Other tests found cases in which the CPU execution time was increased by more than 30 times.

This tremendous increase in execution time is due to the fact that for each arithmetic computation within a FORTRAN arithmetic statement, six to ten FORTRAN instructions must be executed to perform the necessary n-bit wordlength effects. In addition, there is the subroutine linkage time required to call n-bit simulation subroutine. This introduction of new subroutines into a FORTRAN algorithm causes the FORTRAN compiler to save all register values somewhere in storage before jumping to a subroutine, then restore those values after returning. An increase in execution time is not surprising considering these factors.

This additional execution time may greatly impact a user's readiness to take advantage of the n-bit simulation tool if the user's original algorithm required more than 120 seconds for execution. The n-bit simulation version of that algorithm could require anywhere from 1200 to 3600 seconds. Now the user would be dealing in terms of hours instead of minutes and that is only CPU time, not total run time, which would be even more. One method to lessen the execution time explosion factor is discussed in the following chapter.

Conclusions

In conclusion, the n-bit simulation tool was found to be an accurate and fairly versatile tool. Through the employment of various options, a user could simulate the effects of various wordlengths with rounding or truncation effects, double or triple precision arithmetic accuracy, various mantissa and exponent bit lengths, and simulate double wordlength storage variables (for floating point values). The tool does, however, cost the user additional core space and execution time. Nevertheless, it is a very useful tool for those applications where these limiting factors do not come into play to a substantial degree.

VI. Conclusions And Recommendations

This chapter provides the reader with conclusions and recommendations for the n-bit simulation tool.

Conclusions

The n-bit simulation tool provides the user with a means to evaluate the effects of varying wordlength upon the numerical accuracy of an algorithm. The results can then be evaluated against the performance specification to determine the wordlength and precision requirements.

The n-bit simulation tool requires that an algorithm being n-bit simulated is executed on the CDC 6600 or CYBER 74 computer systems and that the algorithm is expressed in a slightly restricted version of FORTRAN IV, as described in Appendix B.

Through the various user options available, the n-bit simulation tool can simulate the numerical effects for computations performed on a fairly wide range of computer types. By exercising various options, the user can specify a computer type which has various wordlengths, varying exponent and mantissa lengths (where the mantissa represents either a fractional or fixed point value), rounding or truncation effects for computations, single, double, or triple wordlength accumulators, and doubleword storage for wordlengths of up to 24-bits.

In addition, the user may specify different n-bit wordlength effects for different portions of an algorithm.

This feature would be especially useful in simulating the algorithms of an integrated system; where different portions of the total integrated system (such as the Kalman Filter, navigation, flight control, and weapon delivery systems) require differing amounts of precision.

Recommendations

Two areas for immediate/near future development are recommended:

1. Include an option for the n-bit simulation tool which would provide the user the exact results obtainable from an n-bit wordlength machine; specifically giving the values that overflow the exact representation which would result from an overflow of the n-bit wordlength machine.
2. Modify the n-bit simulation preprocessor to substitute in-line code, which would perform the n-bit wordlength effects, instead of substitution of subroutine calls. This change would improve the execution time of the n-bit simulated program.

Case One. The result of an overflow varies for different computer types. With the present n-bit simulation tool, maximum values are substituted for values that overflow. In the real world, it may be that the most significant bits are simply truncated off for fixed point addition overflows while multiplication overflows result in the correct mantissa and the maximum exponent possible.

An option to the n-bit simulator which would produce the exact results from an overflow, would enable the user to analyze the exact results obtainable from an algorithm that is processed on an n-bit wordlength machine. This new option would require an analysis of the ways in which different computer types overflow and incorporate simulation techniques into the n-bit simulation tool accordingly.

Case Two. Presently, n-bit simulation effects are performed by substitution of subroutine calls for each arithmetic operation of an algorithm. Upon execution of the n-bit simulation version of an algorithm, the subroutines referenced perform the n-bit wordlength effects. However, the additional time required to simulate these effects for all operations in an algorithm is a very high ratio (15 to 40 times) to the algorithms normal execution time.

It is anticipated that the introduction of in-line code to perform the n-bit wordlength effects would reduce the additional execution time by at least one-third. This proposed enhancement would allow the tailoring of the code, performed for the n-bit simulation task, to the user options in effect at the time. For example, the present system tests within each subroutine for the rounding option. The subroutine executes differently depending upon the option. By tailoring the in-line code to the user option, decision-making time would be saved. In addition, subroutine linkage time (saving and restoring of COMPASS registers) would be saved with the in-line code approach. The total amount of time saved

could be substantial depending on the number of arithmetic operations performed within an algorithm.

Although this enhancement would improve execution time, it would cost the user twice the amount of core space required by the present n-bit simulation tool. However, this additional core space may be more acceptable to the user than the additional execution time of the present system. The changes to the preprocessor required for this modification are described in detail in Appendix D.

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Appendix A

Approaches To Finding Extra COMPASS Registers

The solution to the problem of finding extra registers would be to find some standard value that is always stored in a known register or to devise some method of saving and restoring registers. Several possibilities were investigated:

- 1) The A0 register contents may be predictable. However, no consistencies were found and destroying its contents caused irregular results. A0 is used as the relative starting address in a block copy operation.
- 2) The CDC manuals stated it was a standard practice of CDC to use the B1/B7 registers to store the value 1. If this were the case, that register could be used as a spare register at anytime as long as it was restored to the value 1 when it was no longer needed. However, several cases of exception were found to this standard practice-so it was of no use.
- 3) The Exchange Jump instruction was known to save the contents of all registers, but its use is restricted to the operating system for switching between two central programs, leaving the first program in a useable state for later re-entry.
- 4) If CDC had a standard register saving/restoring convention similar to IBM's, the extra register problem

would be solved. However, CDC does not have one per se; instead before jumping to any external subroutine the register values which will be required later (within that program unit) are stored, then restored when needed.

- 5) Another theory was that since the A and X registers were associated one-for-one, the address in register A_n (where n is 1-5) could be used to recover the corresponding X register. This would enable the current value of the X register to be written over, since it could be restored later. Once one spare 60-bit register was found, all other registers could be saved using a series of register moves and write-outs to a known save area. However, the value of X register does not always correspond to the associated A register. A technique was devised which saved the contents of an X whenever it was changed. In that manner it could be restored later. However, there was still a need to have an associated pair of A6/X6 or A7/X6 registers to make this possible. All this would require a lot of analysis of the COMPASS code which was felt to be beyond the scope of this thesis.

Appendix B

User's Manual For The N-bit Simulation Tool

The n-bit simulation tool enables the user a means to evaluate the numerical effects that varying wordlengths have upon a FORTRAN-programmed algorithm. When applied to a FORTRAN-programmed algortihm, this tool will produce the near exact arithmetic accuracy as a computer of the specified n-bit wordlength. This n-bit simulation tool is useable on either the CDC 6600/CYBER 74 computer systems and requires that algorithms to be n-bit simulated be expressed in a somewhat restricted version of the FORTRAN IV (EXTENDED) programming language.

This User's Manual will describe the seven options available to the user in specifying different n-bit wordlengths, the using programmer's restrictions will be described, and the required control cards will be shown.

User Options

Seven options were considered necessary to provide the user with a versatile n-bit simulation tool which could be used to simulate the characteristics of various computer types. Different options could be applied to every routine within a program. These options are expressed by a call to the SETNBIT subroutine as the first executable statement of each routine (i.e. following the DIMENSION, COMMON, DATA statements, and so on). The general form of this subroutine

call, where each # represents a user option and KEY is a variable name which must be included as the eighth argument, is shown below:

```
CALL SETNBIT (#1,#2,#3,#4,#5,#6,#7,KEY)
```

The variable name KEY must be dimensioned to seven (e.g. DIMENSION KEY(7)) within each program routine. The KEY array is used to store values which are computed by the SETNBIT subroutine based upon the user options specified. These values are used to simulate the effects of n-bit wordlength for each arithmetic operation of an arithmetic statement.

Option #1 specifies the number of bits per wordlength (from 8 to 60 bits).

Option #2 specifies whether arithmetic operations will be performed with rounding or truncation effects. The value 1 would indicate rounding effects and 0 would indicate truncation effects.

Option #3 allows the user to specify whether arithmetic calculations are performed in single, double, or triple n-bit wordlength precision, with the results being stored as single precision quantities. This option essentially simulates the effects of a computer having a single, double, or triple wordlength accumulator but only single precision storage variables. Single is indicated by the value 1, double by 2, and triple precision by 3.

Option #4 gives the user the option to have overflow messages printed when they are detected (i.e. overflows of

the n-bit wordlength). The value 1 would indicate overflow messages will be printed when overflows occur and 0 would indicate the messages would not be printed. An overflow message would tell the user which routine and the line number within that routine where the n-bit wordlength overflow was detected. The line number would reference the line number from the original FORTRAN program (before the n-bit simulation modifications were made). In addition, the arithmetic operation being performed at the time is indicated.

Option #5 indicates the number of bits which are to be used for the mantissa of floating point values.

Option #6 indicates the number of bits which are to be used for the exponent of floating point values. The total number of bits specified for the mantissa and exponent must be one less than the total number of bits specified in option #1 (this other bit is used as the sign bit).

Option #7 specifies the position of the implied decimal point with respect to the mantissa, either to its left (making the mantissa a fractional representation) or to its right (making the mantissa a fixed point representation). The value 1 would position the decimal point on the right and 0 would position the decimal point on the left end of the mantissa.

An example of a typical SETNBIT subroutine call would be

```
CALL SETNBIT (16,1,2,0,9,6,0,KEY)
```

The options shown as arguments of SETNBIT specify a 16-bit wordlength with 9 bits in the mantissa and 6 bits in the exponent. The implied decimal point would be positioned to the left of the mantissa (making it a fractional representation). Rounding effects are specified for arithmetic calculations, double precision arithmetic accuracy is specified, and overflow messages are supposed to be suppressed.

Double precision storage can be simulated by adding the original wordlength to the mantissa specification (Option #5) and doubling the bit length specified in option #1. All other user options should be maintained as they were for the single precision storage word. For instance, for a 16-bit wordlength with 9 bits in the mantissa and 6 bits in the exponent, the double precision storage could be simulated by specifying a 32-bit wordlength, a 25 bit mantissa, and a 6 bit exponent.

Double precision arithmetic accuracy should be distinguished from double precision storage. With double precision arithmetic accuracy (which can be specified under option #3) a FORTRAN expressed arithmetic statement will be computed with double precision accuracy, but after the final computation of the statement is performed the result stored will contain only single precision accuracy. In contrast, double precision storage would have double precision accuracy, in the computation as well as double precision storage of the result. Since the n-bit simulation tool does not presently have a separate option to designate double precision

storage, the double precision can be accomplished by doubling the original wordlength and increasing the mantissa by one wordlength (discussed previously).

User Programming Restrictions

The user programming restrictions which must be adhered to for programs to be n-bit simulated are listed in this appendix. Restrictions are listed in the order in which they are most likely to impact the user. Each restriction will be briefly justified. Other things a user should be aware of are listed in a section following the restrictions.

1. Every routine must have the variable KEY dimensioned to seven (DIMENSION KEY (7)). Also, the first executable statement within each routine must be CALL SETNBIT (seven options and KEY), where the user specifies the seven options to be applied to that particular routine. These additions are required so each routine can perform the n-bit simulation effects.

2. Single assignment statements are required for all values coming from a source of different bit length specifications. Thus all values read in from cards or values passed as parameters from a routine of different n-bit specifications must be set equal to themselves, in order for them to take on their n-bit significance (since all n-bit simulated subroutines assume the values being operated upon have their n-bit significance already). Therefore, if the variable A was read into an n-bit routine then A should be

set equal to itself ($A = A$) before it is used anywhere else in the program.

3. The entire program must be free of syntax errors, before it is run through the n-bit simulation preprocessor. Debugging errors would be easier with the original program and the preprocessor has not been verified to be able to handle all syntax errors.

4. TAPE1 must be an output unit for each n-bit simulated program (overflow messages are printed to the unit number).

5. CONTINUE statements must be used as the last statement for DO Loops due to the way labels are handled by the preprocessor for arithmetic statements.

6. Variable data types must use default data types (Reals (A-H) and (O-Z), Integers (I-N)). Declaration of data types using INTEGER and REAL statements are not allowed, nor are IMPLICIT INTEGER statements allowed. The preprocessor selects handling routines based on default data types.

7. All arrays must be declared by DIMENSION statements (not by COMMON, INTEGER, or REAL declarations). This is so the preprocessor can detect functions from arrays (it finds all arrays in the DIMENSION statement).

8. Logical, double precision, and alphanumeric data handling is not allowed. The n-bit simulated program expects all numeric values in assignment and arithmetic statements.

9. These are certain reserved subroutine names (shown in Figure B-1) and reserved variable names which must not be used by user (except for KEY & SETNBIT). These names are used by the preprocessor.

10. No statement functions are allowed.

11. User should not use shift operations to perform multiplications or divisions (by powers of two). These operations would not be detected as arithmetic operations and would not be n-bit simulated properly.

12. Data statements should not be used for real values or values that might exceed an n-bit simulated word-length. These Data statements are not analyzed by the preprocessor.

13. Reserved labels are 79110 - 79199.

14. Only one statement per line maximum (i.e. no A=B=C=0).

RRADD	R1ADD	R2ADD	IIADD	IASGN	ITTTTAA	RTTTTAA
PRMNS	R1MNS	R2MNS	IIMNS	RASGN	ITTTTAB	RTTTTAB
FRMPY	R1MPY	R2MPY	IIMPY	KEY	:	:
RRDVD	R1DVD	R2DVD	IIDVD	SETNBIT	:	:
RREXP	R1EXP	R2EXP	IIEXP		ITTTTAZ	RTTTTAZ

Figure B-1 Reserved Variable Names For
N-bit Simulated Programs

Other User Considerations:

1. Intrinsic functions (such as SIN,COS,ATAN), other FORTRAN library routines, and all other routines not read into the n-bit simulation preprocessor are not n-bit simu-

lated completely since they are coded in COMPASS and the pre-processor can only handle FORTRAN programs. However, the arguments are n-bit simulated and the resulting value from the function is n-bit simulated.

2. Array arguments are not n-bit simulated .

3. Only assignment and arithmetic statements are affected by the n-bit simulation tool. Expressions occurring elsewhere will not be affected (such as inside IF statements). Also data initialized within DATA statements will not be affected.

4. For exponentiation, n-bit simulation affects are only applied to the result of the exponentiation operation, so the results do not reflect on an n-bit wordlength machine exactly in some cases.

5. Equivalence statements are not recommended, but can be used if the programmer fully understands the impact n-bit simulation may have upon the equivalenced values.

6. Any user should be aware that any arithmetic error which would result in abnormal termination in a CDC FORTRAN program will terminate the n-bit simulated program also. Even though no abnormal termination resulted from execution in a normal CDC execution, the effects of n-bit wordlength may cause a value to be generated which could terminate it (for instance, .000001 might be truncated to 0. for a 16-bit word and would result in underflow termination when used as a divisor).

The pointers shown on the left of Figure B-2 indicate those statements which would not have appeared in the program if the program was not to be n-bit simulated.

Control Cards

The following card sequence could be used to n-bit simulate a FORTRAN programmed algorithm:

```
FTN,L=DUMMY.  
LGO,,,NBIT.  
REWIND,NBIT.  
FTN,L=DUMMY2,B=SUBS.  
REWIND,SUBS.  
EDITLIB.  
LIBRARY(NSUBS).  
FTN,I=NBIT,B=NBITGO,L=DUMMY3.  
REWIND,NBITGO.  
NBITGO.  
7/8/9 Card  
*** Preprocessor source program ***  
7/8/9 Card  
*** USER'S SOURCE PROGRAM ***  
7/8/9 Card  
*** N-bit Simulation Subroutines (Source) ***  
7/8/9 Card  
LIBRARY(NSUBS,NEW)  
ADD(*,SUBS)  
FINISH.  
ENDRUN.  
7/8/9 Card  
6/7/8/9 End-of-Job Card
```

To n-bit simulate the USER'S SOURCE PROGRAM, the programming restrictions previously described must be strictly followed. Then, the USER'S SOURCE PROGRAM is read by the n-bit simulation preprocessor program, which transforms arithmetic statements of the original program into subroutine calls, which are designed to accomplish the n-bit simulation word-length effects. The preprocessor outputs the n-bit simulated version of the USER'S SOURCE PROGRAM. Afterwhich the n-bit


```

→ PROGRAM PREPROCESS(INPUT,OUTPUT,TAPE1=OUTPUT)      000100
  EXPLICIT INTEGER (A-Z)                             000110
  C IMPLICIT INTEGER IS NOT ALLOWED IN THE N-BIT SIMULATOR 000120
  C NOR ARE INTEGER OR REAL VARIABLE DECLARATIONS, VARIABLE 000130
  C TYPE SHOULD USE DEFAULT VALUES I-N FOR INTEGER,AND 000140
  C A-H,I-J FOR REALS 000150
→ DIMENSION C(20),I(15,2),K(1:KEY(7)) 000160
COMMON A,B,DDDD 000170
INTEGER A2,A4 000180
EQUIVALENCE (C5,D5),(I1,D6) 000190
DATA C/20*1.17,11*10/ 000200
C THIS IS ONLY A TEST PROGRAM TO EXAMINE THE PREPROCESSORS 000210
C HANDLING OF THE VARIOUS STATEMENT TYPES 000220
C A "CALL SETUP" MUST BE THE FIRST EXECUTABLE STATEMENT WITHIN 000230
C EACH ROUTINE AND KEY(1) MUST BE DIMENSIONED FOR EACH ROUTINE 000240
C IN THE N-BIT SIMULATED PROGRAM 000250
→ CALL SETUP(20,1,1,1,17,6,C,KEY) 000260
READ *,C5,I1 000270
C ALL EXTERNAL VALUES COMING INTO A N-BIT SIMULATED ROUTINE 000280
C MUST BE SET EQUAL TO THEMSELVES-IN ORDER TO GIVE EACH ITS N-BIT 000290
C SIGNIFICANCE 000300
→ C=C5 000310
I=I1 000320
C EVERY EXTERNAL VARIABLE MUST BE ASSIGNED TO ITSELF TO 000330
C N-BIT SIMULATE ITS ACCURACY 000340
READ (1,7) (K(L),L=1,1) 000350
FORMAT(15F3.7) 000360
C THE ARRAY K (EXTERNAL VALUES) MUST BE SET EQUAL TO ITSELF TO 000370
C GIVE IT ITS N-BIT WORDLENGTH SIGNIFICANCE 000380
→ DO 4 L=1,15 000390
K(L)=K(L) 000400
CONTINUE 000410
I1=(23+7*(PRV-11.0)) 000420
DO 5 M=1,15 000430
C(4)=K(4)*2 000440
CONTINUE 000450
PRINT *, " INPUT " 000460
PRINT *,K 000470
A1=K(1)+7324/K(2) 000480
CALL SUB1(25,6,A1) 000490
A=1+2/3+I+C(12) 000500
CALL MANIPUL (A,B,X,T),RETURNS(13) 000510
DEIND=1 000520
ASSIGN 13 TO ILABEL 000530
DO 7 L=1,15 000540
C(L)=C(L,1)*2+1.37 000550
CONTINUE 000560
GO TO (17,15),AILABEL 000570
IF(I1.EQ.0) GO TO 9 000580
WRITE (1,19) A1,(C(NN),NN=1,20,2) 000590
FORMAT(2Y,=A10) 000600
GO TO 13 000610
PRINT *, " STARTED LAST PART" 000620
13 A1=C(2)*C(19)+K(1)**(-1.7) 000630
15 A2=C(1)*(A1+K(1)) 000640
(A1-B)+1.000 000650
STOP " PREPROCESS TEST COMPLETE" 000660
END 000670
SUBROUTINE SUB1(I1,C,A1) 000680
COMMON A,B,DDDD 000690
DIMENSION C(1),KEY(7) 000700
CALL SETUP(20,1,1,1,17,6,C,KEY) 000710

```

Figure B-2 Sample Program

simulation subroutines which perform the n-bit wordlength effects are compiled and libraried. Then, the n-bit simulated user program is compiled and executed. If the program uses additional input/output files, these files would be put in their proper position along with the NBITGO card.

If the object forms of the preprocessor and n-bit simulation subroutines are used instead of the source versions, the control card sequence would be:

```
COPYBR,,SUBS,1.
INPUT,,,NBIT.
REWIND,SUBS.
EDITLIB.
LIBRARY(NSUBS).
FTN,I=NBIT,B=NBITGO,L=DUMMY3.
REWIND,NBITGO.
NBITGO.
7/8/9 Card
*** N-bit Simulation Subroutines (Object) ***
7/8/9 Card
*** Preprocessor object deck ***
7/8/9 Card
*** USER'S SOURCE PROGRAM ***
7/8/9 Card
LIBRARY(NSUBS,NEW)
ADD(*,SUBS)
FINISH.
ENDRUN.
7/8/9 Card
6/7/8/9 End -of-Job Card
```

If the user wishes to see the n-bit simulated version of the USER'S SOURCE PROGRAM, the L=DUMMY3 parameter on the last FTN card should be removed. Normally this print would probably be of little interest to the user.

The user should be aware that the n-bit simulation of a program will probably increase the core requirements by four to five times and the execution time may increase up to 40 times the program's normal execution time (without n-bit simulation effects).

With the present n-bit simulation tool, values which overflow are replaced by the maximum significant value that the n-bit wordlength can represent. For example, if the resulting value of a computation exceeded the maximum negative number representable by the n-bit word, the value would be replaced by that maximum negative number. In addition, an overflow message would be printed if the user option was on. Similarly, an overflow of the positive maximum would result in the replacement of the value that overflowed with the positive maximum number. Underflow results in the replacement of zero for the value that underflowed. Some results of the n-bit simulation tool are shown in Chapter 5 of the main text.

Appendix C

Maintenance Manual For The N-bit Simulation Tool

The n-bit simulation tool consists of two parts: the preprocessor program and the n-bit simulation subroutines. This manual includes a description of the overall n-bit simulation tool system, general descriptions of the preprocessor and n-bit simulation subroutines, structure charts, interface charts, and module definitions for the preprocessor. In addition, program listings of the preprocessor and the n-bit simulation subroutines are included.

Overall N-bit Simulation Tool

The preprocessor reads a FORTRAN program (to be n-bit simulated) and converts each arithmetic statement into one or more subroutine calls. A subroutine call is generated for each arithmetic operation of an arithmetic statement. The subroutine calls, generated by the preprocessor, reference the n-bit simulation subroutines. When the n-bit simulated version of the original algorithm is executed, the n-bit simulation subroutines perform the n-bit wordlength effects upon each arithmetic statement operation.

The n-bit wordlength effects are specified by the user within each routine of the original FORTRAN program. The results from the execution of the n-bit simulated version are nearly the exact result obtainable from a computer of n-bit wordlength.

Preprocessor

The preprocessor basically reads and syntactically analyzes each statement of the user's program. Each arithmetic statement is transformed into a Polish Notation string, then decomposed one operation at a time. For each operation, a subroutine call is generated. The subroutine referenced will perform the arithmetic computation indicated by the subroutine name. Each combination of variable data types is also handled by a separate subroutine. The operands which the operation is being performed upon may be either integer or real. A different subroutine reference is generated for each possible combination of operand data types and arithmetic operation.

The module structure of the preprocessor is shown in Figure C-1. The module definitions follow, along with a chart describing the preprocessor's module interfaces. The program listing of the preprocessor follow Figure C-1.

Module Definitions And Module Interfaces For The Preprocessor

PREPROCESSOR-EXEC is the main routine of the preprocessor program. It calls the afferent and efferent branches (EVALUATE-STATEMENT and SET-UP-CALLS) which carry out the translation of a FORTRAN algorithm to the n-bit simulation equivalent algorithm in addition to the INITIAL-READ module.

INITIAL-PEAD does the initial read of the FORTRAN algorithm. It reads the contents of the first program card into a temporary array which is used by the GET-STATEMENT module, which functions on a 'look-ahead' basis.

EVALUATE-STATEMENT supervises the evaluation of (legal) arithmetic statements into their Polish Notation equivalent representation.

GET-OBJECT supervises the check of a statement for a legal object, where an arithmetic statement structure consists of object = arithmetic expression .

GET-STATEMENT reads input program's cards, passing a 'statement' a time to its superordinates, GET-OBJECT. A statement may stretch over several cards through the use of continuation cards, so a look-ahead read is performed.

FIND-OBJECT determines if a statement has a legal object.

NON-EXPRESSION-HANDLER calls for the print of a non-expression statement and determines if the non-expression is a key statement type, for dimensioning arrays or the beginning of a new routine.

NEW-ROUTINE-HANDLER extracts the program unit's name and initializes the line counter for that new unit to one. (This information is used for each arithmetic function subroutine call built by the efferent branch). Also it clears the list of variable names declared as arrays for the previous program unit.

DIMENSION-STATEMENT-HANDLER stores all variable names which have been declared arrays.

PRINT-STATEMENT prints the statement which may span several cards to the n-bit simulated program output file.

EVALUATE-EXPRESSION is a recursive module which determines (through a syntax analysis) if the expression is a legal arithmetic expression, and converts the statement into its Polish Notation equivalent representation. It is recursive because parentheses may be used to enclose another complete expression and parentheses may occur throughout an expression.

ANALYZE-OPERAND determines if the syntax of a string of characteristics constitutes a legal variable name (a variable may also be an array or function reference).

ANALYZE-NUMBER determines if the syntax of a string of characters constitutes a legal number.

ANALYZE-POSSIBLE-ARRAY determines if a variable name has been declared an array.

CLASSIFY-CHARACTER types input character as a letter, digit, operator, parenthesis, decimal point, or comma.

STORE-CHARACTER-IN-POLISH-STRING stores a variable name or operator in Polish string according to Reverse Polish Notation and the operator precedence.

PUT-VARIABLE-NAME stores a string of characters in the NAME array and passes back to superordinate a pointer to the position of the string in the NAME array.

PRELIMINARY-OUTPUT outputs a COMMENT statement containing the input statement and outputs preliminary output for labels and IF statements.

SET-UP-CALLS scans Polish string left to right and calls a handling routine for each operator found.

BUILD-SUB-CALLS assembles and merges parts of a subroutine call.

PRINT-SUB-CALL prints 80 character card images to the n-bit simulated program output file.

BUILD-FUNC-ARGUMENTS merges function arguments together, generating calls for single arguments if necessary.

FINISH-FUNCTION merges the function name with its argument list (built by BUILD-FUNC-ARGUMENTS) and generates a single assignment call for the function.

SET-UP-SINGLE-ASSIGNMENT sets up parts for a single assignment.

HANDLE-UNARY places a unary sign in front of the variable name through a call to the PUT-VARIABLE-NAME module.

OUTPUT-SINGLE-ASSIGNMENT merges parts required for a single assignment function call.

GET-VARIABLE-NAME gets a string of characters from the name array according to the pointer given and passes to the caller the string, its data type, and its length.

MERGE merges strings of characters together.

Module Interfaces

Interface # (Reference Fig. 29)	In	Out
1,2,31	-	-
3	STMT, POSITION	-
4,5	-	STMT, POSITION
6	-	STMT
7,11,20,30	-	STMT
8	-	STMT, POSITION
9,15,22,25,28, 60	-	CHARACTER
10,16,24,27,59	POSITION	STMT
12	-	STMT
13,14	POSITION	STMT, POSITION, TYPE
17	POSITION	STMT, POSITION, TYPE
18	-	STMT, POSITION
19	-	STMT, POSITION
21	POSITION	STMT, POSITION, NAME, LENGTH, NMPOINTER, MODE
23,26,29,37,42, 51,54	NAMEPOSITION	NAME, LENGTH, MODE
32	TEMPNAMES	POLISH, POLPOSITION, TEMPNAMES, FINALFLG, TRLR, TRLRLEN
33	TEMPNAMES	POLISH, POLPOSITION, TEMPNAMES, FINALFLG, TRLR, TRLPLEN
34	-	POLISH, POLPOSITION, TEMPNAMES, TRLR, TRLRLEN
35	-	POLISH, POLPOSITION, TRLR, TRLRLEN
36	-	POLISH, POLPOSITION
38,44,50,61	NAME, MODE, LENGTH	NAMEPOSITION
39,44,49,55,58	ARRAY1, LEN1	ARRAY1, LEN1, NAME2, LEN2
40,45,48,53,56	OPENDPOSITION	POLISH, POLPOSITION
41,57	-	OUT, LENGTH
46,47,52	-	NAME, NAMELEN, MODE, OBJECT, OBJLEN, TRLR, TRLRLEN

N-bit Simulation Preprocessor
Program Listing


```

PROGRAM PREPPDS(INPUT,OUTPUT,TAPE1=4002/80,TAPE2=INPUT)      000100
C                                                                000102
C      N-RIT SIMULATION PREPROCESSOR PROGRAM WILL CONVERT AN  000103
C      INPUT FORTRAN PROGRAM INTO ITS EQUIVALENT N-RIT SIMULATED 000104
C      VERSION WHICH USES SUBROUTINE CALLS TO PERFORM THE N-RIT  000105
C      WORDLENGTH EFFECTS UPON THE INPUT PROGRAM'S ARITHMETIC STMTS 000106
C                                                                000109
C      IMPLICIT INTEGER (A-Z)                                  000110
C      TO THE INITIAL READ OF THE INPUT PROGRAM                000118
CALL INITPD, RETURNS(100)                                     000120
C      BUILD POLISH STRING FOR EACH ARITHMETIC STATEMENT       000123
C      ALT RETURN IS USED WHEN AN EOF HAS BEEN ENCOUNTERED    000124
3 CALL EVLSTMT, RETURNS(100)                                   000130
C      OUTPUT N-RIT SIMULATED CODE FOR EACH ARITHMETIC STATEMENT 000134
C      WHICH IS NOW EXPRESSED IN POLISH NOTATION               000135
CALL SETUP                                                    000140
GO TO 3                                                       000150
C      EOF ENCOUNTERED                                         000157
100 STOP " EVLSTMT COMPLETED"                                000160
END                                                            000170
SUBROUTINE EVLSTMT, RETURNS(N1)                               000180
IMPLICIT INTEGER(A-Z)                                         000190
EVLSTMT SUPERVISES THE EVALUATION OF LEGAL ARITHMETIC        000192
STATEMENTS INTO THEIR POLISH NOTATION EQUIVALENT             000193
REPRESENTATION                                                 000194
DIMENSION STMT(4003)                                          000200
CALL GETOBJ(STMT,0), RETURNS(100)                             000210
CALL EVLEXP(STMT,0), RETURNS(5)                               000220
C      PRELIMINARY OUTPUT FOR THE ARITHMETIC STATEMENT FOR SUCH 000224
C      THINGS AS LABELS AND IF STMTS IN ADDITION TO PRINTING A  000225
C      A COMMENT STATEMENT CONTAINING THE ORIGINAL STATEMENT   000226
CALL PRELSTMT(STMT)                                           000230
RETURN                                                         000240
101 RETURN N1                                                 000250
END                                                            000270
SUBROUTINE EVLEXP(STMT,0), RETURNS(N2)                       000280
IMPLICIT INTEGER(A-Z)                                         000284
C                                                                000285
C      02 ALTERNATE RETURN USED AFTER THE DETECTION AND        000286
C      AND HANDLING OF AN ILLEGAL ARITH STATEMENT              000287
C                                                                000290
C      EVALUATE-EXPRESSION IS A RECURSIVE MODULE WHICH         000291
C      DETERMINES(THROUGH SYNTAX ANALYSIS) IF THE EXPRESSION   000292
C      IS A LEGAL ARITHMETIC EXPRESSION, AND CONVERTS THE STATE- 000293
C      MENT INTO ITS POLISH NOTATION EQUIVALENT REPRESENTATION  000294
C      IT IS RECURSIVE BECAUSE PARENTHESES MAY BE USED TO     000295
C      ENCLOSE ANOTHER COMPLETE EXPRESSION, WHICH MUST UNDERGO 000296
C      EXACTLY THE SAME SYNTAX ANALYSIS. PARENTHESES MAY OCCUR  000297
C      THROUGHOUT AN EXPRESSION.                                000298
C                                                                000299
C      DIMENSION STMT(1)                                         000300
DATA UNARY/101/,TWO/102/,THREE/103/,PLUS/104/,MINUS/105/,FUNC/106/ 000302
C      FUNCTION IS A FLAG INDICATING WHETHER A FUNCTION IS BEING 000303
C      EVALUATED                                                 000304
C      EVALUATED                                                 000305
C      INDICATES IT IS VALID TO END EXPRESSION                 000306
1  GO TO 101                                                    000307
C      CALL CLASSIFY(UNARY,0,TYPE), RETURNS(0,500)             000310
C      --- THEN FOR EACH ( STMT(I), I=1,N)                      000311
C      IF(TYPE,0,1) GO TO 101                                    000312
C      ONLY PLUS, MINUS, AND OTHER, SINCE THEY HAVE NO LEFT HANDING 000313
C      ONLY UNARY, SINCE IT IS USED ON TO THE POLISH STRING     000314

```

IF(STMT(P).EQ.PLUS)	GO TO 40	000420
IF(STMT(P).NE.MINUS)	GO TO 600	000430
CALL STPOL(UNARY)		000440
C LOOKING FOR AN OPERAND		000450
40 CALL CLASIFY(STMT,P,TYPE),RETURNS(500,600)		000460
C -- LOOKING FOR THE FIRST VALID CHARACTER OF A STRING		000470
C LETTER # OPERATOR () , . = OTHER		000480
45 GO TO (50,70,100,110,600,600,70,600,600),TYPE		000490
C POSSIBLE OPERAND		000500
50 OK=0		000510
CALL OPND(STMT,P,TYPE),RETURNS(500,600,55)		000520
C ABOVE ALT RETURNS SOS,ERROR,FUNCTION POSSIBILITY		000530
C THE FOLLOWING CHECK FOR LEGAL FOLLOWERS OF OPERANDS		000534
C CHECK FOR OPERATOR		000535
IF(TYPE.EQ.?) GO TO 90		000540
C --- CHECK FOR RIGHT PAREN ---		000550
IF(TYPE.EQ.5) GO TO 130		000560
C --CHECK FOR COMMA--		000570
IF(TYPE.EQ.6) GO TO 150		000580
C OTHERWISE ILLEGAL FOLLOWER OF AN OPERAND		000594
GO TO 400		000590
C --- GOT A POSSIBLE FUNCTION --- SET FUNCTION FLAG		000600
55 CALL STPOL(FUNC)		000610
NPAREN=NPAREN+1		000620
FNDFLG=1		000630
OK=1		000640
C RECURSE TO THE BEGINNING OF ROUTINE TO EVALUATE EXPRESSIONS		000645
C USED AS FUNCTION ARGUMENTS		000646
GO TO 1		000650
C GOT A POSSIBLE #		000660
70 OK=0		000670
CALL NUM(STMT,P,TYPE),RETURNS(500,600)		000680
IF(TYPE.EQ.?) GO TO 90		000690
IF(TYPE.EQ.5) GO TO 170		000700
IF(TYPE.EQ.6) GO TO 150		000710
C ILLEGAL FOLLOWER OF #		000720
GO TO 400		000730
C -- GOT AN OPERATOR		000740
90 OK=1		000750
C CHECK FOR POSSIBLE EXPONENTIATION		000755
IF(STMT(P).EQ.TIMES) GO TO 93		000760
C STORE OPERATOR IN POLISH STRING		000764
91 CALL STPOL(STMT(P))		000770
C --- GO LOOK FOR AN OPERAND		000770
GO TO 40		000800
C --- CHECK FOR EXPONENTIATION ---		000810
92 LP=P		000820
CALL CLASIFY(STMT,LP,TYPE),RETURNS(600,600)		000830
IF(STMT(LP).NE.TIMES) GO TO 91		000840
C -- SET LP UP --		000850
P=LP		000860
C -- STORE EXPONENTIATION IN POLISH STRING --		000870
CALL STPOL(EXP)		000880
GO TO 40		000890
C -- GOT LEFT PAREN -- START OF A NEW EXPRESSION --		000900
110 NPAREN=NPAREN-1		000910
CALL STPOL(STMT(P))		000920
C --- USED AN OBJECT SO SET OK TO 1		000930
OK=1		000940
GO TO 1		000950
C -- GOT A RIGHT PAREN --		000960
170 NPAREN=NPAREN-1		000970
CALL STPOL(STMT(P))		000980
C --- RELO CHECKS TO SEE IF IT HAS A LEGAL PARENTHESIZED EXPRESSION		000990
IF(OK.NE.0) GO TO 40		001000
C CALL NEXT SUBROUTINE WITH P1 TO SEE IF ITS A LEGAL FOLLOWER OF		001010

C	RIGHT PARENTHESIS	001024
	CALL CLASSIFY (STMT, P, TYPE), RETURNS (500, 500)	001030
	IF (TYPE, EQ, 3) GO TO 90	001040
	IF (TYPE, EQ, 5) GO TO 130	001050
	IF (TYPE, EQ, 6) GO TO 150	001060
	GO TO 600	001070
C	GET A COMMA	001090
150	OK=1	001100
C	SHOULD NOT HAVE A COMMA UNLESS IT IS INSIDE A FUNCTION	001110
	IF (FNCLG, EQ, 0) GO TO 600	001120
	CALL STPOL (STMT (P))	001130
C	NEED AN OPERAND	001140
	GO TO 1	001150
C	CHECK FOR END	001160
C	CHECK FOR ENDING PROPERLY	001170
500	IF (OK, NE, 0, 02, NPAREN, NE, 0) GO TO 600	001180
C	NOW CLEAR STACK IN POLISH STORE ROUTINE	001190
	CALL STPOL (STMT (P))	001200
	RETURN	001210
C	ILLEGAL ARITHMETIC EXPRESSION, SO PRINT IT	001216
600	CALL PRNT (STMT)	001220
	RETURN 42	001230
	END	001240
	SUBROUTINE OPERND (STMT, P, TYPE), RETURNS (N1, N2, N3)	001250
	IMPLICIT INTEGER (A-Z)	001256
C		001260
C	OPERND ANALYZES A POSSIBLE OPERAND TO DETERMINE IF THE	001261
C	SYNTAX OF A STRING OF CHARACTERS CONSTITUTES A LEGAL	001262
C	VARIABLE NAME (A VARIABLE MAY ALSO BE AN ARRAY OR	001263
C	FUNCTION REFERENCE).	001264
C	OPERND CHECKS TO SEE IF CHAR SEQUENCE IS A VALID OPERAND	001270
C	IT LIMITS THE SEQUENCE TO 7 CHAR MAX AND WILL DETERMINE THE MODE	001280
C		001289
	DIMENSION STMT(1)	001290
	DATA I/121/, N/124/, 75EO/120/, BLANKS/104 /, FJND/122/	001300
	DATA MAXLEN/7/	001310
C	TESTING TO SEE IF FIRST CHARACTER OF OPERAND IS A LETTER	001320
	IF (TYPE, NE, 1) GO TO 400	001330
C	MODE=1 (REALS) =2 (INTEGERS)	001340
	MODE=1	001350
C	CHECK FOR INTEGER VARIABLE (BASED ON LEADING CHARACTER)	001352
	IF (STMT (P), GE, 1, AND, STMT (P), LE, 9) MODE=2	001360
	MODE=1	001370
	NAME=BLANKS	001380
C	CNT COUNTS # CHARACTERS IN VARIABLE NAME	001392
C	BELOW COMPUTED # WITH CHARACTER BEING STORED MUST BE	001393
C	SHIFTED TO PLACE IN PROPER POSITION WITH RESPECT TO	001394
C	THE OTHER CHARACTERS ALREADY STORED FROM THE VARIABLE NAME	001395
	LEFT=(9-CNT)*3	001396
C	CHANGE TO SAVE BLANKS SOMETIME*****	001400
	NAME=SHIFT (STMT (P), LEFT), OR, (MASK (CNT*3), AND, NAME)	001410
	CNT=CNT+1	001420
	CALL CLASSIFY (STMT, P, TYPE), RETURNS (500, 500)	001430
C	TYPE > 2 INDICATES NEXT CHARACTER IS NOT LETTER OR DIGIT	001431
	IF (TYPE, GT, 2) GO TO 35	001440
C	MAX LENGTH OF A VARIABLE NAME IS SEVEN	001445
	IF (CNT, EQ, 7) GO TO 400	001450
	GO TO 5	001460
C	GET POSSIBLE ARRAY/FUNCTION (TYPE = 1 INDICATES LEFT PAREN)	001470
IF	IF (TYPE, EQ, 1) GO TO 45	001480
C	END OF VARIABLE- GO STORE IT AWAY	001490
	CALL PUTNAME (NAME, CNT, MODE, PTR)	001500
	CALL STPOL (PTR)	001510
	GO TO 23	001520
C	RETURN (0) INDICATES NAME WAS NOT DECLARED AN ARRAY-	001523
C	GO TO 10 POSSIBLE FUNCTION	001524

C	THEN CLASSIFY NEXT CHARACTER FOLLOWING ARRAY	001525
C	ALT RETURN (48) INDICATES EOS, WHICH IS LEGAL END	001526
C	TO AN EXPRESSION	001527
45	CALL PR3APPY (STMT, P, NAME, CNT-1, NPTR, MODE), RETURNS (43)	001530
	CALL CLASSIFY (STMT, P, TYPE), RETURNS (500, 400)	001540
	GO TO 53	001550
C	GOT A FUNCTION	001560
48	CALL PUTNAME (NAME, CNT, MODE, PTR)	001570
	CALL STPOL (PTR)	001580
C	RETURN N3 INDICATES OPERAND WAS A FUNCTION	001590
	RETURN N3	001600
53	RETURN	001610
C	RETURN N2 -ILLEGAL CHARACTER WAS ENCOUNTERED (THUS AN	001612
C	AN ILLEGAL OPERAND)	001614
400	RETURN N2	001620
500	CALL PUTNAME (NAME, CNT, MODE, PTR)	001630
C	EOS WAS ENCOUNTERED SO STORE VARIABLE NAME	001632
	CALL STPOL (PTR)	001640
530	RETURN N1	001650
	END	001660
	SUBROUTINE NUM (STMT, P, TYPE), RETURNS (N1, N2)	001670
	IMPLICIT INTEGER (A-Z)	001680
C	ALT RETURN N1=EOS, N2=ILLEGAL #	001682
C	NUM DETERMINES IF THE SYNTAX OF A STRING OF CHARACTERS	001683
C	CONSTITUTES A LEGAL NUMBER	001684
C	CHECK FOR LEGAL INTEGER/REAL #S	001690
	IMENSION STMT(1), NUM(2)	001700
	DATA DEPT/1P./, BLANKS/10H	001710
C	DEFAULT IS INTEGER (MODE=2)	001720
	MODE=2	001730
C	IF 1ST CHAP IS A DECIMAL PT, SET MODE TO 1 (FOR REALS)	001772
	IF (STMT(1).EQ.1P.) MODE=1	001740
	IF (TYPE.NE.2.AND.TYPE.NE.7) GO TO 600	001750
	N1=1	001760
C	*****MAY WANT TO CHECK FIRST CHAP IN FOR A NUMBER*****	001770
	CNT=0	001780
C	SHIFT CHARACTER INTO STORAGE POSITION	001795
10	LEFT=(CNT-1)*5	001790
	NUM(1)=SHIFT (STMT(1), LEFT).OR.(MASK(CNT*5).AND.NUM(1))	001800
	CNT=CNT+1	001810
	CALL CLASSIFY (STMT, P, TYPE), RETURNS (500, 500)	001820
	IF (TYPE.EQ.2) GO TO 20	001830
C	---CHECK FOR DECIMAL POINT-----	001840
	IF (TYPE.EQ.1) GO TO 15	001850
C	--- CHECK FOR OPERATOR ---	001860
	IF (TYPE.EQ.7) GO TO 57	001870
C	--- ALSO A) OR , IS A LEGAL END TO A NUMBER -----	001880
	IF (TYPE.EQ.3.OR.TYPE.EQ.5) GO TO 53	001890
	GO TO 600	001900
15	IF (MODE.EQ.1) GO TO 600	001910
	MODE=1	001920
20	IF (CNT.LE.10) GO TO 10	001930
C	INCREMENT TO NEXT WORD, IF 1ST WORD IS FULL	001935
	N1=N1+1	001940
	CNT=0	001950
	GO TO 10	001960
57	CNT=(N-1)*10+CNT	001970
	CALL PUTNAME (NAME, CNT, MODE, PTR)	001980
	CALL STPOL (PTR)	002000
	RETURN	002010
C	COMPUTE LENGTH OF THE # THEN RETURN FOR EOS CONDITION	002015
60	CNT=(N-1)*10+CNT	002020
	CALL PUTNAME (NAME, CNT, MODE, PTR)	002030
	CALL STPOL (PTR)	002040
	RETURN N1	002050
C	--- ILLEGAL NUMBER RETURN ---	002060

600	RETURN N2	002070
	END	002080
	SUBROUTINE PROCARRY (STMT,P,NAME,NLEN,NPTR,TYPE), RETURNS (N1)	002090
C		002099
C	ALT RETURN N1 INDICATES A THE VARIABLE WAS NOT DIMENSIONED-	002100
C	--- PROBABLE FUNCTION	002101
C	PROCARRY DETERMINES IF A VARIABLE NAME HAS BEEN	002105
C	DECLARED IN AN ARRAY, IF SO IT STORES THE ARRAY NAME AND ITS	002107
C	ARGUMENT LIST AWAY IN THE NAME ARRAY (VIA PUTNAME)	002108
C		002109
	IMPLICIT INTEGER (A-Z)	002110
	COMMON/ARPOYS/ARRAY(40), ARRYMAX	002120
	DIMENSION STMT(1), ANAME(5)	002130
	DATA LPAREN/12(1,2), RPAREN/12(1,2), BLNK/12(1,2), EOS/32(1,2), MAXLEN/5/	002140
	IF (ARRYMAX.EQ.0) GO TO 6	002150
	DO 5 I=1, ARRYMAX	002160
	IF (NAME.EQ.ARRAY(I)) GO TO 14	002170
5	CONTINUE	002180
C	GO TO A FUNCTION PROBABLY	002190
5	RETURN N1	002200
C	STARTING TO BUILD ARRAY ITEM AND (002210
14	ANAME(1)=NAME	002220
	ACNT=1	002230
	PARCNT=1	002240
17	NLEN=NLEN+1	002250
18	LEFT=(3-NLEN)*6	002260
	ANAME(ACNT)=SHIFT (STMT(P), LEFT).OR. (MASK(NLEN*6).AND. ANAME(ACNT))	002270
C	IF MATCHED PAREN FOUND, THE ARRAY IS COMPLETE-GO ON	002274
	IF (PARCNT.EQ.0) GO TO 500	002280
C	STORE AWAY REST OF ARRAY (UNTIL MATCHING PAREN FOUND)	002285
19	P=P+1	002290
	IF (STMT(P).EQ.BLNK) GO TO 19	002300
	IF (STMT(P).EQ.EOS) GO TO 500	002310
	IF (STMT(P).EQ.RPAREN) PARCNT=PARCNT-1	002320
	IF (STMT(P).EQ.LPAREN) PARCNT=PARCNT+1	002330
	IF (NLEN.LT.3) GO TO 17	002340
C	INCREMENT TO NEXT ARRAY STORAGE WORD AND ZERO THE COUNTER	002345
C	FOR LENGTH	002346
	IF (ACNT.EQ. MAXLEN) GO TO 500	002350
	ACNT=ACNT+1	002370
	NLEN=0	002380
	GO TO 13	002390
C	COMPUTE ARRAY LENGTH AND STORE IT ALL	002395
50	TLEN=10*(ACNT-1)+NLEN+1	002400
	CALL PUTNAME (NAME, TLEN, TYPE, NPTR)	002410
	CALL STPOL (NPTR)	002420
	RETURN N1	002430
	END	002440
	SUBROUTINE CLARIFY (STMT, P, TYPE), RETURNS (N1, N2)	002450
C	-N-EOS, N2-ILLEGAL	002460
C	CLARIFY TYPES INPUT NEXT CHARACTER IN STMT ARRAY	002461
C	AS A LETTER, DIGIT, OPERATOR, PARENTHESIS, SPECIAL POINT, OR COMMA	002462
C	BLANKS WITHIN A STMT ARE IGNORED	002463
	IMPLICIT INTEGER (A-Z)	002470
	DIMENSION STMT(1), CHAR(64)	002480
	DATA EOS/32(1,2), CHAR/9,20(1,2), 10(2,1), 4(3,1), 5(3,1), 8,10,6,7,15(9,1),	002490
	LAST/12(1,2), BLNK/12(1,2)	002500
C		002509
C	TYPE=1 FOR LETTERS	002510
C	TYPE=2 FOR NUMBERS	002511
C	TYPE=3 FOR OPERATORS	002512
C	TYPE=4 FOR (002513
C	TYPE=5 FOR)	002514
C	TYPE=6 FOR ,	002515
C	TYPE=7 FOR .	002516
C	TYPE=8 FOR =	002517

	TYPE=4 F22 OTHER (ILLEGAL CHARACTERS FOR AN ARITHMETIC STMT)	002519
C		002519
5	P=7+1	002520
	IF(STMT(P).EQ.F05) RETURN N1	002530
C	!(LAST) IS THE MAXIMUM ASCII VALUE ALLOWED IN PREPROCESSOR	002533
	IF(STMT(P).GT.LAST) GO TO 10	002540
C	BLANKS ARE IGNORED	002543
	IF(STMT(P).EQ.BLANK) GO TO 5	002550
	TYPE=CHAR(STMT(P)+1)	002550
C	CHECK FOR ILLEGAL CHARACTER	002570
	IF(TYPE.EQ.3) GO TO 10	002580
	RETURN	002590
10	TYPE=3	002600
	RETURN N2	002610
	END	002620
	SUBROUTINE STPOL (CHAR)	002630
	IMPLICIT INTEGER (A-Z)	002640
C	STPOL STOPS A VARIABLE NAME PTR (A LETTER) OR OPERATOR	002642
C	IN POLISH STRING ACCORDING TO REVERSE POLISH NOTATION AND	002643
C	AND THE OPERATOR PRECEDENCE ESTABLISHED.	002644
C	ROUTE ERROR MESSAGES TO A SEPARATE FILE (NOT THE ONE TO BE	002650
C	COMPILED)	002660
C	PUT EOS AT BEGINNING, EOS AT END	002670
	DIMENSION STACK(32,2),OPS(2,3)	002680
	COMMON/POL/POLISH(60,2)	002690
C	COMMON/POL/POLISH(60,2),NPOL??????LATER NOW JUUST DIMENSION**	002700
	COMMON/POL/POLISH(60,2)	002710
	DATA STKMAX/32/,POLMAX/60/,LPAREN/18(/,RPAREN/19/),EOS/33EOS/	002720
C	E= EXPONENTIATION /=UNARY ?=FUNCTION ,=FUNCTION ARGUMENT SEPARATOR	002730
	DATA OPS/12+,-10,18-,-10,18*,-9,18/, -9,18\$, -7,18!, -3,	002740
C	FIRST CHAR MUST BE A # OR .	002743
	A 18,, -12,18?, -17,18=, -13/	002750
	DATA LPRN/18/,RPRS/-50/,STACK(1,1)/-52/,FUNC/18?/,A_P4A/35/	002760
	DATA NSTK/1/,NPOL/1/	002770
	DATA IFYPE/18/,POLISH(1,1)/33EOS/	002780
C	NSTK & NPOL HAVE TO BE INITIALIZED SOMEWHERE	002790
C	*****POLISH(1,1)=EOS ALWAYS,REINITIALIZED TO NPOL=1 EACH TIME	002820
	IF (CHAR.EQ.F05) GO TO 15	002830
C	IF CHAR > ALPHA--ITS NOT AN ALPHANUMERIC CHARACTER (OR OPERAND)	002840
	IF(CHAR.EQ.IFYPE) GO TO 14	002850
	IF(CHAR.GT.35) GO TO 19	002860
C	GET AN OPERAND	002870
C	POLMAX = MAXIMUM LENGTH OF POLISH ARRAY	002874
13	IF(NPOL.GT.POLMAX) GO TO 510	002880
	NPOL=NPOL+1	002890
	POLISH(NPOL,1)=CHAR	002910
	POLISH(NPOL,2)=0	002910
	GO TO 100	002920
19	IF (CHAR.EQ.LPAREN) GO TO 40	002930
	IF (CHAR.EQ.RPAREN) GO TO 50	002940
C	TREAT FUNCTION PAREN LIKE A REGULAR PARENTHESIS	002950
	IF(CHAR.EQ.FUNC) GO TO 40	002960
C	GET AN OPERATOR	002970
	TO 20 IF=1,2	002980
	IF (CHAR.EQ.OPS (1,1)) GO TO 23	002990
20	CONTINUE	003000
	PRINT *, "ERROR--ILLEGAL OPERATION"	003020
	GOTO "167 13"	003030
C	MATCHED AN OPERATOR	003040
21	IF(OPS(2,1).GT.STACK(NSTK,1)) GO TO 35	003050
	IF (NPOL.GT.POLMAX) GO TO 500	003060
	NPOL=NPOL+1	003070
C	NEGATIVE VALUES FOR OPERATORS ARE STORED SO THAT THEY CAN	003075
C	EASILY BE DISTINGUISHED FROM OPERAND NAMES LATER WHEN	003076
C	IS DECOMPILED	003077
	STACK(NPOL,1)=STACK(NSTK,2)	003080

	POLISH(NPOL,2)=0	003090
	NSTK=NSTK-1	003100
	GO TO 23	003110
C	STORE AWAY THIS OPERATOR IN THE STACK	003120
35	IF (NSTK.GT.STKMAX) GO TO 500	003130
	NSTK=NSTK+1	003140
	STACK (NSTK,1)=OPS(2,I)	003150
	STACK (NSTK,2)=CHAR	003160
	GO TO 100	003170
C	GET LEFT PAREN	003180
40	IF (NSTK.GT.STKMAX) GO TO 500	003190
	NSTK=NSTK+1	003200
	STACK (NSTK,1)=LPRNK	003210
C	LPRNK SETS THE LEFT PAREN RANK PRECEDENCE INTO STACK	003213
	STACK(NSTK,2)=CHAR	003220
	GO TO 100	003230
C	GET RIGHT PAREN	003240
50	IF (NSTK.EQ.1) GO TO 500	003250
	IF (STACK(NSTK,2).EQ.LPAREN) GO TO 54	003260
C	DETECT END OF A FUNCTION TOO	003270
	IF (NPOL.GT.POLMAX) GO TO 500	003280
	NPOL=NPOL+1	003290
	POLISH(NPOL,1)=-STACK(NSTK,2)	003300
	POLISH(NPOL,2)=0	003310
C	---IF IT WAS A FUNCTION INDICATOR RETURN, INDICATOR ADDED TO POLISH	003320
	IF (STACK(NSTK,2).EQ.FUNC) GO TO 54	003330
	NSTK=NSTK-1	003340
C	---IF IT WAS A FUNCTION INDICATOR RETURN, INDICATOR ADDED TO POLISH	003350
	GO TO 50	003360
C	NOW REMOVE LEFT PAREN OR FUNCTION FROM STACK	003370
54	NSTK=NSTK-1	003380
	GO TO 100	003390
C	FOR PROCESSING-- CLEAR OPERATOR STACK	003400
75	IF (STACK(NSTK,1).EQ.POS) GO TO 90	003410
	IF (STACK (NSTK,2).EQ.LPAREN) GO TO 500	003420
	IF (NPOL.GT.POLMAX) GO TO 500	003430
	NPOL=NPOL+1	003440
	POLISH(NPOL,1)=-STACK(NSTK,2)	003450
	POLISH(NPOL,2)=0	003460
	NSTK=NSTK-1	003470
	GO TO 75	003480
90	NPOL=NPOL+1	003490
	IF (NPOL.GT.POLMAX) GO TO 500	003500
	POLISH(NPOL,1)=POS	003510
	POLISH(NPOL,2)=0	003520
100	CONTINUE	003530
	RETURN 0.	003560
50	PRINT *, " ERROR IN POLISH NSTK=", NSTK, " NPOL=", NPOL	003570
	GO TO 100	003580
	END	003590
	SUBROUTINE PUTNAME(IN,LEN,TYPE,NMPOS)	003600
	IMPLICIT INTEGER(A-Z)	003610
C	PUTNAME STORES A STRING OF CHARACTERS IN THE NAME ARRAY	003612
C	AND PASSES BACK TO SUPERORDINATE A POINTER TO THE	003613
C	POSITION THROUGH THE STRING WAS STORED IN THE NAME ARRAY.	003614
	COMMON/NAME/ NAME(100,7),LSTN	003620
	DIMENSION IN(1)	003630
	NMPOS=LSTN+1	003640
C	COMPUTE # OF WORDS REQUIRED TO STORE THE STRING	003650
	NMPOS=(LEN-1)/7+1	003660
C	STORE LENGTH OF STRING BESIDE THE 1ST WORD OF THE STRING	003665
	NAME(LSTN+1,0)=LEN	003670
	NAME(LSTN+1,1)=TYPE	003680
C	STORE REST OF STRING	003690
	DO 3 J=1,NMPOS	003700
	LET NAME(LSTN+1,J)=	003710

	DATA EOS/3REOF/,CONT/1R/	004240
	FRST=1	004250
	LST=72	004260
3	PRINT (1,5) (OUT(I),I=FRST,LST)	004270
5	FORMAT(7291,8X)	004280
	IF(OUT(LST+1).EQ.EOS) GO TO 30	004290
C	EOS EXPECTED IN COLUMN 73 OR 145 OR 217,ETC	004294
C	***OUT IN CONTINUATION MARK	004300
	OUT(LST+5)=CONT	004310
	LST=LST+72	004320
	FRST=FRST+72	004330
	GO TO 3	004340
30	RETURN	004350
	END	004360
	SUBROUTINE GETSTMT(STMT),RETURNS(N1)	004370
C	ALT RETURN N1 IS USED WHEN AN EOF HAS BEEN DETECTED	004380
C	GETSTMT READS INPUT PROGRAMS CARDS, PASSING A ' STATEMENT '	004382
C	AT A TIME TO ITS SUPERORDINATE(GETOBJ). A STATEMENT	004384
C	MAY SPAN SEVERAL CARDS THROUGH THE USE OF CONTINUATION CARDS,	004385
C	SO A LOOK-AHEAD READ IS REQUIRED.	004386
	IMPLICIT INTEGER (A-Z)	004390
	DIMENSION STMT(1)	004400
	COMMON/CONTS/LNCONT	004420
	COMMON/TEMP20/TEMP(73)	004430
	COMMON/NAMES/NAME(100,3),LSTN	004440
	DATA LSTCOL/72/,EOS/3RECS/,ZERO/100/,BLNK/1R /,ENDFILE/3REOF/,	004450
	* MAXCOLS/147/,LNCONT/1/	004460
C	USING NO INITIAL READ , JUST A WORD OF PLANKS	004470
1	LST=LSTCOL	004480
C	INITIALIZE LSTN FOR NAMES	004490
	LSTN=0	004500
C	CHECK FOR COMMENT CARD	004502
	IF(TEMP(1).EQ.100.OR.TEMP(1).EQ.1R+.OR.TEMP(1).EQ.1R3) GO TO 70	004510
	DO 2 I=1,LSTCOL	004520
2	STMT(I)=TEMP(I)	004530
4	CONTINUE	004550
C	READ INPUT PROGRAM CARD(COLS 1-72) INTO TEMP	004553
	STMT(2,3) (TEMP(I),I=1,LSTCOL)	004550
	FORMAT(7291)	004570
	IF(EOS(2).NE.0) GO TO 50	004580
	LNCONT=LNCONT+1	004590
C	CHECK FOR COMMENT CARDS-----	004600
	IF(TEMP(1).EQ.100.OR.TEMP(1).EQ.1R+.OR.TEMP(1).EQ.1R3) GO TO 30	004610
C	IF NOT A CONTINUED LINE GO TO 30	004620
	IF(TEMP(5).EQ.7500.OR.TEMP(5).EQ.3BLNK) GO TO 30	004630
C	GOT A CONTINUATION LINE	004640
C	MAXCOLS NOT PRESENT PREPROCESSOR IS 1440(2 CARDS)	004650
C	ALL STATEMENTS LONGER THAN THIS WILL NOT BE ANALYZED	004652
	IF(LAST.GE.MAXCOLS) GO TO 60	004660
	TEMP(6)=BLNK	004670
	DO 14 I=1,LSTCOL	004680
	LST=LST+1	004690
14	STMT(LAST)=TEMP(I)	004700
	GO TO 4	004710
C	PUT EOS MARKED IN LAST COLUMN+1 OF OUTGOING STATEMENT	004713
30	STMT(LAST+1)=EOS	004720
	RETURN	004730
C	IF HAS BEEN FOUND (RETURN N1)	004740
50	STMT(LAST+1)=EOS	004750
	STMT(LAST+2)=ENDFILE	004760
	RETURN N1	004770
C	GOT TOO LONG OF A STATEMENT--DO NOT PROCESS IT	004780
C	(OUTPUT STATEMENT THEN READ NEXT CARD UNTIL A NEW	004782
C	STATEMENT HAS BEEN DETECTED(NO MORE CONTINUED))	004783
60	STMT(LAST+1)=EOS	004790
	CALL PRINT(STMT)	004800

	NAME(LSTN,1)=IN(J)	003710
5	CONTINUE	003720
	RETURN	003730
	END	003740
	SUBROUTINE INIT77, RETURNS(N1)	003750
	IMPLICIT INTEGER (A-Z)	003760
C	INIT77 DOES THE INITIAL READ OF THE FORTRAN ALGORITHM	003764
C	IT READS THE CONTENTS OF THE FIRST PROGRAM CARD INTO	003765
C	A TEMPORARY ARRAY WHICH IS USED BY THE GET-STATEMENT	003765
C	MODULE, WHICH FUNCTIONS ON A LOOK-AHEAD BASIS.	003767
	COMMON/TEMP77/TEMP(73)	003770
	DATA TEMP/77*12,3REDS/	003780
	STAD(2,23) (TEMP(I),I=1,72)	003790
23	FORMAT(7291)	003800
C	CHECK FOR EMPTY FILE	003802
	IF(EOF(2).NE.0) RETURN N1	003810
	RETURN	003830
	END	003840
	SUBROUTINE GETOBJ(STMT,F), RETURNS(N1)	003850
C	ALT RETURN IS USED ONCE AN EOF HAS BEEN DETECTED	003860
C	GETOBJ SUPERSEDES THE CHECK OF A STATEMENT FOR A LEGAL OBJECT	003865
C	WHERE AN ARITHMETIC STATEMENT STRUCTURE CONSISTS OF	003867
C	OBJECT = ARITHMETIC EXPRESSION	003868
	IMPLICIT INTEGER (A-Z)	003870
	COMMON/TEMP77/TEMP(73)	003880
	DATA TEMP/77*12,3REDS/	003890
	STAD(2,23) (TEMP(I),I=1,72)	003900
	FORMAT(7291)	003910
C	INITIALIZE NRIL,NSTK - STKS FOR THE POLISH STRING,SINCE	003930
C	WE WILL BE BEGINNING A NEW ONE FOR EACH STATEMENT	003933
5	NRIL=NSTK=1	003940
	CALL GETSTMT(CTMT), RETURNS(605)	003950
	IF(CTMT.NE.0) GO TO 21	003960
C	CHECK FOR AN IF STMT TYPE, WHICH MAY CONTAIN AN ARITHMETIC	003980
C	STATEMENT TO BE EXECUTED UPON A TRUE CONDITION	003982
	IF(CTMT.EQ.1) GO TO 21	003990
	CALL CLASSIFY(CTMT,J), RETURNS(600,600)	004000
	IF(CTMT(J).EQ.1) GO TO 21	004010
15	CONTINUE	004020
C	GET AN IF STATEMENT, PLACE AN INDICATOR IN POLISH STRING	004030
	PARCNT=1	004040
C	SPACE OVER PREDICATE PORTION OF IF STATEMENT TO GET TO OBJECT	004060
20	IF(CTMT(J).EQ.1) PARCNT=PARCNT+1	004050
	IF(CTMT(J).EQ.2) PARCNT=PARCNT+1	004060
	IF(CTMT(J).EQ.3) GO TO 20	004070
	IF(PARCNT.NE.1) GO TO 20	004080
C	GO BACK UNTIL MATCHING RIGHT PAREN OF IF STMT PREDICATE	004090
C	IS FOUND	004095
	CALL STPOL(IFTYPE)	004100
	IF(IFTYPE.EQ.1) GO TO 21	004110
21	CALL ENDORJ(CTMT,J), RETURNS(6)	004120
	RETURN	004140
C	END FOUND NO ARITHMETIC STATEMENT	004143
201	CALL PRINT(CTMT)	004150
	GO TO 5	004170
C	FOR CONDITION-PRINT LAST STATEMENT (ASSUMED TO BE	004175
C	NON-ARITHMETIC (E.G. "END")	004177
202	CALL PRINT(CTMT)	004180
	RETURN N1	004190
	END	004200
	SUBROUTINE PRINT(OUT)	004210
	IMPLICIT INTEGER (A-Z)	004220
C	PRINT PRINTS OUT STATEMENT, WHICH MAY SPAN SEVERAL CARDS, TO	004223
C	THE UNIT DESIGNATION FROM AN OUTPUT FILE.	004224
	ATTENTION OUT(1)	004230

[illegible]

C	MATCH INPUT STATEMENT AGAINST KEYWORDS (BLANKS ARE IGNORED)	005354
17	P=P+1	005360
	IF (STMT(P).EQ.BLANK) GO TO 17	005370
	IF (STMT(P).NE.KEYWORDS(J,I)) GO TO 20	005380
15	CONTINUE	005390
C	I=1 INDICATES IT WAS A DIMENSION STATEMENT	005395
	IF (I.EQ.1) GO TO 300	005400
	GO TO 400	005410
20	CONTINUE	005420
C	IT DID NOT MATCH ANY KEYWORDS	005430
	GO TO 530	005440
300	CALL DIMENSN (STMT,P)	005450
	GO TO 530	005460
C	GOT A NEW ROUTINE-SO SAVE NAME AND INITIALIZE VIA NWRJTN	005470
400	CALL NWRJTN (STMT,P)	005480
530	CALL PRNT (STMT)	005490
	RETURN	005500
	END	005510
	SUBROUTINE NWRJTN (STMT,P)	005520
	IMPLICIT INTEGER (A-Z)	005530
C		005531
C	NWRJTN EXTRACTS THE PROGRAM UNITS NAME AND INITIALIZES THE	005532
C	LINE COUNTER FOR THAT NEW UNIT TO ONE. (THIS INFORMATION IF	005537
C	IS USED FOR EACH ARITHMETIC FUNCTION SUBROUTINE CALL BUILT	005544
C	BY THE DIFFERENT BRANCH). ALSO IT CLEARS THE LIST OF VARIABLE	005555
C	NAMES DECLARED AS ARRAYS FOR THE PREVIOUS PROGRAM UNIT.	005566
C		005577
C---	SAVES ROUTINE NAME AND REINITIALIZES LINE COUNT	005580
C		005589
	DIMENSION STMT(1)	005590
	COMMON/ARRAYS/ ARRAY(40), ARRYMAX	005600
	COMMON/GOTON/ GOTONUM	005610
	COMMON/ROUTNAM/ ROUTINE	005620
	COMMON/CONT/ LNCNT	005630
	DATA LPAREN/1R(/, RPAREN/1R(/, EOS/3F00/, BLANKS/10H	005640
	ROUTINE=BLANKS	005650
	CONT=0	005660
C	INITIALIZE GOTON # FOR THE NEW ROUTINE JUST DETECTED	005670
C	INITIALIZE LINE COUNT TO 1 FOR NEW ROUTINE	005680
C	CLEAR ARRAY COUNT SINCE THIS NEW ROUTINE DOESN'T HAVE	005690
C	ANY AS YET)	005700
	GOTONUM=10	005710
	LNCNT=1	005720
	ARRYMAX=0	005730
5	P=P+1	005740
C	LEFT PAREN INDICATES END OF ROUTINE NAME	005750
	IF (STMT(P).EQ.LPAREN) GO TO 80	005760
	IF (STMT(P).EQ.BLANK) GO TO 5	005770
	IF (STMT(P).EQ.EOS) GO TO 80	005780
C	STORE CHARACTER INTO ROUTINE NAME SAVE WORD	005790
	LEFT=(3-CONT)*5	005800
	ROUTINE=SHIFT (STMT(P), LEFT).OR. (SHIFT (MASK(54), LEFT).AND.ROUTINE)	005810
	CONT=CONT+1	005820
	GO TO 5	005830
80	RETURN	005840
	END	005850
	SUBROUTINE DIMENSN (STMT,P)	005860
C		005869
C	DIMENSN STORES ALL VARIABLE NAMES WHICH HAVE BEEN DECLARED	005875
C	AS ARRAYS BY THE DIMENSION STATEMENT	005886
C		005897
C	POINTER P IS EXPECTED TO BE SETTING AT LAST LETTER OF DIMENSION WORD	005908
	IMPLICIT INTEGER (A-Z)	005919
	COMMON/ARRAYS/ ARRAY (40), ARRYMAX	005930
	DIMENSION STMT(1)	005940
	DATA LPAREN/1R(/, RPAREN/1R(/, EOS/3F00/, BLANK/1R /,	005950


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      * COMMA/12,/,BLANKS/10H
4  MAY=BLANKS
   NCNT=0
   P=P+1
C  CHECK FOR BEGINNING OF DIMENSION
   IF(STMT(P).EQ.LPAREN) GO TO 63
   IF(STMT(P).EQ.BLNK) GO TO 5
   IF(STMT(P).EQ.EOS) GO TO 80
   LEFT=(9-NCNT)*5
C  STORE AWAY DIMENSIONED VARIABLE NAME (ONE CHARACTER AT A TIME)
   N1=SHIFT(STMT(P),LEFT).OR.(MASK(NCNT*5).AND.NAM)
   NCNT=NCNT+1
   GO TO 5
C  HAS BEEN DETECTED
C WE HAVE FOUND A NEW APRAY SO SAVE IT THEN PROCESS ITS ARGUMENTS
50  PCNT=1
   ARRYMAX=ARRYMAX+1
   APRAY(ARRYMAX)=NAM
C  PROCESSING UNTIL WE FIND RIGHT PAREN
61  P=P+1
   IF(STMT(P).EQ.RPAREN) GO TO 67
C  IF EOS -PROBABLY ERROP STOP PROCESSING
   IF(STMT(P).EQ.EOS) GO TO 80
   IF(STMT(P).EQ.LPAREN) PCNT=PCNT+1
   GO TO 61
67  IF(PCNT.EQ.1) GO TO 74
   PCNT=PCNT-1
   GO TO 51
C COMMA SEPARATES ARRY DECLARATIONS
74  P=P+1
   IF(STMT(P).EQ.COMMA) GO TO 4
   IF(STMT(P).EQ.EOS) GO TO 80
   GO TO 74
80  RETURN
END
SUBROUTINE SETUP
  IMPLICIT INTEGER (A-Z)
  SETUP SOME POLISH STRING LEFT TO RIGHT AND CALLS A HANDLING
  ROUTINE FOR EACH OPERATOR FOUND.
  DIMENSION NAME(5),TARRAY(5),OBJECT(5)
  DIMENSION OPS(9),FILLER(3)
  DIMENSION CNT(2)
  COMMON/GETON/GETONUM
  COMMON/CNTS/LCNT
  COMMON/ROUTINE/ ROUTINE
  COMMON/POL/POLISH(60,2)
  DATA ITEMCH1/'HTTTTTTAA/',ITEMCH2/'HRTTTTAA',EOL/'-44',NEG/'1H-/
C  CHECK OF OPS IS + - * / ** = , ? :
  DATA OPS/'-37,-38,-39,-40,-43,-44,-46,-57,-63/
  DATA ISTYP/'R',GETONUM/10/,CNT/10H CNT, 10HINE /,
  * CNTLEN/14/
  ITEMCH=ITEMCH1
  ITEMCH2=ITEMCH2
C  FNL IS A FLAG WHICH INDICATES THE LAST OPERATION PERFORMED
C  IN THE ARITHMETIC STATEMENT (DESCRIBED BY THE POLISH NOTATION)
C  FNL=)
  FILLLEN=21
C  FILL STANDARD FOLLOWER FOR ALL SUB CALLS DERIVED FROM THE
C  POLISH STRING OF A SINGLE ARITHMETIC STATEMENT
  ENCODE(FILLLEN,1,FILLLEN) ROUTINE,LNCNT,FNL
1  FIRST(1H,7H,1H,1H,1H,1H,1H,1H,1H,1H,1H,1H,1H,1H,1H,1H,1H,1H,1H,1H,1H,1H)
2  P=1
C  -----OPERATORS ARE CODED TO BE LESS THAN 1 IN POLISH STRING
5  IF(POLISH(P).LT.0) GO TO 15
  IF (POLISH (P).EQ. 32665) GO TO 666

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	P=P+1	006410
	GO TO 5	006420
C	FIND NEXT P	006430
15	NP=P+1	006440
17	IF (POLISH (NP,1). LT. C) GO TO 20	006450
C	----- CHECK FOR "=" (SINGLE ASSIGNMENT INDICATION)	006460
	IF (POLISH (NP,1). EQ. 3REOS) GO TO 40	006470
	NP=NP+1	006480
	GO TO 17	006490
20	IF (POLISH (NP,1). EQ. FOL) FNL=1	006500
	ENDCOS (FLLWLEN,1,FLLWER) ROUTINE,LNONT,FNL	006510
	GO TO 70	006520
C	***** "=" SHOULD ALWAYS SET FNL=1 AND STOP IT ALL	006530
C	---SHOULD BE AN EQUAL SIGN IF IT IS THE ONLY OPERATOR IN STRING	006540
40	IF (POLISH (P,1). EQ. FOL) GO TO 42	006550
	PRINT*, "ERROR-NO EQUAL"	006560
	RETURN	006570
C	SET UP A SINGLE ASSIGNMENT	006580
42	CALL GETRND (POLISH,P,RND)	006590
	CALL GETNAME (POLISH (RND,1),NAME,TYPE,LEN)	006600
C	NOW STORE THE OBJECT	006610
	CALL GETRND (POLISH,RND,OBJ)	006620
	CALL GETNAME (POLISH (OBJ,1),OBJECT,OTYPE,OBJLEN)	006630
C	-- NO NEED TO STORE NAME SINCE IT IS THE LAST ONE	006640
	CALL SHGLASS (NAME,LEN,TYPE,OBJECT,OBJLEN,FLLWER,FLLWLEN)	006650
C	--- NOW CLEAR POLISH	006660
	POLISH (OBJ,1)=	006670
	POLISH (RND,1)=	006680
C	--- GO GET NEXT POLISH STRING OPERATOR	006690
	GO TO 240	006700
70	DO 72 I=1,3	006710
	IF (POLISH (P,1). EQ. OPS (I)) GO TO 74	006720
72	CONTINUE	006730
C	***** THE FOLLOWING IS TEMPORARY *****	006740
	PRINT *, "BAD OPERATOR"	006750
	RETURN	006760
74	IF (I.EQ.5) GO TO 230	006770
	IF (I.EQ.7) GO TO 117	006780
	IF (I.EQ.4) GO TO 85	006790
	IF (I.EQ.3) GO TO 125	006800
	RETURN	006810
C	GOT 2 FINISH FUNCTION	006820
85	CALL FINSENO (POLISH,P,ITEMPN,RTMPN,FNL,FLLWER,FLLWLEN)	006830
	GO TO 241	006840
C	GOT "," BUILD PART OF A FUNCTION	006850
C	*****CALL FUNCOLD (POLISH,P) *****	006860
117	CONTINUE	006870
	CALL FUNCOLD (POLISH,P,ITEMPN,RTMPN,FLLWER,FLLWLEN)	006880
	GO TO 240	006890
C	GOT UNARY	006900
125	CALL GETRND (POLISH,P,RND)	006910
	CALL GETNAME (POLISH (RND,1),NAME,TYPE,LEN)	006920
	TARRAY (1)=NEG	006930
	TLEN=1	006940
	CALL MERGE (TARRAY,TLEN,NAME,LEN)	006950
	CALL PUTNAME (TARRAY,TLEN,TYPE,PTR)	006960
	POLISH (RND,1)=PTR	006970
	GO TO 241	006980
230	CALL GETPTR (POLISH,P,ITEMPN,RTMPN,FNL,FLLWER,FLLWLEN)	006990
C	GO TO GET NEXT OPERATOR	007000
240	IF (FNL.EQ.1) GO TO 507	007010
241	DEF	007020
	GO TO 5	007030
507	CONTINUE	007040
C	*****CHECK FOR IF FOLLOWING*****	007050
	IF (POLISH (P,1). NE. IFTYPE) GO TO 525	007060
		007070
		007080
		007090
		007100
		007110
		007120
		007130
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		009110
		009120
		009130
		009140
		009150
		009160
		009170
		009180
		009190
		009200
		009210
		009220
		009230
		009240
		009250
		009260
		009270
		009280
		009290
		009300
		009310
		009320
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		009370
		009380
		009390
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		009970
		009980
		009990
		010000

	END000(CONTLEN,79,CONT) GOTO NUM	007100
90	FORMAT(3H 77,12,9H CONTINUE)	007110
	CALL PRINTF(CONT,CONTLEN)	007120
525	CONTINUE	007130
	RETURN	007140
C	LEVEL 500 MAY WANT TO SET A FLAG OR ALTERNATE RETURN *****	007160
	END	007170
	SUBROUTINE FINSFNC(POLISH,P,ITEMPN,ITEMPN,FNL,FLLWER,FLLWLEN)	007180
	IMPLICIT INTEGER(A-Z)	007190
C		007192
C	FINSFNC MERGES THE FUNCTION NAME WITH ITS ARGUMENT LIST	007193
C	(BUILT BY FUNORLD) AND GENERATES A SINGLE ASSIGNMENT	007194
C	CALL FOR THE FUNCTION.	007195
C		007199
	DIMENSION POLISH(50,2),FARGS(20),FNAM(20),OBJECT(2)	007200
	DIMENSION FLLWER(1)	007210
	DATA LPAREN/1L(/,RPAREN/1L)/,EVAL/1/	007220
	DATA FARGS/20*10H /,FNAM/20*10H	007230
C	GET FUNCTION ARGUMENTS	007240
	CALL GETPND(POLISH,P,PND)	007250
	IF(POLISH(PND,2).EQ.EVAL) GO TO 15	007260
C	---- GET A TEMPORARY OBJECT FOR THE SINGLE ARGUMENT ----	007270
	CALL GETNAME(POLISH(PND,1),FARGS,TYPE,FLEN)	007280
	ORLEN=7	007290
	IF(TYPE.EQ.2) GO TO 7	007300
	OBJECT(1)=ITEMPN	007310
	ITEMPN=ITEMPN+1000000	007320
	GO TO 11	007330
7	OBJECT(1)=ITEMPN	007340
	ITEMPN=ITEMPN+1000000	007350
11	CALL SNGLARG(FARGS,FLEN,TYPE,OBJECT,ORLEN,FLLWER,FLLWLEN)	007360
	FARGS(1)=OBJECT(1)	007370
	ORLEN=ORLEN	007380
	GO TO 17	007390
15	CALL GETNAME(POLISH(PND,1),FARGS,DUMMY,ORLEN)	007400
C	HOW TO GET FUNCTION NAME	007410
17	CALL GETPND(POLISH,PND,FNC)	007420
	CALL GETNAME(POLISH(FNC,1),FNAM,TYPE,FLEN)	007430
	CALL MERGE(FNC,1,FLEN,LPAREN,1)	007440
	CALL MERGE(FNC,1,FLEN,FARGS,ORLEN)	007450
	CALL MERGE(FNC,1,FLEN,RPAREN,1)	007460
	POLISH(PND,1)=	007470
	IF(FNL.EQ.1) GO TO 50	007480
	ORLEN=7	007490
	IF(TYPE.EQ.2) GO TO 31	007500
	OBJECT(1)=ITEMPN	007510
	ITEMPN=ITEMPN+1000000	007520
	GO TO 40	007530
31	OBJECT(1)=ITEMPN	007540
	ITEMPN=ITEMPN+1000000	007550
40	CONTINUE	007560
	ORTYPE=TYPE	007570
	CALL SNGLARG(FNC,FLEN,ORTYPE,OBJECT,ORLEN,FLLWER,FLLWLEN)	007580
C	*** 2 BELOW IS STOPPING AWAY PT2 TO THE WHOLE FUNCTION ***	007590
C	--- WAS THAT THE LAST ASSIGNMENT (FNL=1)? -- IF SO RETURN	007600
	CALL PUTNAME(OBJECT,ORLEN,TYPE,PTR)	007610
	POLISH(FNC,1)=PTR	007620
	GO TO 50	007630
50	CALL GETPND(POLISH,PND,LIST)	007640
	CALL GETNAME(POLISH(LIST,1),OBJECT,DUMMY,ORLEN)	007650
	POLISH(LIST,2)=EVAL	007660
	CALL SNGLARG(FNAM,FLEN,TYPE,OBJECT,ORLEN,FLLWER,FLLWLEN)	007670
60	RETURN	007680
	END	007690
	SUBROUTINE GETPND (PND, PND, TYPE, LND)	007700
	IMPLICIT INTEGER (A-Z)	007710

C	GETNAME GETS A STRING OF CHARACTERS FROM THE NAME ARRAY	007711
C	ACCORDING TO THE POINTER GIVEN AND PASSES TO THE CALLER	007713
C	THE STRING, ITS DATA TYPE, AND ITS LENGTH	007714
C		007715
C	COMMON/NAMES/ NAME(100,3),LSTN	007716
	OTENSION OUT(1)	007720
	LEN=NAME (NPOS,2)	007730
	TYPE=NAME (NPOS,3)	007740
	NWPOS= (LEN-1)/10+1	007750
	DO 5 J=1, NWPOS	007760
	OUT(J) = NAME (NPOS,1)	007770
5	NPOS=NPOS+1	007780
	RETURN	007790
	END	007800
	SUBROUTINE PRINTIT (OUT, LEN)	007810
C		007820
C	PRINTS 80 CHARACTER CARD IMAGES TO THE N-BIT SIMULATED	007823
C	PROGRAM OUTPUT FILE	007824
C	TAKES CARE OF CONTINUATIONS	007825
C	IMPLICIT INTEGER (A-Z)	007830
	DIMENSION OUT(1)	007840
	DATA CONTINU/10H * /	007850
	NWPOS=(LEN-1)/10+1	007860
	J=1	007870
	K=7	007880
	IF(NWPOS.LT.7) K=NWPOS	007890
C	PRINT CARD IMAGE TO OUTPUT FILE	007900
	PRINT(1,5) (OUT(I),I=J,K)	007910
5	FORMAT ('410')	007920
	IF(NWPOS.LE.7) GO TO 20	007930
	K=K+5	007940
C	PRINT SUCCEEDING CARD IMAGES AS CONTINUED STATEMENT	007950
	DO 14 J=8, NWPOS, 5	007954
	IF(K.GT.NWPOS) K=NWPOS	007960
	PRINT (1,5) CONTINU, (OUT(I), I=J,K)	007970
	J=J	007980
	K=K+5	007990
14	CONTINUE	008000
	K=K-5	008010
C	CLEAR OUT SO IT WON'T BE PRINTED ON SUCCEEDING PRINTS	008020
C	OF GREATER THAN ONE CARD	008022
20	DO 23 LL=J,K	008030
23	OUT(LL)=10H	008040
	RETURN	008050
	END	008060
	SUBROUTINE FUNCOLD(COLINH,P,TEMPNM1,TEMPNM2,FLLWEP,FLLWLEN)	008070
	IMPLICIT INTEGER (A-Z)	008080
C		008090
C	FUNCOLD HELPS FUNCTION ARGUMENTS TOGETHER, GENERATING	008091
C	CALLS FOR SINGLE ARGUMENTS IF NECESSARY.	008092
C		008093
	DIMENSION FLLWEP(1)	008099
	OTENSION FLLWEP(1),NAME(20),NAME2(20)	008100
	DATA EVAL/1/, COM1/1H, /	008110
	DATA NAME/20*1H /,NAME2/20*1H /	008120
C	LEN1 AND LEN2 REPRESENT LENGTH OF TEMPORARY NAME	008123
	LEN1=LEN2=7	008130
	CALL GET2(COLINH, P, END)	008140
C	IF ITS BOTH TERMINAL- GO ON, OTHERWISE CREATE A SINGLE	008150
C	STATEMENT CALL FOR THE ARGUMENT (SO THAT IS TAKES	008152
C	IN THE N-BIT SIGNIFICANCE).	008153
C		008154
	IF (COLINH(1),1) EQ. EVAL) GO TO 40	008160
C	GET FUNCTION ARGUMENT	008167
	CALL GETNAME (COLINH(1), NAME, TYPE, LEN)	008170
	IF TYPE EQ. 1) GO TO 30	008175

	TEMPN=TEMPN2	008210
C	INCREMENT TO NEXT REAL TEMPORARY VARIABLE NAME	008212
	TEMPN2=TEMPN2+10000000	008220
	GO TO 29	008230
C	INCREMENT TO NEXT TEMPORARY INTEGER VARIABLE NAME	008231
22	TEMPN=TEMPN1	008240
	TEMPN1=TEMPN1+10000000	008250
29	CONTINUE	008260
C	SET UP CALL FOR STORING AWAY THE TEMPORARY OBJECT NAME	008263
	CALL SNGCLASS(NAME, LEN, TYPE, TEMPN, TLEN2, FLOWER, FLLWLEN)	008270
	CALL PUTNAME(TEMPN, TLEN2, TYPE, PTR)	008280
	POLISH (RND, 1)=PTR	008290
	POLISH (RND, 2)=EVAL	008300
C	SET OTHER FUNCTION ARGUMENT	008302
43	CALL GETPND (POLISH, RND, RND1)	008310
C	SEE IF IT HAS ALREADY BEEN EVALUATED	008320
	IF (POLISH (RND1, 2) .EQ. EVAL) GO TO 83	008330
C	IT HASN'T BEEN EVALUATED, SO EVALUATE IT WITH A SUB CALL	008334
	CALL GETNAME (POLISH (RND1, 1), NAME, TYPE, LEN)	008340
C	CREATE TEMPORARY VARIABLE NAMES	008342
	IF (TYPE.EQ.2) GO TO 45	008350
	TEMPN=TEMPN2	008360
	TEMPN2=TEMPN2+10000000	008370
	GO TO 49	008380
45	TEMPN=TEMPN1	008390
	TEMPN1=TEMPN1+10000000	008400
49	CONTINUE	008410
C	EVALUATE ARGUMENT WITH SINGLE ASSIGNMENT CALL	008414
	CALL SNGCLASS(NAME, LEN, TYPE, TEMPN, TLEN1, FLLWLEN)	008420
	CALL PUTNAME(TEMPN, TLEN1, TYPE, PTR)	008430
	POLISH (RND1, 1)=PTR	008440
	POLISH (RND1, 2)=EVAL	008450
C	NOW TAKE EVALUATED ARGUMENTS AND MERGE THEM TOGETHER	008453
C	TO BUILD PART OF THE FUNCTION ARGUMENT LIST	008455
53	CALL GETNAME (POLISH (RND1, 1), NAME, TYPE, LEN)	008460
	CALL GETNAME (POLISH (RND, 1), NAME2, TYPE, LEN2)	008470
	CALL MERGE(NAME, LEN, NAME2, LEN2)	008480
	CALL MERGE(NAME, LEN, NAME2, LEN2)	008490
	CALL PUTNAME(NAME, LEN1, TYPE, PTR)	008500
C	NAME HAS ALREADY BEEN MARKED AS EVALUATED AND ZERO OUT ARG 2	008510
	POLISH (RND1, 1)=PTR	008520
	POLISH (RND, 1)=0	008530
	RETURN	008540
	END	008550
	ROUTINE SNGCLASS (OPRND, OPRLEN, TYPE, OBJECT, OBJLEN, FLLWLEN)	008560
	IMPLICIT, INTEGER(A-Z)	008570
C		008573
C	SNGCLASS MERGES PARTS REQUIRED FOR A SINGLE ASSIGNMENT	008575
C	FUNCTION CALL	008576
C		008577
	DEFINITION OPER (1), OBJECT (1), OUT (1), FLLWLEN (1)	008580
	DATA IASGN/4=IASGN(1, IASGN/4=IASGN(1, IOP/64	008590
	DATA OUT/64=10	008600
	OUT (1)=10	008610
	OUT=6	008620
C	OUT THE OBJECT	008630
	CALL MERGE (OUT, OLEN, OBJECT, OBJLEN)	008640
	OUT=IASGN	008650
	IF (TYPE.EQ.2) SUB=IASGN	008660
	OUT=OUT	008670
	CALL MERGE (OUT, OLEN, SUB, SUBLEN)	008680
	CALL MERGE (OUT, OLEN, OPRND, OPRLEN)	008690
	CALL CHECKTH (OUT, OUT)	008700
	CALL MERGE (OUT, OLEN, FLLWLEN, FLLWLEN)	008710
	CALL PRINT (OUT, OLEN)	008720
	RETURN	008730

	END	008740
	SUBROUTINE GETRND(POLISH,LP,NXT)	008750
	IMPLICIT INTEGER(A-Z)	008760
	DIMENSION POLISH(1,1)	008770
C		008772
C	GETRND FETCHES (FROM THE POLISH STRING) THE LAST OPERAND	008773
C	AND PASSES BACK A POINTER TO THAT OPERAND	008774
C		008776
	NXT=LP	008780
5	NXT=NXT-1	008790
	IF(POLISH(NXT).GT.0) RETURN	008800
	GO TO 5	008810
	END	008820
	SUBROUTINE MERGE(OUT,OLEN,IN,ILEN)	008830
	IMPLICIT INTEGER(A-Z)	008840
C		008850
C	MERGES STRING ON RIGHT TO THE RIGHT END OF STRING ON THE LEFT	008853
C	OLEN WILL THEN ASSUME THE TOTAL LENGTH OF THE RESULTING STRING	008854
C	AND OUT WILL CONTAIN THE MERGED STRING.	008855
C		008856
	DIMENSION OUT(1),IN(1)	008857
C	WOULD NOT NEED TO INITIALIZE OUT THIS WAY	008860
	OWRDS=(OLEN-1)/10+1	008860
	NWRDS=(ILEN-1)/10+1	008870
	IF(OLEN-OWRDS*10.GT.0) OWRDS=OWRDS+1	008880
	IF(ILEN-NWRDS*10.GT.0) NWRDS=NWRDS+1	008890
C	GET A PARTIAL WORD- COMPUTE WHAT IS LEFT IN OUT (# CHRS)	008900
	MULTPS=OWRDS*10	008910
	DO 30 I=1,NWRDS	008920
	DO 30 J=1,OWRDS	008930
	OUT(MULTPS+J)=IN(I)	008940
	MULTPS=OWRDS*10	008950
30	FORMAT(A=,=(A10))	008960
	GO TO 70	008970
C	TEST FOR OLEN = 0 COVERS CASE WHERE OUT ARRAY WAS EMPTY	008980
66	IF(OLEN.EQ.0) OWRDS=0	008990
C	FILL IN WHOLE WORDS	009000
	DO 37 I=1,NWRDS	009010
	OWRDS=OWRDS+1	009020
67	OUT(OWRDS*10)=IN(I)	009030
70	OLEN=OLEN+ILEN	009040
	RETURN	009050
	END	009060
	SUBROUTINE GENRND(POLISH,P,ITEMN,ITEMN,FUL,FLLWR,FLLWLEN)	009070
	IMPLICIT INTEGER(A-Z)	009080
C		009090
C	GENRND AVAILABLE AND MERGES PARTS OF A SUBROUTINE CALL	009100
C	INVOLVING TWO OPERANDS.	009110
C		009120
	DIMENSION OPND1(5),OPND2(5),OUT(40),FLLWR(1),FLLWLEN(1),	009130
	POLISH(51,2),OBJECT(2),OPS(2,9)	009140
	DATA BLANKS/101 /,FLLWLEN/77,	009150
	OPND1/14,/,FLLWR/77,OPS/2205,OPND2/1/	009160
	DATA SUBN/74=1477(,74=2100(,74=2100(,74=2100(,	009170
	74=2100(,74=2100(,74=2100(,74=2100(,	009180
	74=2100(,74=2100(,74=2100(,74=2100(,	009190
	74=2100(,74=2100(,74=2100(,74=2100(,	009200
	DATA OPS/17=-17,17,-17,17,-17,17,-17,17,-17,17,-17,17,-17,	009210
	17,-17,17,-17,17,-17,17,-17,17,-17,17,-17,17,-17,	009220
	DATA OUT/40=101 /,OPND1/50=101 /	009230
	DATA OPND2/50=101 /	009240
	DO 1 I=1,5	009250
	OUT(I)=0	009260
1	IF(OPS(I).EQ.0) GO TO SUBROUTINE TYPE	009270
C	DO 15 I=1,5	009280
	IF(OPS(I).EQ.0) POLISH(I) GO TO 17	009290
15	CONTINUE	009300

Line	Code	Statement	Address
1	C	PRINT "+, " R00 OP R=", R	003290
2	C	RETURN	003300
3	C	SN=I	003310
4	C	GET OPERAND 2	003320
5	C	CALL GETRND(POLISH,R,RND2)	003330
6	C	GETNAME STORES THE NAME INTO OPND2 ARRAY	003340
7	C	CALL GETNAME(POLISH(RND2,1),OPRND2,TYPE2,LEN2)	003350
8	C	GET OPERAND 1 LOCATION IN POLISH	003360
9	C	CALL GETRND(POLISH,RND2,RND1)	003370
10	C	CALL GETNAME(POLISH(RND1,1),OPRND1,TYPE1,LEN1)	003380
11	C	MAKE A SSIFT	003390
12	C	MODE=(TYPE2-1)+2*(TYPE1-1)+1	003400
13	C	FETCH SUBROUTINE NAME	003410
14	C	SN1=SUBN(MODE,SN)	003420
15	C	IF(FNL.EQ.1) GO TO 70	003430
16	C	IF(MODE.LE.3) GO TO 60	003440
17	C	OBJECT(1)=ITEMPN	003450
18	C	ITEMPN=ITEMPN+10000009	003460
19	C	TYPE=2	003470
20	C	GO TO 65	003480
21	C	IFEN A REAL OBJECT	003490
22	C	OBJECT(1)=ITEMPN	003500
23	C	ITEMPN=ITEMPN+10000003	003510
24	C	TYPE=1	003520
25	C	STORE AWAY NEW OBJECTS NAME	003530
26	C	CALL PUTNAME(OBJECT,TOLFN,TYPE,NPOS)	003540
27	C	POLISH(RND1,1)=NPOS	003550
28	C	POLISH(RND1,2)=EVAL	003560
29	C	CLR OPERAND 2 FROM POLISH STRING ALSO	003570
30	C	POLISH(RND2,1)=0	003580
31	C	POLISH(RND2,2)=0	003590
32	C	OLLEN=TOLFN	003600
33	C	GO TO 75	003610
34	C	LAST ONE--GET FINAL OBJECT	003620
35	C	CALL GETRND(POLISH,RND1,LAST)	003630
36	C	CALL GETNAME(POLISH(LAST,1),OBJECT,TYPE,OLLEN)	003640
37	C	***** JUST IN CASE GET EVAL IN THIS OBJECT*****	003650
38	C	POLISH(LAST,2)=EVAL	003660
39	C	OUT(1)=BLANKS	003670
40	C	FIRST LENGTH WILL BE THE FIRST 6 BLANKS	003680
41	C	OLEN=6	003690
42	C	CALL MERGE(OUT,OLEN,OBJECT,OLLEN)	003700
43	C	CALL MERGE(OUT,OLEN,OPR,OLLEN)	003710
44	C	CALL MERGE(OUT,OLEN,OPRND1,LEN1)	003720
45	C	CALL MERGE(OUT,OLEN,OPRND2,LEN2)	003730
46	C	CALL CHECKTH(OLEN,OUT)	003740
47	C	CALL MERGE(OUT,OLEN,FULLPR,FULLLEN)	003750
48	C	CALL PRINT(OUT,OLEN)	003760
49	C	RETURN	003770
50	C	END	003780
51	C	SUBROUTINE PREPRT(SMT)	003790
52	C	EXPLICIT INTEGER(A-7)	003800
53	C	PREPRT OUTPUTS A COMMENT STATEMENT CONTAINING THE INPUT	003810
54	C	STATEMENT AND OUTPUTS PRELIMINARY OUTPUT FOR LABELS AND	003820
55	C	IF STATEMENTS (WHEN NECESSARY)	003830
56	C	DEFINITION SMT(1),OUT(1-5),IFTYPE(3),GOTO(73),CONT(73)	003840
57	C	CHECK FOR LABELS AND PRINTS THEM	003850
58	C	IF HAVE GOT AN EXPRESSION AT THIS POINT	003860
59	C	CHECK FOR THE EXPRESSION SMT	003870
60	C	CONT(73) SMT GO TO 100	003880
61	C	CALL DEFN(1-7,CONT(73),OUT(73),CONT(73)+1,100,120,140,120)	003890
62	C	SMT(1),OUT(1-5),IFTYPE(3),GOTO(73),CONT(73)+1	003900

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DATA IFTYPE/1R,1R,1R(//,RPAREN/1R)//,LPAREN/1R(//,GOTO/3*1R,1R,6,992J
* 1R0,1R,1R,1R0,1R,1R7,1R7,1R1,1R,1R,53*1R,32R03/ 009930
00 4 I=1,5 009930
IF(STMT(I).NE.BLNK) GO TO 30 009950
CONTINUE 00996J
C --- NO LABELS - JUST COMMENT IT 009970
GO TO 35 009980
30 00 33 I=1,5 009990
33 CONT(I)=STMT(I) 010000
CALL PRNT(CONT) 010010
36 LST=73 010020
37 IF(STMT(LST).EQ.COS) GO TO 33 010030
STMT(LST)=CONT 010040
LST=LST+72 010050
GO TO 37 010060
39 STMT(I)=CONT 010070
00 39 I=2,5 010080
39 STMT(I)=BLNK 010090
CALL PRNT(STMT) 010100
C CHECK FOR AN IF STMT --USE SAME ROUTINE***** 010110
40 J=0 010120
00 43 I=1,3 010130
CALL CLASIFY(STMT,J,TYPE),RETURNS(500,500) 010140
IF(STMT(J).NE.IFYPE(I)) GO TO 500 010150
CONTINUE 010160
C ---- GOT AN IF STMT ---- 010170
C NOW PASSOVER IF(=====)--- 010180
PARCNT=1 010190
45 J=J+1 010200
IF(STMT(J).EQ.RPAREN) PARCNT=PARCNT-1 010210
IF(STMT(J).EQ.LPAREN) PARCNT=PARCNT+1 010220
IF(PARCNT.GT.0) GO TO 45 010230
J=J+1 010240
00 46 I=1,5 010260
45 CONT(I)=BLNK 010270
00 47 K=7,J 010280
47 CONT(K)=STMT(K) 010290
C TO COVER IF COVERING > 1 STMT (LINE) 010300
NINE=((J/7)+1)*72 010310
C -- FILL IN REST OF LINE WITH BLANK, MAKE GO TO PORTION A CONTINUE 010320
00 50 I=J,NINE 010330
50 CONT(I)=BLNK 010340
CONT(NINE+1)=COS 010350
CALL PRNT(OUT) 010360
GOTOIN=GOTOIN+1 010370
GOTO(5)=CONTIN 010380
ENDDEF(2,67,GOUT)GOTOIN 010390
FORMAT(I2) 010400
C GOT TO CONVERT BINARY NUMBER TO ALPHANUMERIC 010410
ENDDEF(2,68,GOUT) (GOTO(I),I=19,10) 010420
68 FORMAT(21,F1) 010430
CALL PRNT(GOTO) 010440
C --- CREATE THE CONTINUE LINKED WITH THE IF GO TO 010450
K=1 010460
00 70 I=10,10 010470
K=K+1 010480
70 CONT(K)=GOTO(I) 010490
GOTO(5)=BLNK 010500
GOTOIN=GOTOIN+1 010510
ENDDEF(2,67,GOUT) GOTOIN 010520
ENDDEF(2,68,GOUT) (GOTO(I),I=19,16) 010530
CALL PRNT(GOUT) 010540
C THE CONTINUE FOR EXPRESSION ---- 010550
CALL PRNT(OUT) 010560
END 010570
END 010580

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	SUBROUTINE CHECK7H(LEN,OUT)	010600
	IMPLICIT INTEGER(A-Z)	010610
C		010612
C	CHECK7H VOIES SPECIAL CASE TO ASSURE "7H" DOES NOT OVERLAP	010613
C	LINES AND CAUSE A SYNTAX ERROR. IT FILLS REMAINING PORTION	010614
C	OF THE LINE WITH BLANKS(SHIFTING THE REST DOWN TEN SPACES)	010615
	DIMENSION OUT(1)	010620
	DATA BLANKS/104	010630
	IF(LEN.LE.50) GO TO 99	010640
	IF(LEN.GT.50.AND.LEN.LT.70) GO TO 5	010650
	IF(MOD(LEN,50).LT.10) GO TO 5	010660
	GO TO 99	010670
5	CALL MERGE(OUT,LEN,BLANKS,10)	010680
99	RETURN	010690
	END	010700

N-bit Simulation Subroutines

Twenty-two n-bit simulation subroutines were built to handle each combination of data types and arithmetic operations. In addition, there were four other subroutines developed (SETNBIT, ROUNDER, OVERFLW, and UNDRFLW).

The SETNBIT subroutine was developed to compute seven key values upon which the other n-bit subroutines base their n-bit wordlength simulation effects. The OVERFLW and UNDRFLW subroutines are used to print overflow messages. A subroutine called ROUNDER is used whenever a special case results where a round up increments the exponent of the floating point value (due to the fixed point add performed to round floating point values).

The 22 n-bit simulation subroutines basically are the same. They can be divided into two classes: those which produce real values and those which produce integer values. After the operation is performed in each subroutine, the rest of the routine is eventually the same as each of the others within its own class.

The source listing of the n-bit simulation subroutines follows:

N-bit Simulation Subroutines
Source Listing

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SUBROUTINE SETNBIT(NBITS,IRNDTR,IPOCSN,MESSGS,MANTSA,IEXPNT,
* IPTPOS,KEY)
C
C THIS ROUTINE ESTABLISHES THE MIN/MAX VALUES THAT COULD BE
C REPRESENTED IN FIXED AND FLOATING PT AS OF THE GIVEN N-BIT
C SPECIFICATIONS-ROUTINE TERMINATES PROGRAM FOR ILL
C
C DIMENSION KEY(7)
C DATA MANTLEN/48/
C IF((MANTSA+IEXPNT+1).NE.NBITS) GO TO 31
C IF(NBITS.GT.60) GO TO 30
C IF(IRNDTR.LT.0) GO TO 32
C *****NOTE THESE CANT USE REGULAR UNITS--DEPENDS ON USER PROGRAM****
C MIGHT WANT INTEGER MAX TO BE NBITS+2 TO MAKE UP FOR SIGN BIT
C KEY(1)=SHIFT(MASK(1),NBITS)-1
C IF(IPOCSN.GT.3.OR.IPOCSN.LT.1) GO TO 33
C IF(IEXPNT.GT.11) GO TO 35
C IF(MANTSA.EQ.MANTLEN.AND.IRNDTR.NE.0) GO TO 34
C MAXMANT=MANTSA+(IPOCSN-1)*NBITS
C IF(MAXMANT.GT.MANTLEN) GO TO 34
C KEY(4)=IRNDTR
C GET MESSAGE PRINT/SUPPRESSION FLAG
C KEY(5)=MESSGS
C KEY(6)=MANTLEN-MANTSA
C KEY(7)=MANTLEN-MAXMANT
C *****NEED 60 BIT ALLOWANCE IN THIS ONE AS OF 17 NOVEMBER****
C IF(IEXPNT.EQ.11.AND.MANTSA.EQ.48) GO TO 25
C IRNGE=SHIFT(MASK(1),IEXPNT)-1
C BIT=2.**IRNGE
C IF(IPTPOS.EQ.0) BIG=2.** (IRNGE-MANTSA)
C PMAX=SHIFT(MASK(MANTSA),MANTSA)
C PMAX=PMAX*BIG
C PMAX=PMAX*2.** (SHIFT(MASK(MANTLEN),MANTLEN))
C PMAX=SHIFT(SHIFT(PMAX,-KEY(7)),KEY(7))
C SHIFTING PMAX DISABLES THE MODE CONVERSION TO INTEGER FOR KEY(2)
C KEY(2)=SHIFT(PMAX,0)
C MAXM=SHIFT(MASK(1),MANTSA)
C PRR=MAXM
C PMIN=PRR*2.** (-IRNGE)
C IF(IPTPOS.EQ.1) PMIN=PRR*2.** (-IRNGE-MANTSA)
C KEY(3)=SHIFT(PMIN,0)
C RETURN
C SPECIAL CASE FOR 60 BIT WORD MAX,ITS REAL MAX WOULD ROUND THE JOB
C WHEN ANY COMPARSES WERE MADE, SO ITS MADE SLIGHTLY LESS
C CONTINUE
C KEY(2)=37767777777777777777
C KEY(3)=00014000000000000000
C GO TO 27
C 60 PRINT *," TOO MANY BITS SPECIFIED: N-BIT MAX IS 60 ",
C " --- ",NBITS," WERE SPECIFIED"
C GO TO 40
C 71 PRINT *," MANTISSA + EXPONENT + 1 SHOULD EQUAL 'N'-BITS"
C GO TO 40
C 72 PRINT *," ILLEGAL ROUNDING/TRUNCATION OPTION VALUE ",
C " ROUNDING = 1/ TRUNCATION = 0"
C GO TO 40
C 73 PRINT *," ILLEGAL AMOUNT OF PRECISION SPECIFIED "
C PRINT *," ( 1=SINGLE PRECISION, 2=DOUBLE, 3=TRIPLE)"
C GO TO 40
C 74 PRINT *," *** NUMBER OF BITS SPECIFIED FOR (MANTISSA*PRECISION)"
C " GREATER THAN MAXIMUM PROGRAM CAN HANDLE"

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GO TO 40
35 PRINT *, " EXPONENT TOO LARGE FOR CDC-- > 11 BITS"
40 PRINT *, " ERROR- ILLEGAL SETBIT DATA PARAMETER FOUND"
STOP " BAD INPUT TO SETBIT ROUTINE"
END
INTEGER FUNCTION IADD(I1,I2,ROUTINE,LNCNT,IFML,KEY)
DIMENSION KEY(7)
IADD=I1+I2
IF(IABS(IADD).GT.KEY(1)) GO TO 100
5 RETURN
100 IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCNT,8HADDITION)
IF(IADD.GT.0) IADD=KEY(1)
IF(IADD.LT.0) IADD=-KEY(1)
GO TO 5
END
SUBROUTINE OVERFLW(ROUTINE,LNCNT,OPER)
IMPLICIT INTEGER (A-Z)
C ***** MAY WANT TO CHANGE THIS PRINT FROM A FORMATTED *****
PRINT(1,4) ROUTINE,LNCNT,OPER
4 FORMAT(" **** OVERFLOW IN ",A8," LINE # ",I5," FOR ",I10," ****")
RETURN
END
SUBROUTINE UNDERFLW(ROUTINE,LNCNT,OPER)
IMPLICIT INTEGER (A-Z)
PRINT(1,4) ROUTINE,LNCNT,OPER
4 FORMAT(" **** UNDERFLOW IN ",A8," LINE # ",I5," FOR ",I10," ****")
RETURN
END
REAL FUNCTION PRADD(R1,R2,ROUTINE,LNCNT,IFNL,KEY)
DIMENSION KEY(7)
PRADD=R1+R2
IF(KEY(4).NE.0) GO TO 30
25 IF(ABS(PRADD).GT.SHIFT(KEY(2),0)) GO TO 100
IF(ABS(PRADD).LT.SHIFT(KEY(3),0).A.PRADD.NE.0) GO TO 110
C CHECK FOR ROUNDING
25 PRADD=SHIFT(SHIFT(PRADD,-KEY(7)),KEY(7))
IF(IFNL.EQ.1) PRADD=SHIFT(SHIFT(PRADD,-KEY(6)),KEY(5))
RETURN
C GO ROUND
30 INC=1
IF(PRADD.LT.0.) INC=-1
PRADD=SHIFT(SHIFT(SHIFT(PRADD,-(KEY(7)-1))+INC,-1),KEY(7))
IF(SHIFT(SHIFT(PRADD,12),-12).EQ.0) CALL ROUNDER(PRADD)
GO TO 25
100 IF(PRADD.GT.0.) PRADD=SHIFT(KEY(2),0)
IF(PRADD.LT.0.) PRADD=-SHIFT(KEY(2),0)
C CHECK FOR MESSAGES
IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCNT,8HADDITION)
GO TO 25
110 PRADD=0
IF(KEY(5).NE.0) CALL UNDERFLW(ROUTINE,LNCNT,8HADDITION)
GO TO 25
END
REAL FUNCTION PRADD(R1,R2,ROUTINE,LNCNT,IFNL,KEY)
DIMENSION KEY(7)
PRADD=R1+I2
IF(KEY(4).NE.0) GO TO 30
25 IF(ABS(PRADD).GT.SHIFT(KEY(2),0)) GO TO 100
IF(ABS(PRADD).LT.SHIFT(KEY(3),0).A.PRADD.NE.0) GO TO 110
25 PRADD=SHIFT(SHIFT(PRADD,-KEY(7)),KEY(7))
IF(IFNL.EQ.1) PRADD=SHIFT(SHIFT(PRADD,-KEY(6)),KEY(5))
RETURN
C GO ROUND IF NON
30 INC=1
IF(PRADD.LT.0.) INC=-1
PRADD=SHIFT(SHIFT(SHIFT(PRADD,-(KEY(7)-1))+INC,-1),KEY(7))
IF(SHIFT(SHIFT(PRADD,12),-12).EQ.0) CALL ROUNDER(PRADD)
GO TO 25
100 IF(PRADD.GT.0.) PRADD=SHIFT(KEY(2),0)
IF(PRADD.LT.0.) PRADD=-SHIFT(KEY(2),0)
C CHECK FOR MESSAGES
IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCNT,8HADDITION)
GO TO 25
110 PRADD=0
IF(KEY(5).NE.0) CALL UNDERFLW(ROUTINE,LNCNT,8HADDITION)
GO TO 25
END

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      R2ADD=SHIFT(R2ADD,12),-12).EQ.0)  CALL ROUNDER(R2ADD)
      GO TO 25
100 IF(R2ADD.GT.0.)      R2ADD=SHIFT(KEY(2),0)
      IF(R2ADD.LT.0.)      R2ADD=-SHIFT(KEY(2),0)
      IF(KEY(5).NE.0)      CALL OVERFLW(ROUTINE,LNCNT,8H10000000)
      GO TO 23
110 R2ADD=0
      IF(KEY(5).NE.0)      CALL UNDRFLW(ROUTINE,LNCNT,8H10000000)
      GO TO 23
      END
      REAL FUNCTION R1ADD(I1,P2,ROUTINE,LNCNT,IFNL,KEY)
      DIMENSION KEY(7)
      R1ADD=I1+P2
      IF(KEY(4).NE.0)      GO TO 30
      IF(ABS(R1ADD).GT.SHIFT(KEY(2),0))      GO TO 100
      IF(ABS(R1ADD).LT.SHIFT(KEY(3),0).A.R1ADD.NE.0)      GO TO 110
28 R1ADD=SHIFT(SHIFT(R1ADD,-KEY(7)),KEY(7))
      IF(IFNL.EQ.1)      R1ADD=SHIFT(SHIFT(R1ADD,-KEY(6)),KEY(6))
      RETURN
C      GO ROUND IT NOW
30 INC=1
      IF(R1ADD.LT.0.)      INC=-1
      R1ADD=SHIFT(SHIFT(SHIFT(R1ADD,-(KEY(7)-1))+INC,-1),KEY(7))
      IF(SHIFT(SHIFT(R1ADD,12),-12).EQ.0)  CALL ROUNDER(R1ADD)
      GO TO 25
110 IF(RRADD.GT.0.)      RRADD=SHIFT(KEY(2),0)
      IF(RRADD.LT.0.)      RRADD=-SHIFT(KEY(2),0)
      IF(KEY(5).NE.0)      CALL OVERFLW(ROUTINE,LNCNT,8H10000000)
      GO TO 28
110 R1ADD=0
      IF(KEY(5).NE.0)      CALL UNDRFLW(ROUTINE,LNCNT,8H10000000)
      GO TO 28
      END
      INTEGER FUNCTION IIMPY(I1,I2,ROUTINE,LNCNT,IFNL,KEY)
      DIMENSION KEY(7)
      IIMPY=I1*I2
      IF(ABS(IIMPY).GT.KEY(1))      GO TO 100
      RETURN
5      CALL OVERFLW(ROUTINE,LNCNT,8H10000000)
120 IF(IIMPY.GT.0)      IIMPY=KEY(1)
      IF(IIMPY.LT.0)      IIMPY=-KEY(1)
      GO TO 5
      END
      INTEGER FUNCTION IIMNS(I1,I2,ROUTINE,LNCNT,IFNL,KEY)
      DIMENSION KEY(7)
      IIMNS=I1-I2
      IF(ABS(IIMNS).GT.KEY(1))      GO TO 100
      RETURN
5      CALL OVERFLW(ROUTINE,LNCNT,8H10000000)
100 IF(IIMNS.GT.0)      IIMNS=KEY(1)
      IF(IIMNS.LT.0)      IIMNS=-KEY(1)
      GO TO 5
      END
      INTEGER FUNCTION IIOVD(I1,I2,ROUTINE,LNCNT,IFNL,KEY)
      DIMENSION KEY(7)
      IIOVD=I1/I2
      IF(ABS(IIOVD).GT.KEY(1))      GO TO 100
      RETURN
5      CALL OVERFLW(ROUTINE,LNCNT,8H10000000)
110 IF(IIOVD.GT.0)      IIOVD=KEY(1)
      IF(IIOVD.LT.0)      IIOVD=-KEY(1)
      GO TO 5
      END
      INTEGER FUNCTION IISXP(I1,I2,ROUTINE,LNCNT,IFNL,KEY)
      DIMENSION KEY(7)
      IISXP=I1**I2

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	IF(IABS(IIEXP).GT.KEY(1)) GO TO 100	002030
5	RETURN	002040
100	IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCNT,8HEXPONENT)	002050
	IF(IIEXP.GT.0) IIEXP=KEY(1)	002060
	IF(IIEXP.LT.0) IIEXP=-KEY(1)	002070
	GO TO 5	002080
	END	002090
	REAL FUNCTION RMPY(R1,P2,ROUTINE,LNCNT,IFNL,KEY)	002100
	DIMENSION KEY(7)	002110
	RMPY=R1*P2	002120
	IF(KEY(4).NE.0) GO TO 30	002130
25	IF(ABS(RMPY).GT.SHIFT(KEY(2),0)) GO TO 100	002140
	IF(ABS(RMPY).LT.SHIFT(KEY(3),0).A.RMPY.NE.0) GO TO 110	002150
28	RMPY=SHIFT(SHIFT(RMPY,-KEY(7)),KEY(7))	002160
	IF(IFNL.EQ.1) RMPY=SHIFT(SHIFT(RMPY,-KEY(5)),KEY(5))	002170
	RETURN	002180
C	GO ROUND	002190
30	INC=1	002200
	IF(RMPY.LT.0) INC=-1	002210
	RMPY=SHIFT(SHIFT(SHIFT(RMPY,-(KEY(7)-1)+INC,-1),KEY(7))	002220
	IF(SHIFT(SHIFT(RMPY,12),-12).EQ.0) CALL ROUNDER(RMPY)	002230
	GO TO 25	002240
100	IF(RMPY.GT.0) RMPY=SHIFT(KEY(2),0)	002250
	IF(RMPY.LT.0) RMPY=-SHIFT(KEY(2),0)	002260
C	CHECK FOR MESSAGES	002270
	IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCNT,8MULTIPLY)	002280
	GO TO 28	002290
110	RMPY=0	002300
	IF(KEY(5).NE.0) CALL UNDFLW(ROUTINE,LNCNT,8MULTIPLY)	002310
	GO TO 28	002320
	END	002330
	REAL FUNCTION RPNMS(R1,P2,ROUTINE,LNCNT,IFNL,KEY)	002340
	DIMENSION KEY(7)	002350
	RPNMS=R1/P2	002360
	IF(KEY(4).NE.0) GO TO 30	002370
25	IF(ABS(RPNMS).GT.SHIFT(KEY(2),0)) GO TO 100	002380
	IF(ABS(RPNMS).LT.SHIFT(KEY(3),0).A.RPNMS.NE.0) GO TO 110	002390
28	RPNMS=SHIFT(SHIFT(RPNMS,-KEY(7)),KEY(7))	002400
	IF(IFNL.EQ.1) RPNMS=SHIFT(SHIFT(RPNMS,-KEY(5)),KEY(5))	002410
	RETURN	002420
C	GO ROUND	002430
30	INC=1	002440
	IF(RPNMS.LT.0) INC=-1	002450
	RPNMS=SHIFT(SHIFT(SHIFT(RPNMS,-(KEY(7)-1)+INC,-1),KEY(7))	002460
	IF(SHIFT(SHIFT(RPNMS,12),-12).EQ.0) CALL ROUNDER(RPNMS)	002470
	GO TO 25	002480
100	IF(RPNMS.GT.0) RPNMS=SHIFT(KEY(2),0)	002490
	IF(RPNMS.LT.0) RPNMS=-SHIFT(KEY(2),0)	002500
C	CHECK FOR MESSAGES	002510
	IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCNT,8SUBTRACT)	002520
	GO TO 28	002530
110	RPNMS=0	002540
	IF(KEY(5).NE.0) CALL UNDFLW(ROUTINE,LNCNT,8SUBTRACT)	002550
	GO TO 28	002560
	END	002570
	REAL FUNCTION RPDV(R1,P2,ROUTINE,LNCNT,IFNL,KEY)	002580
	DIMENSION KEY(7)	002590
	RPDV=R1/P2	002600
	IF(KEY(4).NE.0) GO TO 30	002610
25	IF(ABS(RPDV).GT.SHIFT(KEY(2),0)) GO TO 100	002620
	IF(ABS(RPDV).LT.SHIFT(KEY(3),0).A.RPDV.NE.0) GO TO 110	002630
28	RPDV=SHIFT(SHIFT(RPDV,-KEY(7)),KEY(7))	002640
	IF(IFNL.EQ.1) RPDV=SHIFT(SHIFT(RPDV,-KEY(5)),KEY(5))	002650
	RETURN	002660
C	GO ROUND	002670
30	INC=1	002680

	IF(RRQVD.LT.0.) INC=-1	002696
	RRQVD=SHIFT(SHIFT(SHIFT(RRQVD,-(KEY(7)-1))+INC,-1),KEY(7))	002700
	IF(SHIFT(SHIFT(RRQVD,12),-12).EQ.0) CALL ROUNDER(RRQVD)	002710
	GO TO 26	002720
100	IF(RRQVD.GT.0.) RRQVD=SHIFT(KEY(2),0)	002730
	IF(RRQVD.LT.0.) RRQVD=-SHIFT(KEY(2),0)	002740
C	CHECK FOR MESSAGES	002750
	IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCNT,8HODIVISION)	002760
	GO TO 28	002770
110	RRQVD=0	002780
	IF(KEY(5).NE.0) CALL UNDERFLW(ROUTINE,LNCNT,8HODIVISION)	002790
	GO TO 28	002800
	END	002810
	REAL FUNCTION R2MNS(R1,I2,ROUTINE,LNCNT,IFNL,KEY)	002820
	DIMENSION KEY(7)	002830
	R2MNS=R1-I2	002840
C	CHECK FOR ROUNDING	002850
	IF(KEY(4).NE.0) GO TO 30	002860
26	IF(ABS(R2MNS).GT.SHIFT(KEY(2),0)) GO TO 100	002870
	IF(ABS(R2MNS).LT.SHIFT(KEY(3),0).A.R2MNS.NE.0) GO TO 110	002880
28	R2MNS=SHIFT(SHIFT(R2MNS,-KEY(7)),KEY(7))	002890
	IF(IFNL.EQ.1) R2MNS=SHIFT(SHIFT(R2MNS,-KEY(6)),KEY(6))	002900
	RETURN	002910
C	GO ROUND IT NOW	002920
30	INC=1	002930
	IF(R2MNS.LT.0.) INC=-1	002940
	R2MNS=SHIFT(SHIFT(SHIFT(R2MNS,-(KEY(7)-1))+INC,-1),KEY(7))	002950
	IF(SHIFT(SHIFT(R2MNS,12),-12).EQ.0) CALL ROUNDER(R2MNS)	002960
	GO TO 26	002970
100	IF(R2MNS.GT.0.) R2MNS=SHIFT(KEY(2),0)	002980
	IF(R2MNS.LT.0.) R2MNS=-SHIFT(KEY(2),0)	002990
	IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCNT,4HSUBTRACT)	003000
	GO TO 28	003010
110	R2MNS=0	003020
	IF(KEY(5).NE.0) CALL UNDERFLW(ROUTINE,LNCNT,4HSUBTRACT)	003030
	GO TO 28	003040
	END	003050
	REAL FUNCTION R1MNS(I1,F2,ROUTINE,LNCNT,IFNL,KEY)	003060
	DIMENSION KEY(7)	003070
	R1MNS=I1-F2	003080
C	CHECK FOR ROUNDING OPTION	003090
	IF(KEY(4).NE.0) GO TO 30	003100
26	IF(ABS(R1MNS).GT.SHIFT(KEY(2),0)) GO TO 100	003110
	IF(ABS(R1MNS).LT.SHIFT(KEY(3),0).A.R1MNS.NE.0) GO TO 110	003120
28	R1MNS=SHIFT(SHIFT(R1MNS,-KEY(7)),KEY(7))	003130
	IF(IFNL.EQ.1) R1MNS=SHIFT(SHIFT(R1MNS,-KEY(6)),KEY(6))	003140
	RETURN	003150
C	GO ROUND IT NOW	003160
30	INC=1	003170
	IF(R1MNS.LT.0.) INC=-1	003180
	R1MNS=SHIFT(SHIFT(SHIFT(R1MNS,-(KEY(7)-1))+INC,-1),KEY(7))	003190
	IF(SHIFT(SHIFT(R1MNS,12),-12).EQ.0) CALL ROUNDER(R1MNS)	003200
	GO TO 26	003210
100	IF(R1MNS.GT.0.) R1MNS=SHIFT(KEY(2),0)	003220
	IF(R1MNS.LT.0.) R1MNS=-SHIFT(KEY(2),0)	003230
	IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCNT,4HSUBTRACT)	003240
	GO TO 28	003250
110	R1MNS=0	003260
	IF(KEY(5).NE.0) CALL UNDERFLW(ROUTINE,LNCNT,4HSUBTRACT)	003270
	GO TO 28	003280
	END	003290
	REAL FUNCTION R2MNY(R1,I2,ROUTINE,LNCNT,IFNL,KEY)	003300
	DIMENSION KEY(7)	003310
	R2MNY=R1+I2	003320
C	CHECK FOR ROUNDING	003330
	IF(KEY(4).NE.0) GO TO 30	003340


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25 IF(ABS(R2MPY).GT.SHIFT(KEY(2),0)) GO TO 100
IF(ABS(R2MPY).LT.SHIFT(KEY(3),0).A.R2MPY.NE.0) GO TO 110
28 R2MPY=SHIFT(SHIFT(R2MPY,-KEY(7)),KEY(7))
IF(IFNL.EQ.1) R2MPY=SHIFT(SHIFT(R2MPY,-KEY(6)),KEY(5))
RETURN
C GO ROUND IT NOW
30 INC=1
IF(R2MPY.LT.0.) INC=-1
R2MPY=SHIFT(SHIFT(SHIFT(R2MPY,-(KEY(7)-1))+INC,-1),KEY(7))
IF(SHIFT(SHIFT(R2MPY,12),-12).EQ.0) CALL ROUNDER(R2MPY)
GO TO 25
100 IF(R2MPY.GT.0.) R2MPY=SHIFT(KEY(2),0)
IF(R2MPY.LT.0.) R2MPY=-SHIFT(KEY(2),0)
IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCHT,8MULTIPLY)
GO TO 25
110 R2IPY=0
IF(KEY(5).NE.0) CALL UNDRFLW(ROUTINE,LNCHT,8MULTIPLY)
GO TO 25
END
REAL FUNCTION R1MPY(I1,F2,ROUTINE,LNCHT,IFNL,KEY)
DIMENSION KEY(7)
R1MPY=I1*F2
C CHECK FOR ROUNDING
IF(KEY(4).NE.0) GO TO 30
25 IF(ABS(R1MPY).GT.SHIFT(KEY(2),0)) GO TO 100
IF(ABS(R1MPY).LT.SHIFT(KEY(3),0).A.R1MPY.NE.0) GO TO 110
28 R1MPY=SHIFT(SHIFT(R1MPY,-KEY(7)),KEY(7))
IF(IFNL.EQ.1) R1MPY=SHIFT(SHIFT(R1MPY,-KEY(6)),KEY(5))
RETURN
C GO ROUND IT NOW
30 INC=1
IF(R1MPY.LT.0.) INC=-1
R1MPY=SHIFT(SHIFT(SHIFT(R1MPY,-(KEY(7)-1))+INC,-1),KEY(7))
IF(SHIFT(SHIFT(R1MPY,12),-12).EQ.0) CALL ROUNDER(R1MPY)
GO TO 25
100 IF(R1MPY.GT.0.) R1MPY=SHIFT(KEY(2),0)
IF(R1MPY.LT.0.) R1MPY=-SHIFT(KEY(2),0)
IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCHT,8MULTIPLY)
GO TO 25
110 R1MPY=0
IF(KEY(5).NE.0) CALL UNDRFLW(ROUTINE,LNCHT,8MULTIPLY)
GO TO 25
END
REAL FUNCTION R2DVO(R1,I2,ROUTINE,LNCHT,IFNL,KEY)
DIMENSION KEY(7)
R2DVO=R1/I2
C CHECK FOR ROUNDING
IF(KEY(4).NE.0) GO TO 30
25 IF(ABS(R2DVO).GT.SHIFT(KEY(2),0)) GO TO 100
IF(ABS(R2DVO).LT.SHIFT(KEY(3),0).A.R2DVO.NE.0) GO TO 110
28 R2DVO=SHIFT(SHIFT(R2DVO,-KEY(7)),KEY(7))
IF(IFNL.EQ.1) R2DVO=SHIFT(SHIFT(R2DVO,-KEY(6)),KEY(5))
RETURN
C GO ROUND IT NOW
30 INC=1
IF(R2DVO.LT.0.) INC=-1
R2DVO=SHIFT(SHIFT(SHIFT(R2DVO,-(KEY(7)-1))+INC,-1),KEY(7))
IF(SHIFT(SHIFT(R2DVO,12),-12).EQ.0) CALL ROUNDER(R2DVO)
GO TO 25
100 IF(R2DVO.GT.0.) R2DVO=SHIFT(KEY(2),0)
IF(R2DVO.LT.0.) R2DVO=-SHIFT(KEY(2),0)
IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCHT,8DIVISION)
GO TO 25
110 R2DVO=0
IF(KEY(5).NE.0) CALL UNDRFLW(ROUTINE,LNCHT,8DIVISION)
GO TO 25

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```


	END	004000
	REAL FUNCTION R10VD(I1,I2,ROUTINE,LNCNT,IFNL,KEY)	004010
	DIMENSION KEY(7)	004020
	R10VD=I1/I2	004030
C	CHECK FOR ROUNDING	004040
	IF(KEY(4).NE.0) GO TO 30	004050
26	IF(ABS(R10VD).GT.SHIFT(KEY(2),0)) GO TO 100	004060
	IF(ABS(R10VD).LT.SHIFT(KEY(3),0).A.R10VD.NE.0) 30 TO 110	004070
28	R10VD=SHIFT(SHIFT(R10VD,-KEY(7)),KEY(7))	004080
	IF(IFNL.EQ.1) R10VD=SHIFT(SHIFT(R10VD,-KEY(5)),KEY(5))	004090
	RETURN	004100
C	GO ROUND IT NOW	004110
30	INC=1	004120
	IF(R10VD.LT.0.) INC=-1	004130
	R10VD=SHIFT(SHIFT(SHIFT(R10VD,-(KEY(7)-1))+INC,-1),KEY(7))	004140
	IF(SHIFT(SHIFT(R10VD,12),-12).EQ.0) CALL ROUNDER(R10VD)	004150
	GO TO 25	004160
100	IF(R10VD.GT.0.) R10VD=SHIFT(KEY(2),0)	004170
	IF(R10VD.LT.0.) R10VD=-SHIFT(KEY(2),0)	004180
	IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCNT,8HODIVISION)	004190
	GO TO 25	004200
110	R10VD=0	004210
	IF(KEY(5).NE.0) CALL UNDRFLW(ROUTINE,LNCNT,8HODIVISION)	004220
	GO TO 25	004230
	END	004240
	REAL FUNCTION R2EXP(R1,I2,ROUTINE,LNCNT,IFNL,KEY)	004250
	DIMENSION KEY(7)	004260
	R2EXP=R1**I2	004270
	IF(KEY(4).NE.0) GO TO 30	004280
26	IF(ABS(R2EXP).GT.SHIFT(KEY(2),0)) GO TO 100	004290
	IF(ABS(R2EXP).LT.SHIFT(KEY(3),0).A.R2EXP.NE.0) GO TO 110	004300
28	R2EXP=SHIFT(SHIFT(R2EXP,-KEY(7)),KEY(7))	004310
	IF(IFNL.EQ.1) R2EXP=SHIFT(SHIFT(R2EXP,-KEY(6)),KEY(5))	004320
	RETURN	004330
C	GO ROUND IT NOW	004340
30	INC=1	004350
	IF(R2EXP.LT.0.) INC=-1	004360
	R2EXP=SHIFT(SHIFT(SHIFT(R2EXP,-(KEY(7)-1))+INC,-1),KEY(7))	004370
	IF(SHIFT(SHIFT(R2EXP,12),-12).EQ.0) CALL ROUNDER(R2EXP)	004380
	GO TO 25	004390
100	IF(R2EXP.GT.0.) R2EXP=SHIFT(KEY(2),0)	004400
	IF(R2EXP.LT.0.) R2EXP=-SHIFT(KEY(2),0)	004410
	IF(KEY(5).NE.0) CALL OVERFLW(ROUTINE,LNCNT,8HEXPONENT)	004420
	GO TO 25	004430
110	R2EXP=0	004440
	IF(KEY(5).NE.0) CALL UNDRFLW(ROUTINE,LNCNT,8HEXPONENT)	004450
	GO TO 25	004460
	END	004470
	FUNCTION IASGN(I1,ROUTINE,LNCNT,IFNL,KEY)	004480
	DIMENSION KEY(7)	004490
	IASGN=I1	004500
	IF(I1.GT.KEY(1)) IASGN=KEY(1)	004510
	IF(I1.LT.-KEY(1)) IASGN=-KEY(1)	004520
	RETURN	004530
	END	004540
	FUNCTION PARCN(I1,ROUTINE,LNCNT,IFNL,KEY)	004550
	DIMENSION KEY(7)	004560
	PARCN=I1	004570
	IF(ABS(PARCN).GT.SHIFT(KEY(2),0)) PARCN=SHIFT(KEY(2),0)	004580
	IF(ABS(PARCN).LT.SHIFT(KEY(3),0).A.PARCN.NE.0) PARCN=0	004590
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004600
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004610
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004620
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004630
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004640
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004650
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004660
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004670
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004680
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004690
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004700
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004710
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004720
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004730
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004740
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004750
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004760
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004770
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004780
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004790
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004800
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004810
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004820
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004830
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004840
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004850
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004860
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004870
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004880
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004890
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004900
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004910
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004920
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004930
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004940
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004950
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004960
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004970
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004980
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	004990
	IF(PARCN.LT.0.) PARCN=-SHIFT(KEY(2),0)	005000

Appendix D

Proposed In-line Modification

A proposed change to the existing n-bit simulation tool would replace all subroutine calls generated by the present preprocessor program with in-line code which would perform the n-bit wordlength effects. Although this modification would result in large memory requirements for the resulting preprocessed program, the execution time should be reduced by one-third. The in-line code would be tailored to the user options in effect for that particular program module. This would mean fewer decisions would be made than in the present simulation tool, during execution of the n-bit simulated version of the algorithm. Also, subroutine linkage time would be save since the n-bit wordlength effects were no longer performed through calls to subroutines.

This modification would require the detection of the CALL SETNBIT subroutine for each program unit. The NON-EXPRESSION-HANDLER module could be modified to detect CALL SETNBIT, then call a handling routine which could be developed to analyze the user options expressed as arguments in SETNBIT. The same SETNBIT subroutine which is used presently could compute the key values and pass them via COMMON to the subroutines of the preprocessor which output the n-bit simulation code for the arithmetic expressions.

The in-line code which could perform the n-bit simulation effects for the various arithmetic operations and operand data types is shown in Figure D-1. The underlined words in the figure would be filled in by the handling routine with each n-bit simulation situation.

```

For Real
Results:  object = varb1 operation varb2
If Round
Option   { object = SHIFT(SHIFT(SHIFT(object ,-(key7 -1)))+
is on    (-1**(SHIFT(object .AND. MASK(1),1)),
          -1),key7 )
          IF(SHIFT(SHIFT(object ,12),-12).eq.0)
              CALL RCUNDER(object)
          IF( ABS(object).GT.key2 .OR.(ABS(object).LT.key3
              .AND.ABS(object).NE.0)
              CALL OVRFLWR(object ,operation ,routine ,
                  line# ,key2 )
          object = SHIFT(SHIFT(object ,-key6/7), key6/7)
          -----
For Integer
Results:  object = varb1 operation varb2
          IF(IABS(object).GT.key1)
              CALL OVRFLWI(object ,operation ,routine ,
                  line# ,key1)

```

Figure D-1 Possible In-line Code For Real and Integer Arithmetic Operation Results

To output the code shown in Figure D-1, the GETPRTS and SNGLASG modules would have to be modified slightly. Specifically, the merge operations of each module would be replaced by the actual arithmetic operation build followed by a sub-routine call to a new module. This new module would have standard outputs (shown by the non-underlined portions of Figure D-1). Key values (computed by SETNBIT) would be placed in the standard positions. The calling module would have to pass to this new module the object name, object data type, and the final assignment flag. In addition, the new routine must

be passed values computed by the SETNBIT subroutine. In Figure D-1, if the final flag was turned on, key6 would be the value positioned in the last statement. If this was not the final assignment key7 would be placed there, so a multi-precision mantissa would be saved when the in-line code is executed.

Vita

Gary A. Klein was born on 18 June 1951 in Newton, Iowa, the son of Clifford M. Klein and Doris G. Klein. After graduating from Newton High School in 1969, he attended Iowa State University in Ames, Iowa. In May, 1973 he graduated with distinction from Iowa State with a Bachelor of Science degree in Computer Science and received his commission as a graduate of the ROTC program. One month after graduation, he entered the Air Force as a Second Lieutenant and was stationed at Offutt Air Force Base for three years as a computer programmer. In May, 1976 he entered graduate studies in Electrical Engineering at the Air Force Institute of Technology.

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The objective of this thesis was to process algorithms written in FORTRAN on the Control Data Corporation (CDC) 6600/ CYBER 74 computer systems such that the results obtained were similar to those obtainable on an n-bit machine. The problem of n-bit simulation was addressed from two levels: the assembly language level and the FORTRAN language level. Each level involved designing a preprocessor which would modify the algorithm's code so that the numerical effects of n-bit wordlength could be realized. The assembly language (COMPASS) level approach failed, however, an n-bit simulation tool was successfully developed and implemented at the FORTRAN level. Several user options were incorporated into the n-bit simulation tool to enable the user to simulate the characteristics of various computer types.

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